#### Combinatorics of non-ambiguous trees

#### Adrien Boussicault

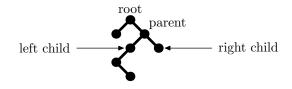
#### Joint work with J.C. Aval, M. Bouvel, M. Silimbani

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Université du Quebec à Trois-Rivières - DMI - 6/5/2015

## The binary trees

A binary tree is a tree where vertices have 0, 1 or 2 children.



size = nb of vertices = 7.

All the binary trees of size 4

# $\frac{\langle \langle \langle \langle \langle \langle \langle \rangle \rangle \rangle \rangle}{\langle \langle \rangle \rangle \langle \langle \langle \rangle \rangle \rangle}$

Adrien Boussicault (LaBRI)

Combinatorics of non-ambiguous trees

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#### The number of binary trees: the Catalan numbers

There are

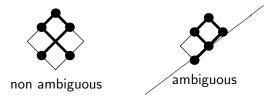
$$C_n = rac{1}{n+1} {2n \choose n} \qquad \sim \qquad rac{1}{\pi n \sqrt{n}} 4^n$$

binary trees of size n.

$$(C_n)_{n\geq 1} = (1, 2, 5, 14, 42, \dots)$$

## Intuitive definition of non-ambiguous trees

A non-ambiguous tree is just a way of drawing a tree in a grid without "ambiguity".

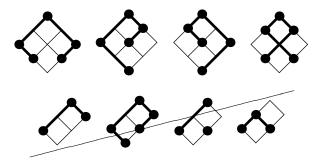


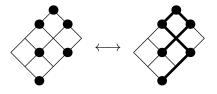
# We forbid the pattern

### The non-ambiguous trees

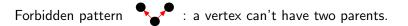
A non-ambiguous tree of size n is a binary tree of size n embedded in a grid in such a way that:

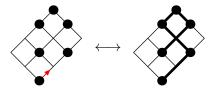
- the vertices are on the grid
- the edges follow the lines of the grid
- $\bullet$  the pattern  $\checkmark$  is forbidden
- each line contains at least one vertex

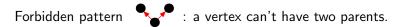


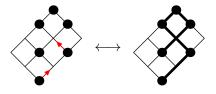


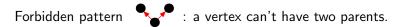
Size = nb of vertices = 
$$6$$

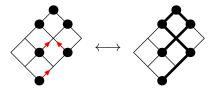


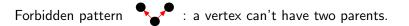


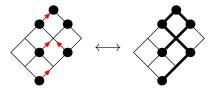


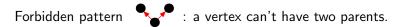


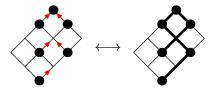


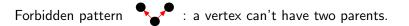




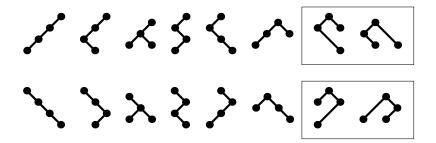








All the non-ambiguous trees of size 4



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The number of non-ambiguous trees

 $A136127 = (a_k)_{k\geq 1} = (1, 2, 5, 16, 63, 294, 1585, 9692, \dots)$ 

#### The number of non-ambiguous trees

$$\texttt{A136127} = (a_k)_{k \ge 1} = (1, 2, 5, 16, 63, 294, 1585, 9692, \dots)$$

The number  $a_n$  of non-ambiguous trees of size n is equal to the number of permutations having all its excedences placed at the begining. [Steingrímsson, Williams, 07] + [Burstein, 07] + [Aval, B., Nadeau, 11]

$$a_n = \sum_{k+l=n-1}^n \left( \sum_{i=1}^k (-1)^{k-i} S(k,i) i! i^{l-1} \right)$$

where S(k, i) are the Stirling numbers of the second kind. [Ehrenborg and Steingrímsson, 2000] [Combinatorial proof with non ambiguous trees and tree-like tableaux]

## The number of non-ambiguous trees

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$$a_{w,h} = \sum_{p \ge 0} (p+1)! p! S(w, p+1) S(h, p+1)$$

where S(k, i) are the Stirling numbers of the second kind.,  $a_{w,h}$  is the number of nat with width w and height h.

Patxi Laborde-Zubieta - October/November 2015 - Trois-Rivières ?

#### The context

Permutation tableaux and alternative tableaux are objects used to study :

- the PASEP model in physics;
- 2-31 pattern inside a permutation;
- excedences and cycles of a permutation;
- Laguerre polynomials;

• ...

(Burstein, Corteel, Dasse-Harteau, Hitczenko, Josuat-Verges, Nadeau, Postnikov, Steingrímsson, Viennot, Williams, Kim, 2005-2011)

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Tree-like tableaux have been introduced to simplify and to explain lots of previous results. (Aval, B., Nadeau, 2011)

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Tree-like tableaux have been introduced to simplify and to explain lots of previous results. (Aval, B., Nadeau, 2011)

The internal structure of tree-like tableaux is described by the non-ambiguous trees.



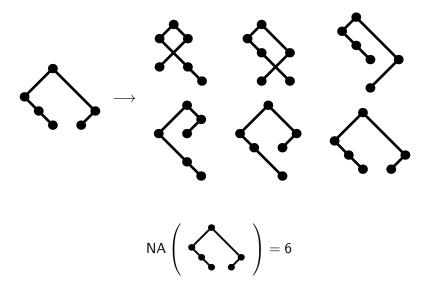
tree-like tableau

<sup>• ...</sup> 

## Outline

- Enumeration of non-ambiguous trees with a fixed underlying tree
- 2 The complete non-ambiguous trees and the Bessel function
- 3 Bijection between prallelogram polyominoes and binary trees

The non-ambiguous trees associated with a fixed tree



## Outline

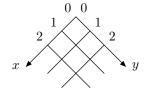
Enumeration of non-ambiguous trees with a fixed underlying tree

- Coding a non-ambiguous tree with the coordinates of its vertices
- All these coordinates are described by linear extensions of particular posets
- The non-ambiguous trees are counted by a new hook formula

#### 2) The complete non-ambiguous trees and the Bessel function

3 Bijection between prallelogram polyominoes and binary trees

#### Some conventions



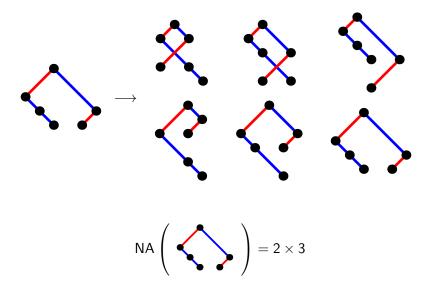
The x and y coordinates



We choose some arbitrary indices for the vertices

#### A column is a *x*-oriented line A row is a *y*-oriented line

#### The non-ambiguous trees associated with a fixed tree



# Study of the forbidden pattern

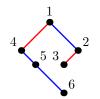


forbidden pattern

- The ending points of two left edges can be in the same row;
- The x-coordinates of the ending points of two different edges are different;
- The x-coordinates of the ending points of the left edges form an interval [1, k].

These properties are the same for the right edges.

Total order on the set of the left (resp. right) edges



X(3) < X(4)Y(5) < Y(2) < Y(6)

We get two total orders on the ending points of the left (resp. right) edges

$$\alpha_L = 34$$
  $\alpha_R = 526$ 

The non-ambiguous trees are represented by  $(\alpha_L, \alpha_R)$ 

$$A = \overset{1}{4} \underbrace{5}_{6} \overset{2}{3} \underbrace{0}_{6} \overset{2}{\leftarrow} \left( T = \overset{1}{4} \underbrace{3}_{5} \overset{2}{6} , \alpha_{L} = 34, \alpha_{R} = 526 \right)$$

$$T = tree(A)$$

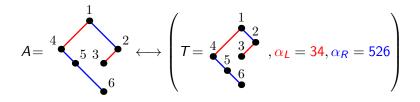
$$X(0) = Y(0) = 0$$

$$\forall i \text{ an ending point of left edges}$$

$$\forall j \text{ an ending point of right edges}$$

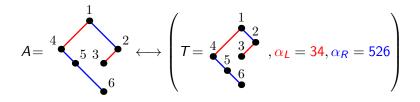
$$\begin{cases} \alpha_L^{-1}(i) = Y(i) \text{ et } X(i) = X(parent(i)) \\ \alpha_R^{-1}(j) = X(j) \text{ et } Y(j) = Y(parent(j)) \end{cases}$$

Question:



What are the constraints verified by  $\alpha_L$  and  $\alpha_R$  ?

Question:



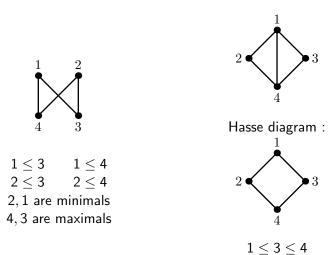
What are the constraints verified by  $\alpha_L$  and  $\alpha_R$ ?

Answer: there are linear extensions of two particular posets:

- the poset of the ending points of the left edges;
- the poset of the ending points of the right edges.

#### Poset and Hasse diagram

Poset = partialy ordered set



1 < 2 < 4

Linear extensions of a poset

It is a totaly ordered set with respect with the oder of P



Linear extensions of the previous poset : 1234, 1243 , 2134, 2143 Notation:  $\mathcal{L}(P)$  is the set of all the linear extensions of P. The constraints verifyed by the ending points of the left edges (  $\alpha_{\rm L}$  )

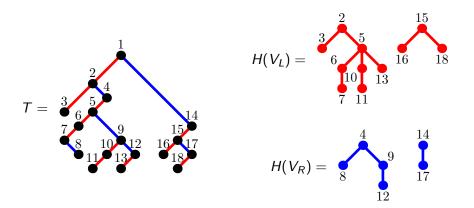


 $i = ancestor(j) \implies X(i) < X(j) \implies i <_{\alpha_L} j$ 

 $\alpha_h$  is a linear extension of the poset  $V_L$  defined on the ending points of the left edges with the following relation order:

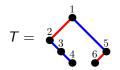
$$i = ancestor(j) \iff i <_{V_L} j$$

The Hasse diagrams of the poset  $V_L$  and  $V_R$ 

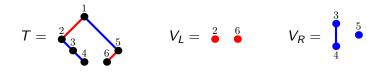


 $\alpha_L$  (resp.  $\alpha_R$ ) are a linear extension of  $V_L$  (resp.  $V_R$ ).

Non-ambiguous trees with a fixed underlying tree



Non-ambiguous trees with a fixed underlying tree



Non-ambiguous trees with a fixed underlying tree

$$T = \begin{array}{c} 2 \\ 3 \\ 4 \\ 4 \\ 6 \end{array} \begin{array}{c} 5 \\ 6 \\ 6 \end{array} \qquad V_L = \begin{array}{c} 2 \\ \bullet \end{array} \begin{array}{c} 6 \\ \bullet \end{array} \qquad V_R = \begin{array}{c} 3 \\ 4 \\ 4 \\ 4 \end{array} \begin{array}{c} 5 \\ \bullet \end{array}$$

$$\begin{array}{c} \alpha_L \times \alpha_R \\ \in \\ \mathcal{L}(V_L) \times \mathcal{L}(V_R) \end{array}$$

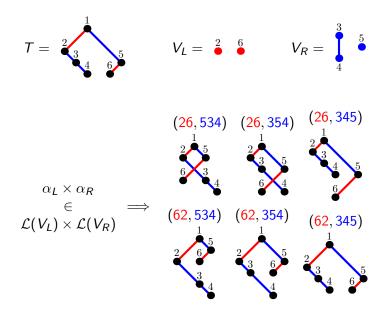
Non-ambiguous trees with a fixed underlying tree

$$T = \begin{array}{c} 2 \\ 3 \\ 4 \\ 4 \end{array} \begin{array}{c} 6 \\ 5 \\ 6 \end{array} \qquad V_L = \begin{array}{c} 2 \\ 6 \\ 6 \end{array} \qquad V_R = \begin{array}{c} 3 \\ 4 \\ 4 \end{array} \begin{array}{c} 5 \\ 6 \\ 4 \end{array} \begin{array}{c} 5 \\ 6 \\ 6 \end{array}$$

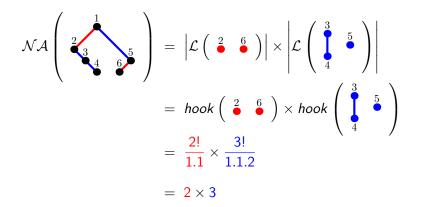
(26, 534) (26, 354) (26, 345)

$$\begin{array}{c} \alpha_L \times \alpha_R \\ \in \\ \mathcal{L}(V_L) \times \mathcal{L}(V_R) \end{array} \implies (62, 534) \quad (62, 354) \quad (62, 345) \end{array}$$

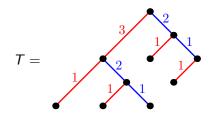
Non-ambiguous trees with a fixed underlying tree



The number of non-ambiguous trees with a fixed underlying tree



## A hook formula for the non-ambiguous trees



We associate each edge with an integer  $n_e$ . Given a left (resp.right) edge e, the integer  $n_e$  is the number of left (resp. right) edges contained in the subtree whose root is the ending point of e, plus 1.

$$\mathcal{NA}(T) = \frac{\{\{\text{left edges}\} : \{\{\text{right edges}\}\}}{\prod\{\{\text{left labels}\} \cdot \prod\{\text{right labels}\}\}}$$
$$= \frac{5! \cdot 4!}{1 \cdot 1 \cdot 1 \cdot 3 \cdot 1 \cdot 1 \cdot 2 \cdot 2}$$
$$= 240$$

# Outline

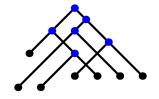
### Enumeration of non-ambiguous trees with a fixed underlying tree

#### 2 The complete non-ambiguous trees and the Bessel function

- The complete non-ambiguous trees are related with the formal power series of the logarithm of the Bessel function of order 0
- The root suppression gives a recursive enumeration
- Another identities

# The compete non-ambiguous trees

A complete non-ambiguous tree is a non-ambiguous tree whose vertices have only 0 or 2 children.



- size = nb of internal vertices
  - = nb of left edges
  - = nb of right edges
  - = 6

# The number of non-ambiguus trees

The number  $b_n$  of complete non-ambiguous trees of size is referenced on OEIS:

$$\texttt{A136127} = (b_n)_{n \ge 0} = (1, 1, 4, 33, 456, 9460, \dots)$$

# The number of non-ambiguus trees

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$$\texttt{A136127} = (b_n)_{n \ge 0} = (1, 1, 4, 33, 456, 9460, \dots)$$

$$-\ln (J_0(x)) = \sum_{k\geq 0} b_k \frac{x^{2(k+1)}}{((k+1)!2^{k+1})^2}$$

where  $J_O$  is the Bessel function of order 0. [Carlitz, 1963]

## Proof

The Bessel function of order 0 is the solution of the differential equation:

$$\frac{d^2y}{dx^2} + \frac{1}{x}\frac{d^2}{dx} + y = 0,$$

such that the first coefficients in its series expansion are  $j_0 = 1$  and  $j_1 = 0$ .

#### We try

$$B(x) = exp\left(-\sum_{k\geq 0} b_k \frac{x^{2(k+1)}}{((k+1)!2^{k+1})^2}\right)$$

# Proof

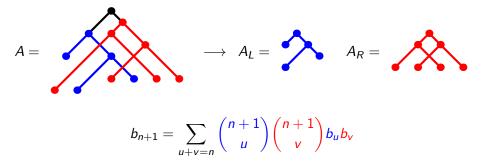
The function

$$B(x) = exp\left(-\sum_{k\geq 0}b_krac{x^{2(k+1)}}{((k+1)!2^{k+1})^2}
ight)$$

satisfies the Bessel differential equation if and only if

$$b_{n+1} = \sum_{u+v=n} \binom{n+1}{u} \binom{n+1}{v} b_u b_v.$$

The root suppression gives a recursive enumeration



The coefficients  $b_k$  are the number of complete non-ambiguous trees of size k.

$$\sum_{k=0}^{n-1} (-1)^k \binom{n}{k+1} \binom{n-1}{k} b_k = 1$$

[Carlitz, 1963]

$$\sum_{k=0}^{n-1} (-1)^k \binom{n}{k+1} \binom{n-1}{k} b_k = 1$$

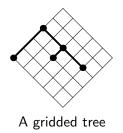
[Carlitz, 1963]

Combinatorial proof:

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[Carlitz, 1963]

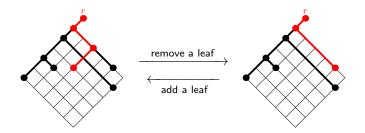
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Combinatorial proof:



# Outline

D Enumeration of non-ambiguous trees with a fixed underlying tree

2 The complete non-ambiguous trees and the Bessel function

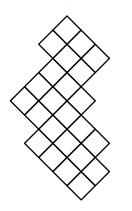
- Definition of a prallelogram polyomino
- Parallelogram polyominoes are the shapes of particular non-ambiguous trees

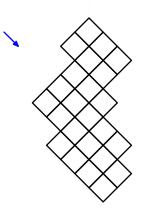
# The parallelogram polyominoes

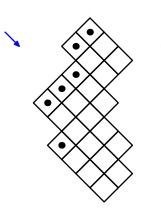
A parallelogram polyomino of size n is a pair of lattice paths of lengths n+1 with south-west and south-east steps starting at the same point, ending at the same point and never meeting each other.

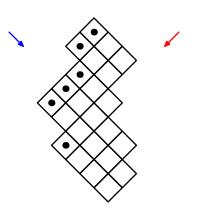


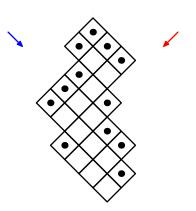
Parallelogram polyominoes and binary trees are in bijection. [Delest and Viennot, 1984]

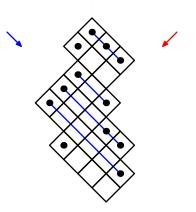


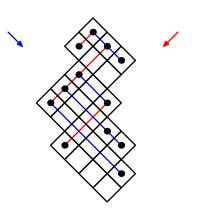


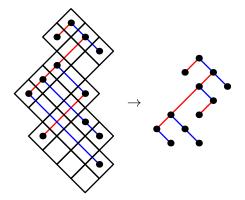












# In the future ...

We can try to

- find a combinatorial interpretation of the Bessel functions for any order;
- study the link between tree-like tableau background and the non-ambiguous trees.
- define some Hopf Algerbras for non-ambiguous trees.
- generalize in other dimensions (in progress ...) (Patxi Laborde-Zubieta in October 2015 - Trois-Rivières ?)

# Thank you for your attention!