# Distributed Algorithms in Dynamic Networks (a short introduction) 

Arnaud Casteigts

The content of this talk can be found in
https://arxiv.org/abs/1807.07801

February 8, 2019

## Networks?

- Set of nodes V (a.k.a. entities, vertices)
- Set of links $E$ among them (a.k.a. relations, edges) $\rightarrow$ A network (or graph) $G=(V, E)$



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Complex networks
$\rightarrow$ compute global metrics
$\rightarrow$ explain and reproduce phenomena


Communication networks
$\rightarrow$ design interactions among entities
$\rightarrow$ study what can be done from within


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$\rightarrow$ study what can be done from within
$\rightarrow$ distributed algorithms...


## Distributed Algorithms



Collaboration of distinct entities to perform a common task.
No centralization available. No global knowledge.
(Think globally, act locally)

## Examples of problems

Broadcast


Propagating a piece of information from one node to all others.
$\longrightarrow$

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Election


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Distinguishing exactly one node among all.


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Spanning tree


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Determining how many participants there are.


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Consensus, naming, routing, exploration, ...

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## Dynamic Networks



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- Nature of the system. Change is normal.
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## Basic graph concepts



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$\rightarrow$ Paths become temporal (journey)

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\text { Ex: }\left(\left(a c, t_{1}\right),\left(c d, t_{2}\right),\left(d e, t_{3}\right)\right) \text { with } t_{i+1} \geq t_{i} \text { and } \rho\left(e_{i}, t_{i}\right)=1
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$G_{1}$
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$\rightarrow$ Hierarchy of classes of dynamic networks (basic examples)
$1-*$
$\mathcal{E}^{1 \forall}$
$1 \stackrel{s t}{\stackrel{s t}{*}} *$
$\mathcal{J}^{1 \forall>}$
$1 \rightsquigarrow *$
$\mathcal{J}^{1 \forall}$
$\underset{*-*}{\mathcal{K}}$

$$
\begin{aligned}
& \mathcal{T C} \mathcal{C}^{>} \\
& * \stackrel{s t}{\sim}
\end{aligned}
$$

$$
\underset{* \sim *}{\mathcal{T C}}
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\mathcal{J}^{\forall 1}
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## Classes of dynamic networks

## Network algorithms



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## Zoom: Optimal broadcast in DTNs?

What optimality ?
(Bui-Xuan, Jarry, Ferreira, 2003)


Which way is optimal from a to $d$ ?

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$\rightarrow$ Computing shortest, foremost and fastest journeys in dynamic networks

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Difficulty: Depends on the starting date! (E.g. for in the foremost case)

(a) The network.

(b) Foremost broadcast trees from a depending on the starting date.

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What about the distributed version?
$\rightarrow$ Can we broadcast a message to all the nodes in a foremost, shortest, or fastest way? (with termination detection at the emitter and without knowing the schedule)

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$\mathcal{T C}^{\mathcal{R}}:=$ All nodes can reach each other through journeys infinitely often (Formally, $\mathcal{T C}^{\mathcal{R}}:=\forall t, \mathcal{G}_{[t,+\infty)} \in \mathcal{T C}$ )

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Ex: MinimalDominatingSet and MaximalindependentSet


Robust MDS


Non robust MDS



Local algorithm for $\forall$ MIS

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Centralized algorithms

Classes of dynamic networks

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Centralized algorithms
Movement synthesis

## Algorithmic movement synthesis

1) Collective movements which induce temporal structure

$\rightarrow$ Synthesizing collective movements (a.k.a. mobility models) that satisfy temporal properties on the resulting communication graph (combined with a target mission).
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Discrete acceleration models
Vector Racer
(paper and pencil game)


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Vector Racer (paper and pencil game)

$\rightarrow$ Impact on problems, e.g. TSP $\nearrow$


Acceleration does impact the visit order!

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