Short vertex labels for connectivity checking in planar graphs with forbidden parts Bruno Courcelle, LaBRI, Bordeaux, and Institut Universitaire de France, courcell@labri.fr

Let P(x, y, Z) be the graph property, called a *connectivity query*, that holds iff Z does not separate x from y, i.e., iff there a path from x to y that avoids Z.

Main result: An algorithm A can define for each vertex u of a connected planar graph G a bit sequence J(u) of length $O(\log(n))$ called a *label*, and another algorithm B can decide if P(x, y, Z) is true by using only J(x), J(y) and J(Z), the sets of labels J(z) for z in Z. This algorithm has no other knowledge of G than these labels.

There are other results of this type that construct *labeling schemes*, i.e., pairs (A,B) as above, for logically specified graph properties and particular graph classes :

(1) For monadic second-order properties and graphs of bounded tree-width or clique-width (Courcelle and Vanicat [CV]),

(2) for first-order properties (with set arguments) and graphs that can be covered by graphs of bounded tree-width or clique-width with limited overlaps (Courcelle, Gavoille and Kanté [CGK]),

(3) for connectivity queries and the associated distance computations, in graphs of bounded tree-width or clique-width (Courcelle and Twigg [CT]); this is a special case of (1) but with a proof that does not use the huge automata constructed from logical formulas.

The connectivity query is monadic second-order expressible but not first-order expressible, and the class of planar graphs has unbounded clique-width. Hence, this result is not a consequence of the above cited ones. Here is the structure of the construction. A labeling scheme for connectivity queries in 3-connected planar graphs (and even in 2-connected planar graphs such that any two vertices belong to at most m faces for some fixed m) is first constructed. It uses a *straightline embedding* of the given graph with positive integer coordinates of the vertices of value O(n). Such values are encoded by bit sequences of length $\log(n)$. The label J(u) of a vertex u consists of its coordinates and of those of a bounded number of neighbour vertices and of points in the plane that are "in the middle" of certain faces. The decoding by B uses algorithms of computational geometry for the *planar point location problem* because the validity of P(x, y, Z) is equivalent to the fact that the points of the plane corresponding to x and y are not separated by certain line segments linking vertices to points representing faces, the coordinates of which are in J(x), J(y)and J(Z).

For general connected planar graphs, this geometric method cannot detect all cases where Z separates x from y. We decompose the given graph into a tree of 3-connected components. On this tree, one can detect *some cases of separation*, for instance the case where Z contains a separating vertex s of the graph G such that every path from x to y goes through s. A labeling K(u) using the result of [CV] can be constructed to detect this case, and also some cases where x and y are separated by a pair of vertices in Z. By combining this labeling and the one based on a straightline embedding of G, one obtains the desired labeling.

This construction can be extended to graphs that are planar gluings of graphs of bounded clique-width with limited overlaps (Kanté [K]).

All results presented in the lecture are available in the following publications.

[CGK] B. Courcelle, C. Gavoille, M.M. Kanté : Efficient First-Order Model-Checking using Short Labels, Proceedings of Frontiers of Algorithms Workshop, June 2008, Changsha, China, Lec. Notes Comput. Sci. 5059 (2008) 159-170

[CGKT08] B. Courcelle, C. Gavoille, M.M. Kanté, A. Twigg : Connectivity check in 3-connected planar graphs with obstacles, Electronic Notes in Discrete Mathematics 31 (2008) 151-155

[CGKT09] B. Courcelle, C. Gavoille, M.M. Kanté, A. Twigg : Optimal labeling for connectivity checking in planar networks with obstacles, Preprint, December 2008.

[CV] B. Courcelle, R. Vanicat : Query efficient implementations of graphs of bounded cliquewidth, Discrete Applied Mathematics 131 (2003) 129-150.

[CT] B. Courcelle, A. Twigg : Compact forbidden-set routing, Proceedings STACS 2007, Lec. Notes. Comput. Sci. 4393 (2007) 37-48.

[K] : M. M. Kanté : Graph structurings : some algorithmic applications, Ph.D., December 2008, Bordeaux 1 University (written in English).