

Programming Metasystems with Active Objects

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Introduction and motivations

*Models for programming metasystems

*A brief introducton to HiMM

- SThe Hierarchical Metacomputer architecture (HiM)
- The Customizable architecture of a Node
- ₲ The API of HiMM
- *The integration of ProActive with HiMM

Serformance analisys

Conclusions and future work



Introduction and motivations

- * The presence of a huge amount of computers and mobile devices results in an increased focus on the interconnection of systems
- * This evolution of the computer scenario has promoted two new trends in distributed computing:
 - **Grid computing**, for scientific applications
 - **Web Services**, for e-commerce and business applications
- * In both domains, it is important to
 - define suitable programming models in order to better exploit large-scale distributed systems
 - improve customizability of middleware platforms



Models for programming metasystem

- Most of the middleware platforms for Grid computing does not define a model for programming parallel applications
- The send/receive model is used as the de-facto standard for communication both in cluster environments and in large heterogeneous distributed systems
- However, to handle the unpredictability of resource avaiibility and behavior more dynamic programming models are required

S Agents

 completely change the way distributed applications are designed and deployed

Active Objects

* allow for taking advantage of well-defined theories and design techniques



Our proposal (1)

Even though many active objects libraries have been developed, the diffusion of Grids has introduced new issues to be taken into account during program development

heterogeneity, scalability, unpredictability and adaptability

- * We think that the separation of aspects gives the proper flexibility for programming Grid applications
 - computational entities should exploit communication, security and management features of a middleware platform without affecting the functional aspects of an application
- * So, we have chosen an approach based on
 - A generalized, customizable middleware platform
 - A customization to support an active object programming model

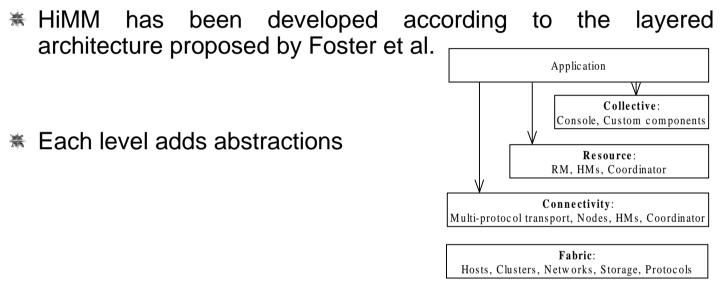


Our proposal (2)

- * The middleware adopted is HiMM
 - *G* Hierarchical Metacomputer Middleware
- The programming model adopted is the one provided by ProActive
 - Currently ProActive is implemented on top of RMI that does not offer specialized services for Grid computing
- ★ This way
 - an application can be programmed by exploiting the asynchronous remote method invocation model
 - set metasystems can be transparently managed by using meta-objects
 - through meta-objects, programmers can exploit the services of the underlying middleware without affecting the functional code of the application.



A brief introduction to HiMM (1) layered architecture

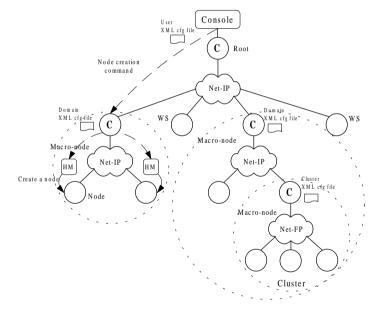


* At fabric layer, it allows a user to exploit collections of hosts, which can be workstations or computing units of parallel systems, interconnected by heterogeneous networks



A brief introduction to HiMM (2) connectivity layer

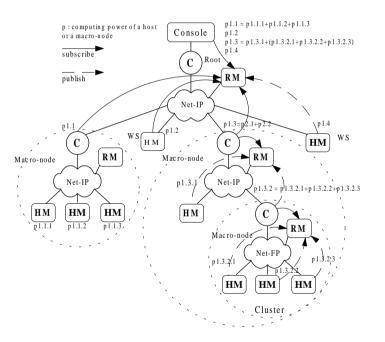
- HiMM manages resources according to a hierarchical topology (Hierarchical Metacomputers) in order to
 - meet the constraints of the Internet organization
 - scheme state state
 - 🧯 improve scalability
- Clusters of computers hidden from the Internet or intra-connected by fast networks are seen as *macro-nodes*
 - A macro-node is a high-level concept that allows clusters to be transparently used as a single powerful machine
 - A macro-node can in turn contain other macronodes
 - A metacomputer can be organized according to a recursive tree topology
- The Coordinator interfaces the macronode with the metacomputer network





A brief introduction to HiMM (3) resource layer

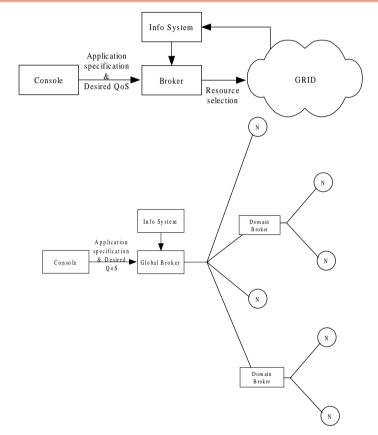
- A macro-node is characterized by two main components:
 - S The Host Manager (HM)
 - The Resource Manager (RM)
- The HM runs on each node wanting to donate CPU cycles
- The RM is use to publish the available computing power at each level
- A macro-node manages another important component:
 - the Distributed Class Storage System (DCSS)
 - This component allows a HiM to run applications even if application code is not present on nodes





A brief introduction to HiMM (4) collective layer

- The broker acts as a mediator between the user (console) and grid resources (HiM) by using middleware services (Info System and HMs)
 - the broker becomes responsible for resource discovery, resource selection, process/task mapping, task scheduling, and presents the Grid (HiM) as a single, unified resource
- The architecture of the broker is distributed and hierarchical in order to mach the architectural organization of a HiM



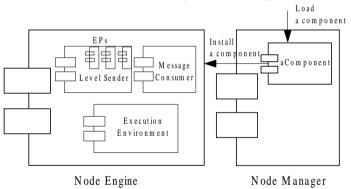


A brief introduction to HiMM (5) node architecture

- HiMM implements its services in several software components whose interfaces are designed according to Component а Framework approach
- Both nodes and coordinators are 촕 processes in which a set of software components are loaded either at start-up or at run-time
- The main components are: 鮝
 - the Node Manager (NM)
 - guarantees macro-node consistency ÷
 - provides users with services for writing distributed applications
 - * takes charge of some system tasks, such as the creation of new nodes at run-time
 - the Node Engine (NE) 6
 - + The Message Consumer (MC)
 - The Execution Environment (EE)
 - The Level Sender (LS)

Node Manager (NM) Node Engine (NE) Level Sender (LS)

: manages a node in order to guarantee the HiM consistency : defines the behavior of a node : sends a message (an object) to a node or a group of nodes Message Consumer (MC) : receives messages from the network Execution Environment (EE) : stores a component of a distributed application

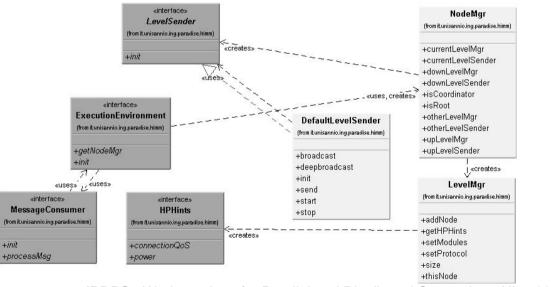


(Level Manager)

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* A component is a Java class implementing one or more interfaces according to the pattern "Inversion of control"





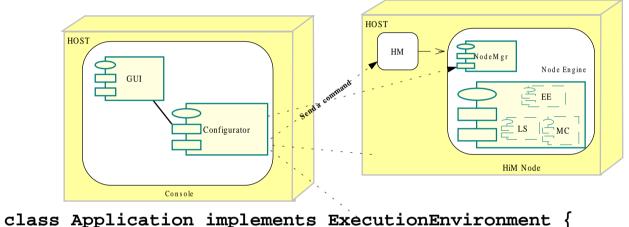
A brief introduction to HiMM (7) communication API

- HiMM allows nodes to communicate using the component LevelSender
- * A LevelSender can be customized even if a default one is always available
 - Control Con

class DefaultLevelSender impements LevelSender {
public void send(Object m, int node) { ... }
public void broadcast(Object m) { ... }
public void deepBroadcast (Object m) { ... }



A brief introduction to HiMM (8) interaction among components

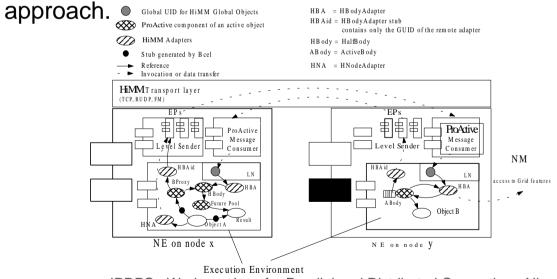


```
class Application implements ExecutionEnvironment {
  void init(NodeMgr nm) throws ... {
      nodeMgr = nm;
  }
  void start() {
      ...
      nodeMgr.downLevelSender().send(msg, node);
      nodeMgr.downLevelMgr().addNode(...);
      Object o = nodeMgr.downLevelMsgConsumer().receive(); ...
  }}
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```



Integration HiMM / ProActive (1)

- The implementation of P/H has been made easy thanks to the particular organization of both software systems
 - Service ProActive is heavily based on the Adapter Pattern
 - G HiMM is implemented following the Component Framework





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Integration HiMM / ProActive (2)

public class MatrixMultiply extends ProActiveExecutionEnvironment { private NodeMgr nodeMgr; private volatile boolean stop = false; public void init (NodeMgr nm) { super.init(nm); nodeMgr = nm; NodeFactory.setFactory("himm", new HNodeFactory()); public void start() { stop = false;if(nodeMgr.isRoot()) HMatrix[] mDxActiveGroup=null; int dim=0; System.out.println("Insert matrix size:"); <read dim>; Matrix mDx = new Matrix(dim); LevelMgr lm = nodeMgr.downLevelMgr(); int size = lm.size()-1; Object[] po = new Object[]{mDx.getTab()}; mDxActiveGroup = new Matrix[size]; for(int i =0; i<size; i++)</pre> mDxActiveGroup[i] = (HMatrix) ProActive.newActive(HMatrix.class.getName(), po, NodeFactory.getNode("himm://"+(i+1)), null, HMetaObjectFactory.newInstance()); Matrix mSx = new Matrix(dim); while (!stop) { Matrix[] results = multiply(mSx, mDxActiveGroup); Matrix result = Matrix.rebuild(results, dim); . . .

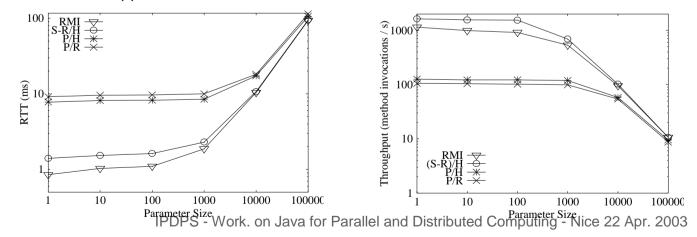
public void runActivity(Body body) { Service service = new Service(body); while (body.isActive()) { Request r = service.blockingRemoveOldest(); HBodyAdapter hba =(HBodyAdapter) body.getRemoteAdapter(); NodeMgr nodeMgr = hba.getNodeMgr(); if(r.getMethodName().equals("multiply")&& nodeMgr.isCoordinator()){ HiMMRequest rr = (HiMMRequest)r; float[][]mSxtab =(float[][])rr.methodCall.getParameter(0); Matrix mSx = new Matrix(mSxtab); Matrix[] results = new Matrix[subMats.length]; Object[] subSxMats = mSx.createSubMatrixes(subMats.length); for(int i=0; i< subMats.length; i++)</pre> results[i] = subMats[i].multiply((float[][])subSxMats[i]); Matrix result = Matrix.rebuild(results, mSxtab.length); Object[] params = new Object[]{result}; Class[] classes = new Class[]{Matrix.class}; try { Method m = HMatrix.class.getMethod("getResult", classes); rr.methodCall = MethodCall.getMethodCall(m, params); } catch (NoSuchMethodException e) { ... } service.serve(rr); } else service.serve(r);



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Performance analisys (1)

- * We have conducted a first analysis to compare RMI with HiMM and P/R with P/H
- * The benchmark is the invocation of a remote method (with one parameter and an empty body) and the return of an integer value, for a varying size of the parameter
- The figure shows that for a small size (1 byte) of the parameter, RMI RTT (0.9 ms) is smaller than HiMM RTT (1.3 ms), but P/H RTT (7.8 ms) is smaller than P/R RTT (9.2 ms)
 - Even if the HiMM transport layer is not optimized, P/H behaves better than P/R due to the use of asynchronous messaging that allows low-level ProActive operations to be overlapped with communication

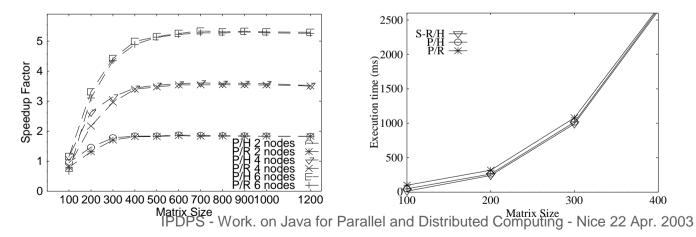




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Performance analisys (2)

- * A further analysis has aimed to measure the speedup factor by running a simple application benchmark
- The benchmark is the product of two square matrixes with different sizes of the matrixes.
- The performance obtained with P/H is slightly better than the one obtained with P/R, especially when the size of the matrixes is small
 - 5 This is mainly due to the improvement of remote method invocation implemented by HiMM
 - When the size of the matrixes is large, the execution time is dominated by the time of the matrix serialization, which is the same in P/H and P/R





Conclusions

- I have described the integration of ProActive with HiMM to create a new framework for easily programming Grid applications with the active object model
- The separation of concerns and the use of meta-objects allow programmers to exploit the underlying features of HiMM to program distributed aspects of applications without affecting functional code
 - Chis approach enables code reuse and makes the object oriented approach effective for the development of distributed and parallel applications
- * HiMM assures better performance than RMI to ProActive
 - S The overhead introduced by the new framework is small



Future work

- * In the future, we will conduct further experiments with a larger, hierarchical distributed system
 - Characterized by different computational power, communication hardware and protocols
- We will complete the definition and the implementation of the broker architecture
- We will define a framework for automatically transforming sequential object oriented applications in parallel and distributed applications when the master/slave parallel model can be adopted