Distributed Algorithms in Dynamic Networks (a short introduction)

Arnaud Casteigts

The content of this talk can be found in

https://arxiv.org/abs/1807.07801

February 8, 2019

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- Set of nodes V (a.k.a. entities, vertices)
- Set of links E among them (a.k.a. relations, edges)

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Complex networks

- \rightarrow compute global metrics
- \rightarrow explain and reproduce phenomena



Communication networks

- \rightarrow design interactions among entities
- \rightarrow study what can be done from within



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 \rightarrow distributed algorithms...

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Distributed Algorithms



Collaboration of distinct entities to perform a common task.

No centralization available. No global knowledge.

(Think globally, act locally)

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Broadcast

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Propagating a piece of information from one node to all others.



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Propagating a piece of information from one node to all others.



Election

Distinguishing exactly one node among all.

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Spanning tree

Selecting a cycle-free set of edges that interconnects all nodes.

Broadcast

Propagating a piece of information from one node to all others.

Election

Distinguishing exactly one node among all.

Spanning tree

Selecting a cycle-free set of edges that interconnects all nodes.

Counting

Determining how many participants there are.

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Counting

Consensus, naming, routing, exploration, ...

Determining how many participants there are.

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In fact, highly dynamic networks.

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How changes are perceived?

- Faults and Failures?
- Nature of the system. Change is normal.
- Partitioned network

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Example of scenario (say, exploration by mobile robots)

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As a sequence

Sequence of static graphs $\mathcal{G} = \mathcal{G}_0, \mathcal{G}_1, \dots$ [+table of dates]

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With a presence function $\mathcal{G} = (V, E, \mathcal{T}, \rho),$ with ρ being a presence function $\rho : E \times \mathcal{T} \to \{0, 1\}$

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 \rightarrow Paths become temporal (*journey*)

Ex: $((ac, t_1), (cd, t_2), (de, t_3))$ with $t_{i+1} \ge t_i$ and $\rho(e_i, t_i) = 1$

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What optimality ?

Which way is optimal from a to d?

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-min hop? (shortest)

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(Bui-Xuan, Jarry, Ferreira, 2003)

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Which way is optimal from *a* to *d*? -min hop? (shortest) -earliest arrival? (foremost) -fastest traversal? (fastest)

 \rightarrow Computing shortest, foremost and fastest journeys in dynamic networks

Which way is optimal from *a* to *d*? -min hop? (shortest) -earliest arrival? (foremost) -fastest traversal? (fastest)

Difficulty: Depends on the starting date! (E.g. for in the foremost case)

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 \rightarrow Can we broadcast a message to all the nodes in a foremost, shortest, or fastest way? (with termination detection at the emitter and without knowing the schedule)

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Ex: MINIMALDOMINATINGSET and MAXIMALINDEPENDENTSET

Robust MDS Non robust MDS

Network algorithms

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Algorithmic movement synthesis

1) Collective movements which induce temporal structure

 \rightarrow Synthesizing collective movements (*a.k.a. mobility models*) that satisfy temporal properties on the resulting communication graph (combined with a target mission).

 $\leftarrow \mathsf{Ex: this network} \in \mathcal{E^R}$ (Credit video: Jason Schoeters / JBotSim)

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 \rightarrow Impact on problems, e.g. $\mathrm{TSP}\nearrow$

Acceleration does impact the visit order!

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