Programming with Coq

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August 2009

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In this class, we introduce the basic commands for defining new values and programs. We shall present :

- The defining commands
- The typing discipline : how to construct a well-formed value

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A collection of basic types and functions

The Definition command

```
Attach a name to an expression
```

```
Definition three := 3.
three is defined
```

Verify that an expression is well-formed

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Check three. three : nat

Compute a value

```
Eval compute in three.
= 3 : nat
```

Defining functions

Expressions that depend on a variable

Check 2 + 3. 2 + 3 : nat

Check 5 + 3.

5 + 3 : nat

Definition add3 (x : nat) := x + 3. add3 is defined

The type of values

The command $\ensuremath{\mathsf{Check}}$ is used to verify that an expression is well-formed

- It returns the type of this expression
- ▶ The type says in which context the expression can be used

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```
Check 2 + 3.

2 + 3 : nat

Check 2.

2 : nat

Check (2 + 3) + 3.

(2 + 3) + 3 : nat
```

The type of functions

The value add3 is not a natural number

Check add3. add3 : nat -> nat

The value add3 is a function

- It expects a natural number as input
- It outputs a natural number

Check add3 + 3. Error the term "add3" has type "nat -> nat" while it is expected to have type "nat"

Applying functions

Function application is written only by juxtaposition

Parentheses are not mandatory

```
Check add3 2.

add3 2 : nat

Eval compute in add3 2.

= 5 : nat

Check add3 (add3 2).
```

add3 (add3 2) : nat

Eval compute in add3 (add3 2).
= 8 : nat

Functions with several arguments

```
At definition time, just use several variables
```

Definition s3 (x y z : nat) := x + y + z. s3 is defined

Check s3. s3 : nat -> nat -> nat -> nat

Functions with one argument that return functions.

```
Check s3 2.

s3 2 : nat -> nat -> nat

Check s3 2 1.
```

```
s3 2 1 : nat -> nat
```

Functions are values

- The value add3 2 is a natural number,
- The value s3 2 1 is a function, like add3
- Functions can also expect functions as argument

Definition rep2 (f : nat -> nat)(x:nat) := f (f x).
rep2 is defined

```
Check rep2.
rep2 : (nat -> nat) -> nat -> nat
```

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Anonymous functions

Abstraction : remove a sub-expression from another expression, make it a variable

```
Check 2 + 3.

2 + 3 : nat

Check fun (x : nat) => x + 3.

fun x : nat => x + 3 : nat -> nat
```

The new expression is a function, usable like add3 or $s3 \ 2 \ 1$.

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Type verification strategy (function application)

Function application is well-formed if types match :

- Assume a function f has type A -> B
- Assume a value a has type A
- then the expression f a is well-formed and has type B

```
Check rep2.
rep2 : (nat -> nat) -> nat -> nat
```

```
Check add3.
add3 : nat -> nat
```

```
Check rep2 add3.
rep2 add3 : nat -> nat
```

Type verification strategy (abstraction)

An anonymous function is well-formed if the body is well formed

- add the assumption that the variable has the input type
- add the argument type in the result
- Example, verify : fun x : nat => x + 3
- x + 3 is well-formed when x has type nat, and has type nat

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Result : fun x : nat => x + 3 has type nat -> nat

Putting data together

- Grouping several pieces of data : tuples,
- fetching individual components : pattern-matching,

```
Check (3,4).

(3, 4) : nat * nat

Check

fun v : nat * nat =>

match v with (x, y) => x + y end.

fun v : nat * nat => let (x, y) := v in x + y

: nat * nat -> nat
```

Numbers

As in programming languages, several types to represent numbers

- natural numbers (non-negative), relative integers, more efficient reprentations
- Need to load the corresponding libraries
- Same notations for several types of numbers : need to choose a scope

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- By default : natural numbers
 - Good properties to learn about proofs
 - Not adapted for efficient computation

Focus on natural numbers

```
Require Import Arith.
Open Scope nat_scope.
```

```
Check (3, 0, S, S (S 0)).
(3, 0, S, 2) : nat * nat * (nat -> nat) * nat
```

```
Fixpoint fact (x : nat) :=
  match x with 0 => 1 | S p => x * fact p end.
fact is recursively defined (...)
```

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Eval compute in fact 6. = 720 : nat

Recursive programming with natural numbers

- Recursive definition keyword : Fixpoint
- Choice of a principal argument
- Recursive calls allowed only on numbers that are smaller than the initial argument
- Smaller numbers found by pattern-matching

Fixpoint fact (x:nat) : nat :=
 match x with 0 => 1 | S p => x * fact p end

This simple schema only for natural numbers! More complex schemas for other datatypes.

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Focus on integers

```
Require Import ZArith.
Open Scope Z_scope.
```

```
Check (3, Z0, xH, xI, x0, Zpos (x0 xH)).

(3, 0, 1%positive, xI, x0, 2)

: Z * Z * positive * (positive -> positive) *

(positive -> positive) * Z
```

```
Eval compute in
    iter 6
        (Z*Z)
        (fun p => let (x,r) := p in (x+1, (x+1)*r))
        (0, 1).
= (6, 720) : Z * Z
```

Mathematical notations

- Depending on top opened scope
- Numbers read as natural numbers or integers
- Force the scope with %nat or %Z
- Infix notation for usual binary operations
- Understand notations with the command Locate

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Understanding notations

```
Locate "_ * _".
Notation Scope
"x * y" := prod x y : type_scope
"x * y" := Pmult x y : positive_scope
"n * m" := mult n m : nat_scope
"x * y" := Zmult x y : Z_scope (default interpretation)
"x * y" := Nmult x y : N_scope
```

Large collections of data

- If A is a datatype then list A is also a datatype
- Theoretically no size limit,
- Written a₁::a₂::...:nil
- ▶ a₁::/ adds the element a₁ in front of I
- nil is an empty list
- Need to load a package to use this, provides many functions

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Examples using lists

```
Require Import List.
```

```
Check 1::2::3::nil.
1::2::3::nil : list Z
Eval compute in (1::2::3::nil) ++ (4::5::nil).
= 1::2::3::4::5::nil : list Z
Eval compute in length (1::2::nil).
= 2 : nat
Eval compute in map (fun x \Rightarrow x + 1) (1::2::3::nil).
= 2::3::4::nil : list Z
Eval compute in
   fold_right (fun x y => x + y) 0 (1::2::3::nil).
= 6 : Z
```

Boolean values

- Values true and false
- Usable in if .. then .. else .. statements

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- comparison function provided for numbers
- To find them : use the command Search