

Profile-Guided Java Program Partitioning for Power Aware Computing

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Mobile Devices are very powerful.

- The iPAQ handheld
 - ✤ 400 MHz. Intel Xscale Processor
 - ✤ 128 MB RAM
 - Wireless communication capabilities

Battery Power is a bottleneck.

- Processing and Communication hog the battery resource, drive handheld at peak power.
- No-op just needs idle power.
- Inability to run power hungry applications.



Proliferation of Wireless access points.

• Wi-Fi Hotspots are now available widely.

Scenario

- The handheld can connect wirelessly to remote servers and use their computing cycles.
- Opportunity to partly compute remotely and save power.
- Relevant program parts need to be transmitted DRAWBACK !!
- Look at overall energy savings.



- Is it possible to execute Java applications remotely and save power ?
- Execute Java applications partly remotely and reduce overall power.
 - Remote computation runs CPU at idle power.
 - Transmission needs additional energy.

□ We form a partition of the Java program.

- Distribute objects between local and remote machines.
- Local objects need peak CPU power.
- Remote objects need idle CPU power.
- Remote object creation needs transmission.



- Problem Statement
- **•** What is a Java program partition ?
- **D** Partitioning Algorithm I OPTIMAL
 - Based on MIN-CUT and OPTIMAL
- **Drawbacks of Algorithm I**
- **D** Partitioning Algorithm II FAST
 - Greedy Algorithm and APPROXIMATE.
- Experimental Results



HANDHELD







Object

Communication

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Executing a Partitioned Java Application

HANDHELD SERVER

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Partitioning Algorithms Overview

Profiling Step

- Number of objects of each class created.
- The amount of communication involved between object pairs.
- The CPU time spent in the context of each object.

Partitioning Step – Graph Representation

- OPTIMAL Method
 - Gives optimal partitions, very poor running time.
- FAST Method
 - Does not give the best partitions but is very fast.

Executing the Application as per the Partition



- The number of objects of each class created.
- The amount of communication involved between object pairs.
 - This gives an estimate of the transmission energy needed if these object pairs were separated.

Energy Spent= Transmission Energy per byte X Bytes Transmitted

- The CPU time spent in the context of each object.
 - This gives an estimate of the CPU energy saved if this object was executed remotely.

Energy Saved= (Peak - Idle) **X** CPU Time taken



Partitioning Algorithm I OPTIMAL





- The partition outputs a ratio a : b of objects of every class.
- For every 'a' objects that are placed in the server, we place 'b' objects in the client.
- This way, a different execution instance can still use the partition.



Graph Representation

- Each node is an object.
- Edges represent communication [C-Edge] and computation energies [P-Edge].

Partitioning Algorithm

• MIN-CUT of this graph will give the OPTIMAL partition – expressed as a ratio.

Time and Space Complexity

- SPACE : Graph is as large as the number of objects created in the application, HUGE !!
- TIME : runs in time cubic in the number of objects created, HUGE !!



- Partitioning is dependent on network variables and handheld variables.
 - Bandwidth.
 - Communication Energy per byte.
 - Battery Power.
- Change of variables requires partitioning to be done again.
- TIME SPACE inefficient algorithm might in itself need a lot of energy.



- Problem Statement
- **\Box** What is a partition ? \checkmark
- Partitioning Algorithm I OPTIMAL
 - Based on MIN-CUT and OPTIMAL
- Drawbacks of Algorithm I \checkmark
- **D** Partitioning Algorithm II FAST
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Graph Representation – Minimize space

- Each node is a class.
- Node weight is the number of objects of that class.
- Edges represent communication [C-Edge] and computation energies [P-Edge], as before.

Partitioning Algorithm – Minimize time

• A greedy style algorithm gives a partition that tends to reduce the energy spent.



Partition Algorithm II



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Benefit in moving an object to the server. **Benefit = (Peak Energy – Idle Energy)** Cost in moving an object to the server. **Cost = Transmission Energy Spent** $Cost = ((n_A - 1) * O_{AA} + n_B * O_{AB}) * \alpha$ **O**_{AA} is the average bytes transferred between objects of Class A.

If Benefit > Cost, the object can be moved to the server side.



- 1. Compute benefit [v] and cost [w] for all objects.
- 2. While (there exists some object whose v/w > 0) do
- 3. Choose the class whose object has the maximum v/w ratio.

4.	if max	f maximum v/w > 1 then			
5.		Move all objects of this class.			
6.		recompute v/w for all objects, CHECKPOINT.			
	else				
		repeat			
7.		move one object to the server.			
8.		recompute v/w for all objects.			
9.		Find maximum v/w			
10.		until v/w >1			
	end-if.				
	end-while				



Partition Algorithm II FAST

TIME Complexity

• runs in time linear in the number of objects created in the worst-case.

SPACE Complexity

• Number of node and edges in the graph is a function of the number of classes, SMALL !!.



Java Application that finds a sub-matrix within a Large Matrix

• Image Processing involves similar operations.

Number of	Time Taken in secs.		Energy Savings %	
Objects	OPTIMAL	FAST	OPTIMAL	FAST
3111	16.7	0.188	43.0	29.29
3111	16.7	0.188	44.08	29.29
3111	18.3	0.172	44.6	29.29
4443	95.1	0.133	43.6	32.22
6663	324	0.344	44.5	33.88

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Java Application that finds strings within large text documents

• Important operation in word-processing applications.

Number of	Time Taken in secs.		Energy Savings %	
Objects	OPTIMAL	FAST	OPTIMAL	FAST
523	1.62	0.003	45.29	37.84
1043	2.03	0.003	44.25	40.8
2043	8.94	0.06	42.1	38.06
3063	36.58	0.128	46.02	38.53
4563	41	0.172	47.5	43.06

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Graphs for Energy Savings



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- Energy savings in embedded devices by running Java programs partly remotely.
- We have presented an optimal algorithm that generates good partitions.
- We have presented a fast algorithm that can partition quickly.
- Experiments show that the fast algorithm is reasonably good in practice.