Computing Parameters of Sequence-based Dynamic Graphs

Ralf Klasing

LaBRI, CNRS, University of Bordeaux, France

**This is a joint work with Arnaud Casteigts, Yessin M. Neggaz, and Joseph G. Peters.

















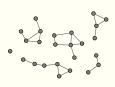








- How changes are perceived?
 - Faults and Failures?
 - Nature of the system
 - Change is normal









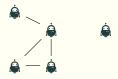






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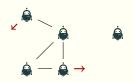






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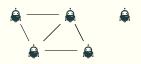






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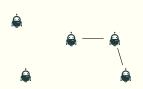






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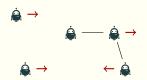






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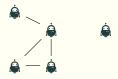




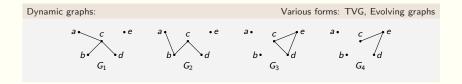


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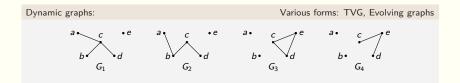


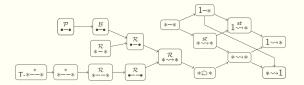
Dynamic Graphs



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Dynamic Graphs





Dynamic graphs classes: [Casteigts, Flocchini, Quattrociocchi et Santoro, 2011]

$$\mathcal{G} = (V, E_i)$$

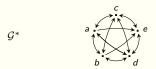
$$\stackrel{a \leftarrow c \leftarrow e}{\downarrow b \leftarrow G_1} \stackrel{a \leftarrow c \leftarrow e}{\downarrow b \leftarrow G_2} \stackrel{a \leftarrow c \leftarrow e}{\downarrow b \leftarrow G_3} \stackrel{a \leftarrow c \leftarrow e}{\downarrow b \leftarrow G_4} \stackrel{a \leftarrow c \leftarrow e}{\downarrow b \leftarrow G_4}$$

Temporal connectivity $\iff \forall u, v \in V, u \rightsquigarrow v.$

$$\mathcal{G} = (V, E_i)$$

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- **Temporal connectivity** $\iff \forall u, v \in V, u \leadsto v.$
- Transitive closure of the journeys: reachability over time [Bhadra and Ferreira, 2003]



 $\mathcal G$ is temporally connected \Leftrightarrow Transitive closure $\mathcal G^*$ is complete









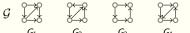






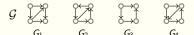


- ➡ High-level strategies that work directly at the graph level
- Elementary graph-level operations





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- Elementary graph-level operations

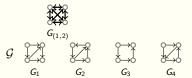


TEMPORAL-DIAMETER

Finding the *temporal diameter* of a given dynamic graph \mathcal{G} , i.e. the smallest duration in which there exists a journey from any node to all other nodes.



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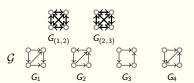


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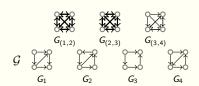
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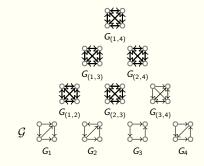




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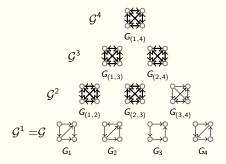
- Transitive closures
- Completeness test



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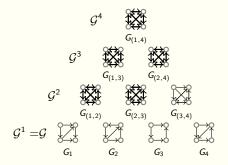
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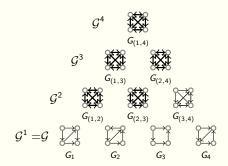


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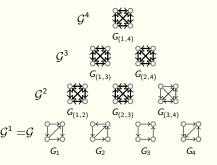
Finding the *temporal diameter* of a given dynamic graph \mathcal{G} , i.e. the smallest duration in which there exists a journey from any node to all other nodes.

Finding the smallest d such that every super node in row \mathcal{G}^d is a complete graph (i.e. every subsequence of length d is temporally connected).

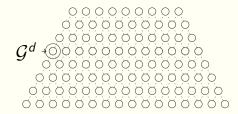
- Transitive closures
- Completeness test
- Transitive closures concatenation



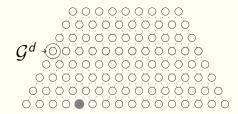
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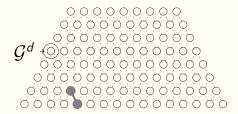
Decision version (given d)



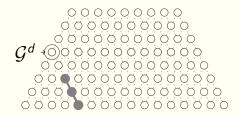
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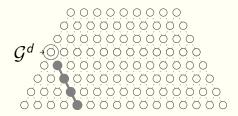
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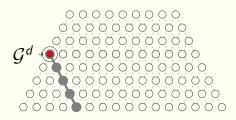


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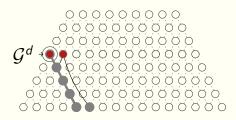
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A ladder of length / costs I − 1 concatenation



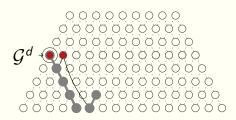
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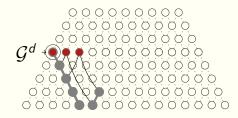
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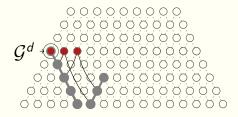
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- A ladder of length / costs / − 1 concatenation
- Use left and right ladders



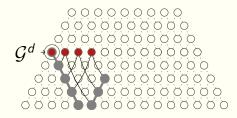
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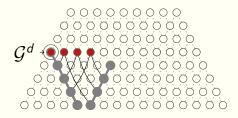
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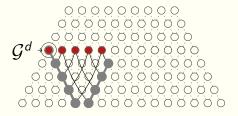
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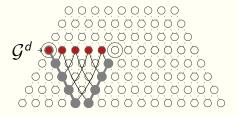
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- Any graph "between" two ladders (red graphs) can be computed by a single binary concatenation



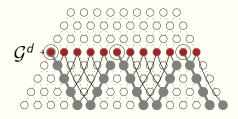
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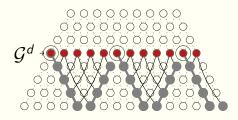
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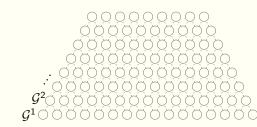
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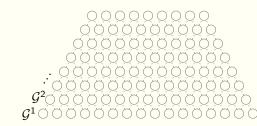
 $O(\delta)$ elementary operations per row

Minimization version (find the temporal diameter d)

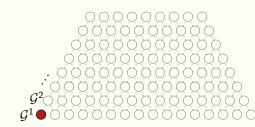
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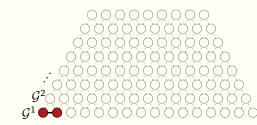
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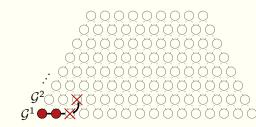
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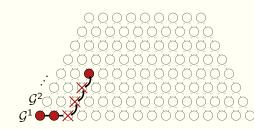


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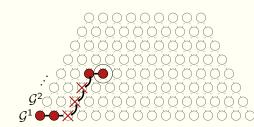
Minimization version (find the temporal diameter d)

Strategy: ascending walk



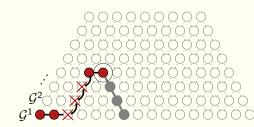
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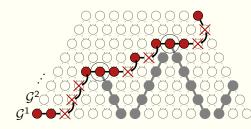
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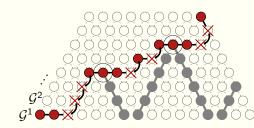
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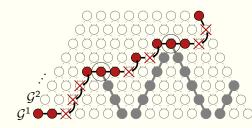
Minimization version (find the temporal diameter d)

- Strategy: ascending walk
- The total length of the ladders is $O(\delta)$
- At most O(δ) binary concatenation and completeness tests



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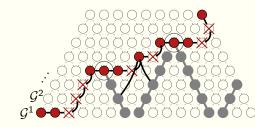
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Disjointness property: $cat(G_{(i,j)}, G_{(i',j')}) = G_{(i,j')}$

Minimization version (find the temporal diameter d)

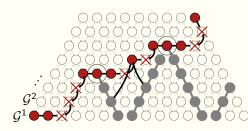
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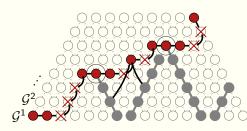
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- If $G_{(i,j)}$ is complete, then $G_{(i',j')}$ is complete, for all $i' \leq i$ and $j' \geq j$

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Temporal-Diameter is solvable with $O(\delta)$ elementary operations

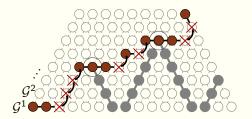
Online Algorithms

Online Algorithms

- The optimal algorithms can be adapted to an online setting
- The sequence of graphs $G_1, G_2, G_3, ...$ of \mathcal{G} is processed in the order of reception
- **Amortized cost of** O(1) elementary operations per graph received
- Dynamic version: consider only the recent history

Solve other problems using the same framework

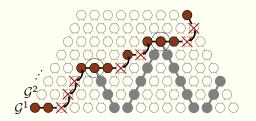
Solve other problems using the same framework



Framework generalization

- Transitive closures concatenation
- Completeness test
- Transitive closure

Solve other problems using the same framework

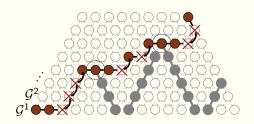


Framework generalization

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Composition operation
Test operation
Super node

Solve other problems using the same framework



Minimization problems

Find the smallest value

Framework generalization

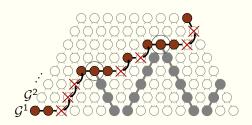
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Composition operation Test operation

Super node

13

Solve other problems using the same framework



Minimization problems

Maximization problems

Find the largest value

Framework generalization

Find the smallest value

- Transitive closures concatenation
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 \rightarrow

Composition operation Test operation

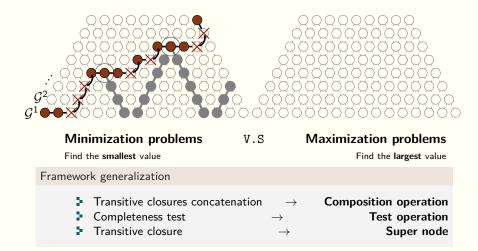
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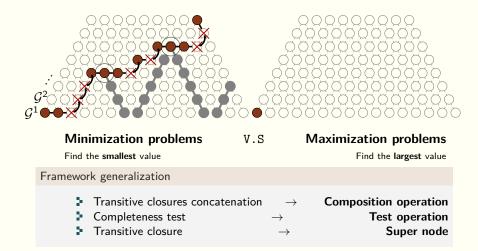
Temporal Connectivity

V.S

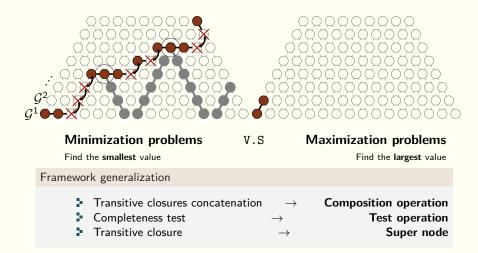
Solve other problems using the same framework



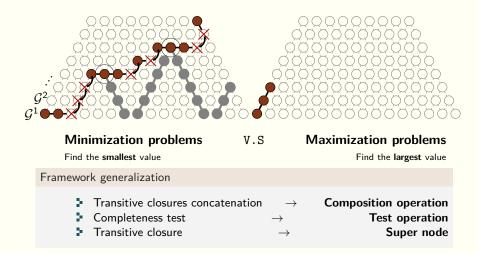
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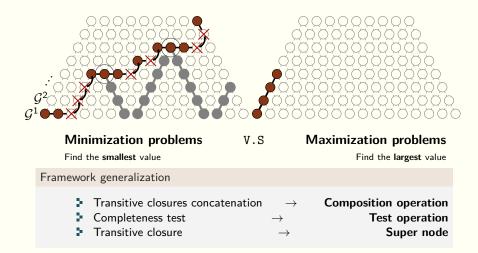
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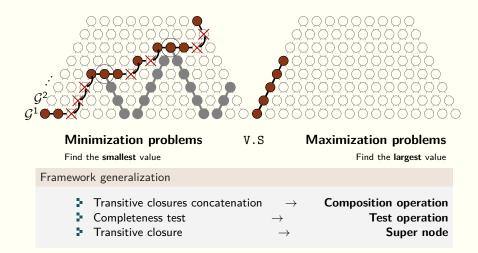
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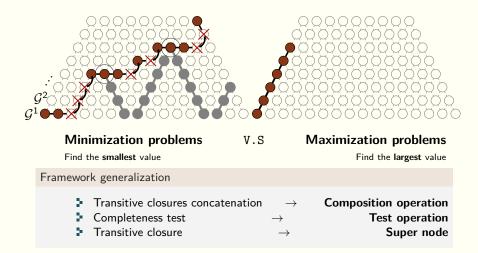
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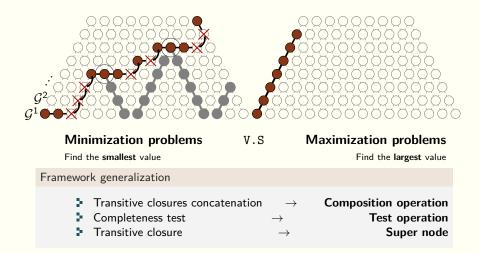
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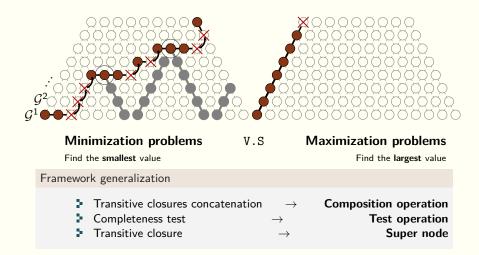
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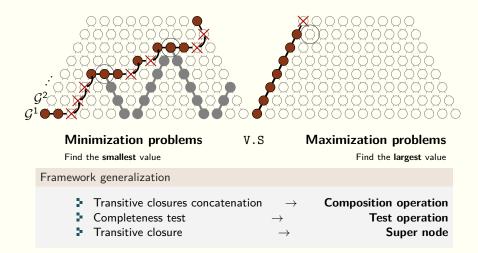
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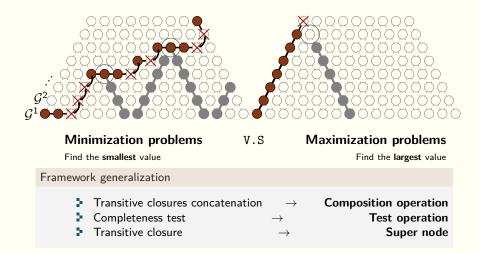
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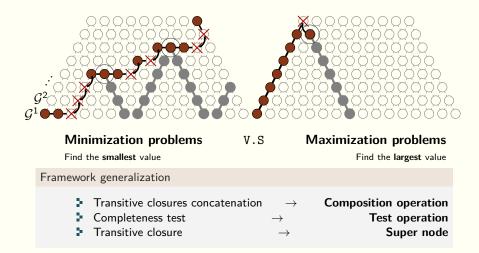
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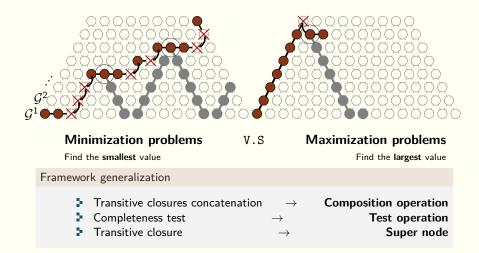
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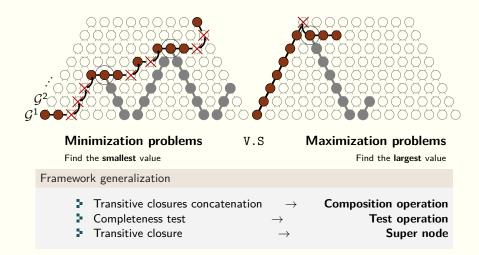
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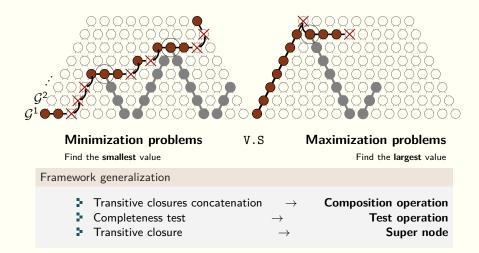
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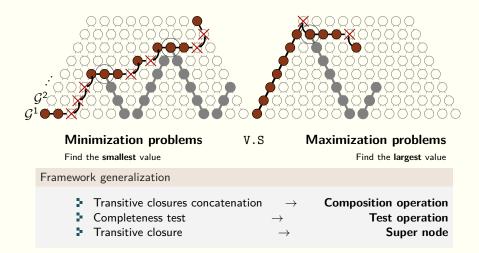
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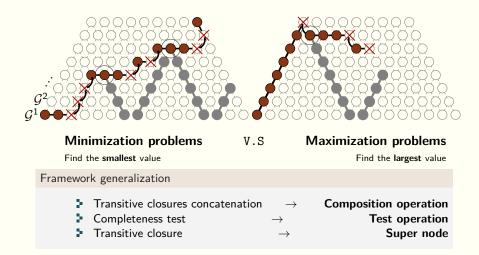
Solve other problems using the same framework



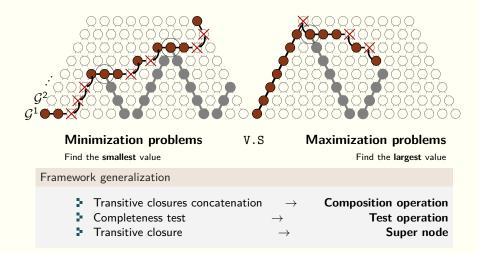
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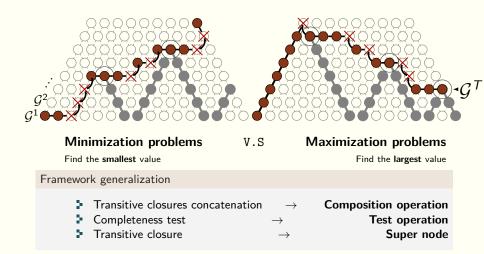
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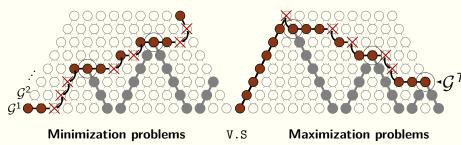
Solve other problems using the same framework



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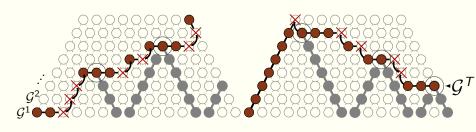
Solve other problems using the same framework



Find the smallest value

Find the largest value

Solve other problems using the same framework



Minimization problems

Find the smallest value

V.S

Maximization problems

Find the largest value

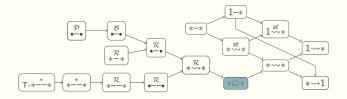
Requirements

- test $(G_{(i,j)}) = true \Leftrightarrow \{G_i, G_{i+1}, \dots, G_j\}$ satisfies the property P
- The composition operation is associative
- Only minimization: If $test(G_{(i,j)}) = true$ then $test(G_{(i',j')}) = true, \forall i' \leq i, j' \geq j$
- Only maximization: If $test(G_{(i,j)}) = true$ then $test(G_{(i',j')}) = true, \forall i' \geq i, j' \leq j$

Round-trip Temporal Connectivity

Round-trip Temporal Connectivity

A dynamic graph ${\cal G}$ is round-trip temporal connected if and only if a back-and-forth journey exists from any node to all other nodes.



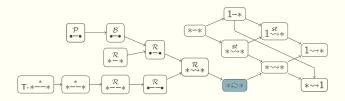
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ROUND-TRIP-TEMPORAL-DIAMETER(minimization)

Finding the smallest duration in which there exists a back-and-forth journey from any node to all other nodes.



Round-trip Temporal Connectivity

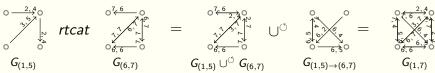
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ROUND-TRIP-TEMPORAL-DIAMETER(minimization)

Finding the smallest duration in which there exists a back-and-forth journey from any node to all other nodes.

- Super node: Round-trip transitive closure
- **Composition operation:** Round-trip transitive closure concatenation

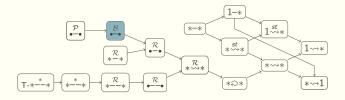


Test operation: Round-trip completeness

Bounded Realization of the footprint

Time-bounded edge reappearance

A dynamic graph $\mathcal G$ has a time-bounded edge reappearance with a bound b if the time between two appearances of the same edge is at most b.



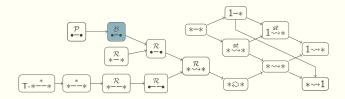
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BOUNDED-REALIZATION-OF-THE-FOOTPRINT(minimization)

Finding the smallest b such that in every subsequence of length b in the sequence \mathcal{G} , all the edges of the footprint appear at least once.



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BOUNDED-REALIZATION-OF-THE-FOOTPRINT(minimization)

Finding the smallest b such that in every subsequence of length b in the sequence G, all the edges of the footprint appear at least once.

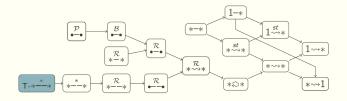
Super node: Union graphs

Composition operation: Union

Test operation: Equality to the footprint

Definition: *T*-interval connectivity

A dynamic graph $\mathcal G$ is T-interval connected if and only if every T length sequence of graphs has a common connected spanning sub-graph.



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T-Interval-Connectivity (maximization)

Finding the largest T for which the graph is T-interval connected.



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$$\mathcal{G}^2$$
 \mathbb{Z} \mathbb{Z}

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Finding the largest T for which the graph is T-interval connected.

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- Composition operation: Intersection
 - **Test operation:** Connectivity test

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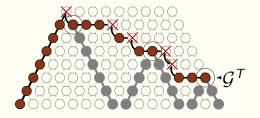
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Symmetric Problems

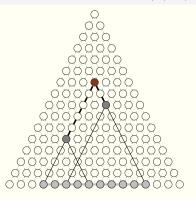
Symmetric Problems

A minimization or maximization problem is symmetric if: for all $i, j, i', j' \leq \delta$, $i \leq i' \leq j$, $composition(G_{(i,j)}, G_{(i',j')}) = G_{(i,j')}$.

Symmetric Problems

Symmetric Problems

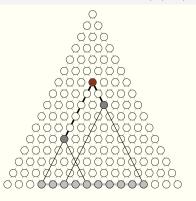
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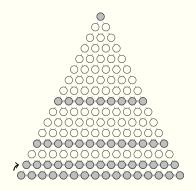
e.g T-Interval-Connectivity and Bounded-Realization-of-the-Footprint

Row-Based Strategy

Symmetric problems (maximization)

Symmetric problems (maximization)

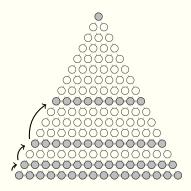
Symmetric problems (maximization)



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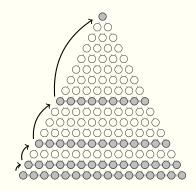
Symmetric problems (maximization)

- $O(\delta)$ composition per row
- $O(\delta)$ tests per row



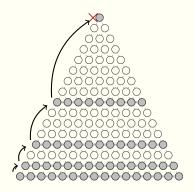
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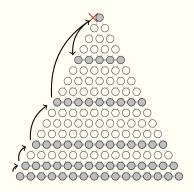
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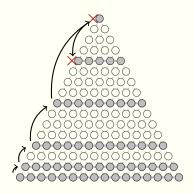
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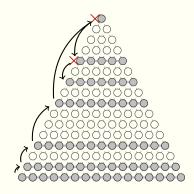
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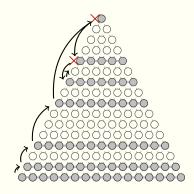
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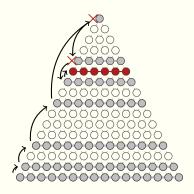
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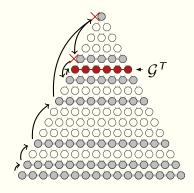
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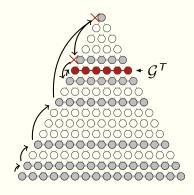
Symmetric problems (maximization)

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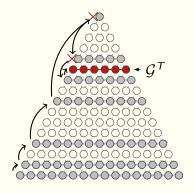
Symmetric problems (maximization)

- $O(\delta)$ composition per row
- $O(\delta)$ tests per row
- $O(\log \delta)$ rows



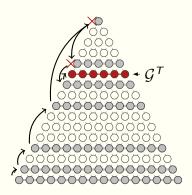
Symmetric problems (maximization) $O(\delta \log \delta)$ elementary operations

- $O(\delta)$ composition per row
- $O(\delta)$ tests per row
- $O(\log \delta)$ rows



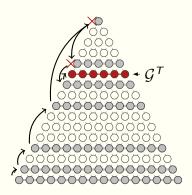
Parallel Version

On EREW PRAM



Parallel Version

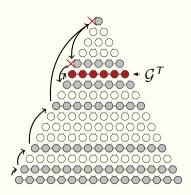
On EREW PRAM



Parallel Version

On EREW PRAM

Symmetric problems are solvable in $O(\log^2 \delta)$ on an **EREW PRAM** with $O(\delta)$ processors



Conclusion

Conclusion

- High-level strategies for computing minimization and maximization parameters
- Algorithms that use only $O(\delta)$ elementary operations
- Parallel versions on PRAM (in Nick's class)
- Online algorithms with amortized cost of O(1) elementary operations per graph received
- Perspectives
 - How about other classes?
 - Generic Framework
 - What if the evolution of the dynamic graph is constrained?

Thank you!