# Combinatorics of non-ambiguous trees 

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## The binary trees

A binary tree is a tree where vertices have 0,1 or 2 children.

size $=n b$ of vertices $=7$.

## All the binary trees of size 4



 $s$

## The number of binary trees: the Catalan numbers

There are

$$
C_{n}=\frac{1}{n+1}\binom{2 n}{n} \quad \underset{+\infty}{\sim} \quad \frac{1}{\pi n \sqrt{n}} 4^{n}
$$

binary trees of size $n$.

$$
\left(C_{n}\right)_{n \geq 1}=(1,2,5,14,42, \ldots)
$$

## Intuitive definition of non-ambiguous trees

A non-ambiguous tree is just a way of drawing a tree in a grid without "ambiguity".

non ambiguous


We forbid the pattern


## The non-ambiguous trees

A non-ambiguous tree of size $n$ is a binary tree of size $n$ embedded in a grid in such a way that:

- the vertices are on the grid
- the edges follow the lines of the grid
- the pattern

生 is forbidden

- each line contains at least one vertex



## The edges of a non-ambiguous tree are not necessary



$$
\text { Size }=\mathrm{nb} \text { of vertices }=6
$$

Forbidden pattern : a vertex can't have two parents.

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## All the non-ambiguous trees of size 4



## The number of non-ambiguous trees

$$
\text { A136127 }=\left(a_{k}\right)_{k \geq 1}=(1,2,5,16,63,294,1585,9692, \ldots)
$$

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$$

The number $a_{n}$ of non-ambiguous trees of size $n$ is equal to the number of permutations having all its excedences placed at the begining. $[$ Steingrímsson, Williams, 07] $+[$ Burstein, 07] $+[$ Aval, B., Nadeau, 11]

$$
a_{n}=\sum_{k+l=n-1}^{n}\left(\sum_{i=1}^{k}(-1)^{k-i} S(k, i) i!i^{l-1}\right)
$$

where $S(k, i)$ are the Stirling numbers of the second kind.
[Ehrenborg and Steingrímsson, 2000]
[Combinatorial proof with non ambiguous trees and tree-like tableaux]

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$$
a_{w, h}=\sum_{p \geq 0}(p+1)!p!S(w, p+1) S(h, p+1)
$$

where $S(k, i)$ are the Stirling numbers of the second kind., $a_{w, h}$ is the number of nat with width $w$ and height $h$.
Patxi Laborde-Zubieta - October/November 2015 - Trois-Rivières ?

## The context

Permutation tableaux and alternative tableaux are objects used to study :

- the PASEP model in physics;
- 2-31 pattern inside a permutation;
- excedences and cycles of a permutation;
- Laguerre polynomials;
(Burstein, Corteel, Dasse-Harteau, Hitczenko, Josuat-Verges, Nadeau, Postnikov, Steingrímsson, Viennot, Williams, Kim, 2005-2011)


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Tree-like tableaux have been introduced to simplify and to explain lots of previous results. (Aval, B., Nadeau, 2011)

The internal structure of tree-like tableaux is described by the non-ambiguous trees.


## Outline

(1) Enumeration of non-ambiguous trees with a fixed underlying tree
(2) The complete non-ambiguous trees and the Bessel function
(3) Bijection between prallelogram polyominoes and binary trees

The non-ambiguous trees associated with a fixed tree


## Outline

(1) Enumeration of non-ambiguous trees with a fixed underlying tree

- Coding a non-ambiguous tree with the coordinates of its vertices
- All these coordinates are described by linear extensions of particular posets
- The non-ambiguous trees are counted by a new hook formula
(2) The complete non-ambiguous trees and the Bessel function
(3) Bijection between prallelogram polyominoes and binary trees


## Some conventions



The $x$ and $y$ coordinates


We choose some arbitrary indices for the vertices

A column is a $x$-oriented line
A row is a $y$-oriented line

The non-ambiguous trees associated with a fixed tree


$N A\left(\infty_{0}^{\infty}\right)=2 \times 3$

## Study of the forbidden pattern


forbidden pattern

- The ending points of two left edges can be in the same row;
- The $x$-coordinates of the ending points of two different edges are different;
- The $x$-coordinates of the ending points of the left edges form an interval $[1, k]$.

These properties are the same for the right edges.

## Total order on the set of the left (resp. right) edges



$$
\begin{gathered}
X(3)<X(4) \\
Y(5)<Y(2)<Y(6)
\end{gathered}
$$

We get two total orders on the ending points of the left (resp. right) edges

$$
\alpha_{L}=34 \quad \alpha_{R}=526
$$

The non-ambiguous trees are represented by $\left(\alpha_{L}, \alpha_{R}\right)$


$$
\begin{gathered}
T=\operatorname{tree}(A) \\
X(0)=Y(0)=0
\end{gathered}
$$

$\forall i$ an ending point of left edges
$\forall j$ an ending point of right edges

$$
\left\{\begin{array}{l}
\alpha_{L}^{-1}(i)=Y(i) \text { et } X(i)=X(\operatorname{parent}(i)) \\
\alpha_{R}^{-1}(j)=X(j) \text { et } Y(j)=Y(\operatorname{parent}(j))
\end{array}\right.
$$

## Question:



What are the constraints verified by $\alpha_{L}$ and $\alpha_{R}$ ?

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Answer: there are linear extensions of two particular posets:

- the poset of the ending points of the left edges;
- the poset of the ending points of the right edges.


## Poset and Hasse diagram

Poset $=$ partialy ordered set


Hasse diagram :


$$
\begin{aligned}
& 1 \leq 3 \leq 4 \\
& 1 \leq 2 \leq 4
\end{aligned}
$$

## Linear extensions of a poset

It is a totaly ordered set with respect with the oder of $P$


Linear extensions of the previous poset : 1234, 1243, 2134, 2143 Notation: $\mathcal{L}(P)$ is the set of all the linear extensions of $P$.

The constraints verifyed by the ending points of the left edges $\left(\alpha_{L}\right)$

$\alpha_{h}$ is a linear extension of the poset $V_{L}$ defined on the ending points of the left edges with the following relation order:

$$
i=\operatorname{ancestor}(j) \quad \Longleftrightarrow \quad i<v_{L} j
$$

The Hasse diagrams of the poset $V_{L}$ and $V_{R}$

$\alpha_{L}\left(\right.$ resp. $\left.\alpha_{R}\right)$ are a linear extension of $V_{L}\left(\right.$ resp. $\left.V_{R}\right)$.

## Non-ambiguous trees with a fixed underlying tree



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$$
V_{L}={ }_{\bullet}^{2} \stackrel{6}{\bullet}^{6} \quad V_{R}=\int_{4}^{3}{ }_{5}^{5}
$$

## Non-ambiguous trees with a fixed underlying tree



$$
v_{L}=2_{0}^{6} \quad V_{R}=\int_{4}^{3} \stackrel{5}{5}
$$

$$
\begin{gathered}
\alpha_{L} \times \alpha_{R} \\
\in \\
\mathcal{L}\left(V_{L}\right) \times \mathcal{L}\left(V_{R}\right)
\end{gathered}
$$

Non-ambiguous trees with a fixed underlying tree

$$
\begin{aligned}
& T= \\
& V_{L}={ }_{\bullet}^{2}{ }_{6}^{6} \quad V_{R}=\int_{4}^{3} \stackrel{5}{\bullet} \\
& (26,534) \quad(26,354) \quad(26,345) \\
& \alpha_{L} \times \alpha_{R} \\
& \underset{\mathcal{L}\left(V_{L}\right) \times \mathcal{L}\left(V_{R}\right)}{\in} \Longrightarrow \quad(62,534) \quad(62,354) \quad(62,345)
\end{aligned}
$$

Non-ambiguous trees with a fixed underlying tree




$$
\mathcal{L}\left(V_{L}\right) \times \mathcal{L}\left(V_{R}\right) \Longrightarrow \begin{array}{ccc}
(62,534) & \Longrightarrow & (62,354)
\end{array}
$$

The number of non-ambiguous trees with a fixed underlying tree

$$
\begin{aligned}
& =\operatorname{hook}\left(\begin{array}{ll}
2 & 6 \\
\bullet & 6
\end{array}\right) \times \operatorname{hook}\left(\begin{array}{ll}
3 & 5 \\
& \bullet \\
4 &
\end{array}\right) \\
& =\frac{2!}{1.1} \times \frac{3!}{1.1 .2} \\
& =2 \times 3
\end{aligned}
$$

## A hook formula for the non-ambiguous trees



We associate each edge with an integer $n_{e}$. Given a left (resp.right) edge $e$, the integer $n_{e}$ is the number of left (resp. right) edges contained in the subtree whose root is the ending point of $e$, plus 1.

$$
\begin{aligned}
\mathcal{N} \mathcal{A}(T) & =\frac{\sharp\{\text { left edges }\}!\cdot \sharp\{\text { right edges }\}!}{\prod\{\text { left labels }\} \cdot \prod\{\text { right labels }\}} \\
& =\frac{5!\cdot 4!}{1 \cdot 1 \cdot 1 \cdot 1 \cdot 3 \cdot 1 \cdot 1 \cdot 2 \cdot 2} \\
& =\quad 240
\end{aligned}
$$

## Outline

(1) Enumeration of non-ambiguous trees with a fixed underlying tree
(2) The complete non-ambiguous trees and the Bessel function

- The complete non-ambiguous trees are related with the formal power series of the logarithm of the Bessel function of order 0
- The root suppression gives a recursive enumeration
- Another identities
(3) Bijection between prallelogram polyominoes and binary trees


## The compete non-ambiguous trees

A complete non-ambiguous tree is a non-ambiguous tree whose vertices have only 0 or 2 children.

size $=\mathrm{nb}$ of internal vertices
$=\mathrm{nb}$ of left edges
$=\mathrm{nb}$ of right edges
$=6$

## The number of non-ambiguus trees

The number $b_{n}$ of complete non-ambiguous trees of size is referenced on OEIS:

$$
\mathrm{A} 136127=\left(b_{n}\right)_{n \geq 0}=(1,1,4,33,456,9460, \ldots)
$$

## The number of non-ambiguus trees

The number $b_{n}$ of complete non-ambiguous trees of size is referenced on OEIS:

$$
\begin{aligned}
\mathrm{A} 136127=\left(b_{n}\right)_{n \geq 0} & =(1,1,4,33,456,9460, \ldots) \\
-\ln \left(J_{0}(x)\right) & =\sum_{k \geq 0} b_{k} \frac{x^{2(k+1)}}{\left((k+1)!2^{k+1}\right)^{2}}
\end{aligned}
$$

where $J_{O}$ is the Bessel function of order 0 .
[Carlitz, 1963]

## Proof

The Bessel function of order 0 is the solution of the differential equation:

$$
\frac{d^{2} y}{d x^{2}}+\frac{1}{x} \frac{d^{2}}{d x}+y=0
$$

such that the first coefficients in its series expansion are $j_{0}=1$ and $j_{1}=0$.

We try

$$
B(x)=\exp \left(-\sum_{k \geq 0} b_{k} \frac{x^{2(k+1)}}{\left((k+1)!2^{k+1}\right)^{2}}\right)
$$

## Proof

The function

$$
B(x)=\exp \left(-\sum_{k \geq 0} b_{k} \frac{x^{2(k+1)}}{\left((k+1)!2^{k+1}\right)^{2}}\right)
$$

satisfies the Bessel differential equation if and only if

$$
b_{n+1}=\sum_{u+v=n}\binom{n+1}{u}\binom{n+1}{v} b_{u} b_{v}
$$

The root suppression gives a recursive enumeration


The coefficients $b_{k}$ are the number of complete non-ambiguous trees of size $k$.

## Another identity involving complete non-ambiguous trees

$$
\sum_{k=0}^{n-1}(-1)^{k}\binom{n}{k+1}\binom{n-1}{k} b_{k}=1
$$

[Carlitz, 1963]

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Combinatorial proof:

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[Carlitz, 1963]
Combinatorial proof:


A gridded tree

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Combinatorial proof:



## Outline

(1) Enumeration of non-ambiguous trees with a fixed underlying tree
(2) The complete non-ambiguous trees and the Bessel function
(3) Bijection between prallelogram polyominoes and binary trees

- Definition of a prallelogram polyomino
- Parallelogram polyominoes are the shapes of particular non-ambiguous trees


## The parallelogram polyominoes

A parallelogram polyomino of size $n$ is a pair of lattice paths of lengths $n+1$ with south-west and south-east steps starting at the same point, ending at the same point and never meeting each other.


Parallelogram polyominoes and binary trees are in bijection.
[Delest and Viennot, 1984]

Bijection between prallelogram polyominoes and binary trees


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## In the future ...

We can try to

- find a combinatorial interpretation of the Bessel functions for any order;
- study the link between tree-like tableau background and the non-ambiguous trees.
- define some Hopf Algerbras for non-ambiguous trees.
- generalize in other dimensions (in progress ...) (Patxi Laborde-Zubieta in October 2015 - Trois-Rivières ?)


## Thank you for your attention!

