Université Bordeaux 1. Master Sciences & Technologies, Informatique. Examen "Games and synthesis". Responsables H. Gimbert, A. Muscholl Session 1. January 7, 2015, 14h-17h.

Documents: lecture notes and exercice class material.

Exercise 1 A non-deterministic Büchi ω -automaton is called *weak* if every strongly connected component of the automaton contains only accepting states or only non-accepting states (a strongly connected component is a maximal, non-empty set C of vertices such that for every two vertices u, v in C, there is a path from u to v). The figure below represents such an automaton \mathcal{A} with $F = \{s_0, s_3\}$:



- 1. Describe the language accepted by \mathcal{A} by a ω -regular expression.
- 2. Give a *deterministic* automaton that accepts $L(\mathcal{A})$ (you can choose the acceptance condition).
- 3. Show how to complement a weak Büchi automaton into a Büchi automaton (*Hint*: use the subset construction.)

Exercise 2 A Büchi automaton accepts an infinite, Σ -labelled tree t, if there is a run on t such that on every path from the root some final state is seen infinitely often.

Propose a polynomial time algorithm for checking emptiness of Büchi tree automata. Justify your answer.

Exercise 3 A request-response game is a 2-player game played on a graph arena $(V = V_0 \cup V_1, E)$, where the winning condition consists of pairs $(Q_i, R_i)_{i=1}^k$ of requests $Q_i \subseteq V$ and responses $R_i \subseteq V$. A play $\pi = v_0, v_1, \ldots$ is winning if for all $1 \leq i \leq k$, whenever a state in Q_i is visited, a state in R_i is visited now or later:

$$\bigwedge_{i=1}^{k} \mathbf{G}(Q_i \to \mathbf{F}R_i)$$

1. Consider the arena

 $V_0 = \{1, 4, 5\}, V_1 = \{2, 3\}, E = \{(1, 2), (1, 3), (2, 4), (2, 5), (3, 4), (3, 5), (4, 2), (5, 3)\}$

The winning condition is $(Q_1, R_1), (Q_2, R_2)$ with $Q_1 = \{1\}, R_1 = \{4\}, Q_2 = \{2\}, R_2 = \{4, 5\}.$

Determine the winning regions of the two players.

- 2. We will now reduce a request-response game to a Büchi game.
 - (a) Let $A, B \subseteq \Sigma$ be two alphabets, not necessarily disjoint. Give a *deterministic* Büchi automaton accepting the language $\mathbf{G}(A \to \mathbf{F}B)$.
 - (b) Let $A_i, B_i \subseteq \Sigma$ be subalphabets, not necessarily disjoint. Give a *deterministic* Büchi automaton accepting the language $\bigwedge_{i=1}^{k} \mathbf{G}(A_i \to \mathbf{F}B_i)$.
 - (c) Describe the complete reduction of the a request-response game to a Büchi game.