## Graph Exploration by Oblivious Agents

David ILCINKAS

CNRS, Bordeaux, France

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### **Problem**

#### Model/context

- Team of mobile entities
  - sensing the environment by taking a snapshot of it
  - that do not communicate
  - that are anonymous and oblivious

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#### In this talk: exploration

Each node must be visited.

## The Look-Compute-Move cycle

#### Look

The agent takes an egocentric instantaneous snapshot of the network and its agents

#### Compute

Based on this observation, it decides to stay idle or to move to some neighboring node.

#### Move

In the latter case it <u>instantaneously</u> moves towards its destination.

## Identical oblivious asynchronous agents

#### Identical

Agents have no IDs. They execute the same program.

#### **Oblivious**

The agents have no memory of observations, computations and moves made in previous cycles.

#### Possibly asynchronous

The time between Look, Compute, and Move operations is finite but unbounded.

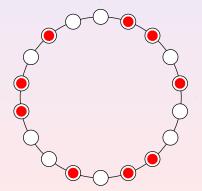
Reminder:

#### Non-communicating

No communication mechanisms between agents, even locally.

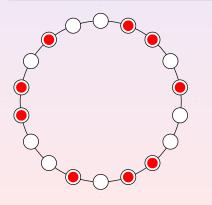
### Anonymous setting

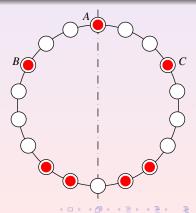
Agents and nodes have no IDs



### Anonymous setting

- Agents and nodes have no IDs
- There are no port numbers, and no chirality



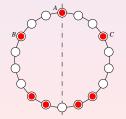


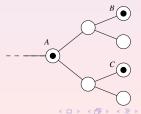
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#### In case of symmetry

- Compute : choice of an equivalence class of neighbors
- Actual choice: made by the adversary (i.e. worst case)





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#### Weak multiplicity detection

"zero", "one", or "more than one" agents



### Two types of exploration

#### Exploration with stop

- Each node must be visited by at least one agent.
- All agents must stop after finite time.

- Each node must be visited by every agent infinitely ofter
- Agents must not collide (meet at a node or cross each other on an edge).
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### Usual focus

#### n: number of nodes

#### **Smallest** exploring team

 $\rho^-(n) = \text{Minimum}$  number of agents that can explore any n-node graph of a given family.

#### Largest exploring team

 $\rho^+(n) = \text{Maximum}$  number of agents that can explore any n-node graph of a given family.

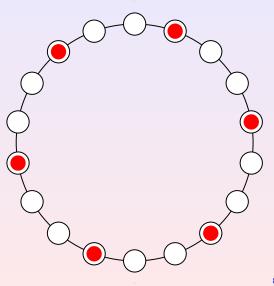
Families: Rings, (Partial) Grids, Trees, Tori, etc.

## A simple impossibility result

#### Observation (in rings)

Impossible for both types of exploration when *k* divides *n*.

n: number of nodesk: number of agents

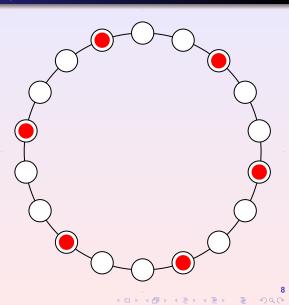


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## Focus on Exploration with stop

#### Exploration with stop (reminder)

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### Family: Rings

- Additional assumption
- n and k are coprime.
  - n: number of nodes
- k: number of agents



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Family: Rings

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n and k are coprime.

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## Some additional definitions

#### Interdistance

Minimum distance taken over all pairs of distinct agents.

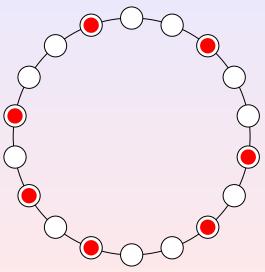
Here interdistance=2

Maximal set of agents of size at least 2

forming a line with an

agent every *d* nodes.

(*d*=interdistance)

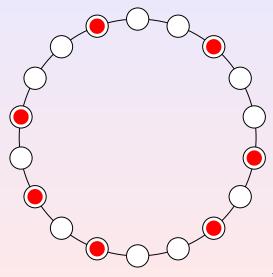


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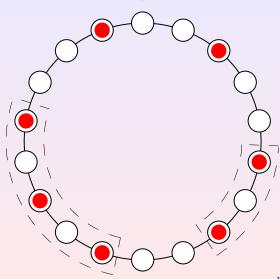
#### Interdistance

Minimum distance taken over all pairs of distinct agents.

Here interdistance=2.

#### Block

Maximal set of agents, of size at least 2, forming a line with an agent every d nodes. (d=interdistance)



## The algorithm

#### Set-Up Phase

Goal: to transform the (arbitrary) initial configuration into a configuration of interdistance 1 where there is a single block or two blocks of the same size.

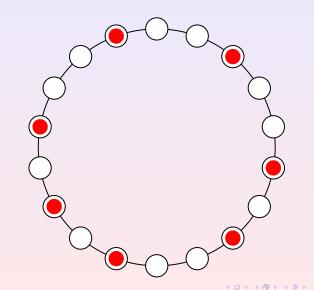
Method: decrease the number of blocks whenever possible. Otherwise, decrease the interdistance.

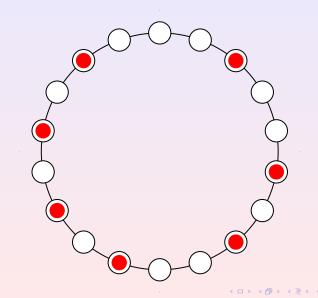
#### Tower-Creation Phase

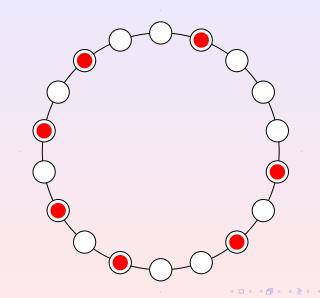
Goal: to create one or two multiplicities (towers) inside each block; furthermore a number of agents become uniquely identified as explorers.

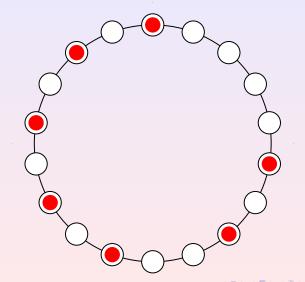
#### **Exploration Phase**

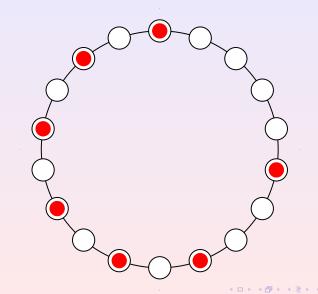
Goal: to perform exploration thanks to the explorers until reaching an identified final configuration.

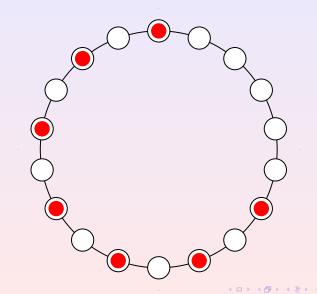


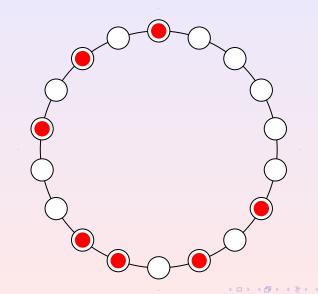


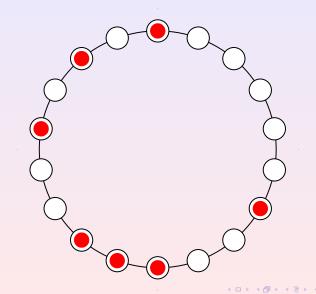


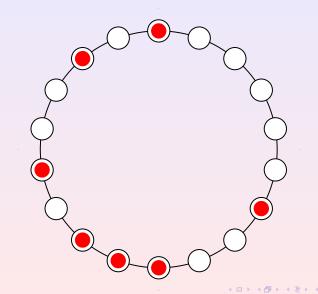


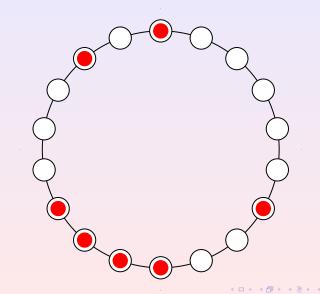


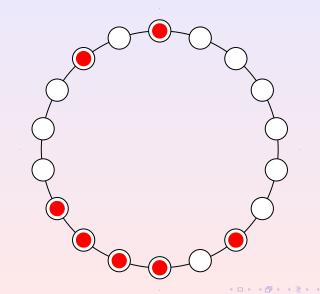


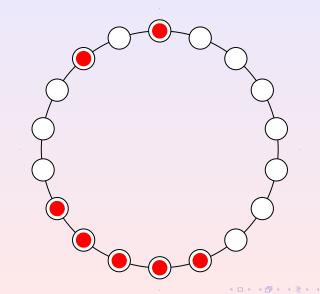


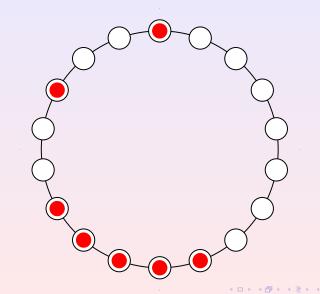


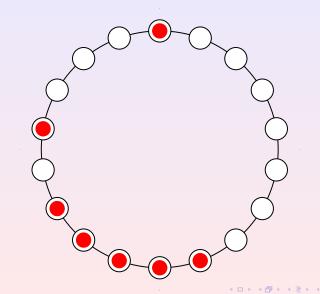


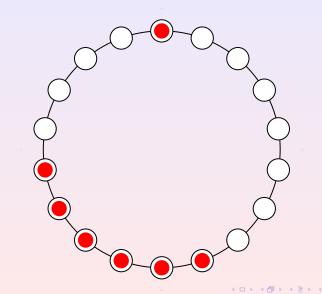


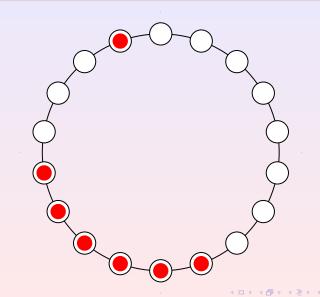


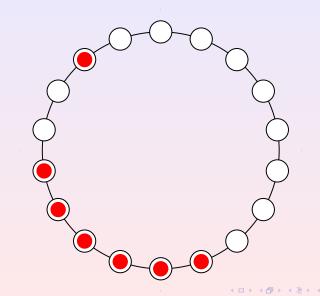


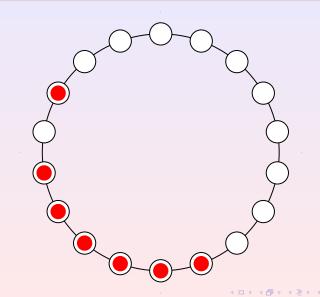


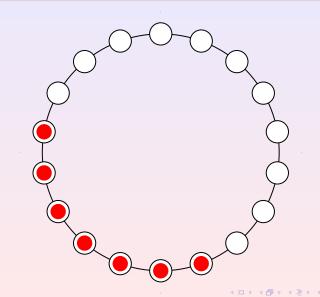


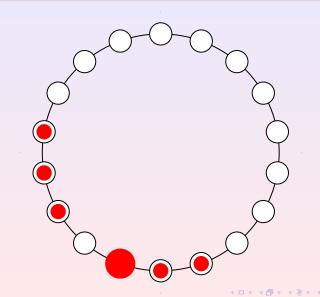


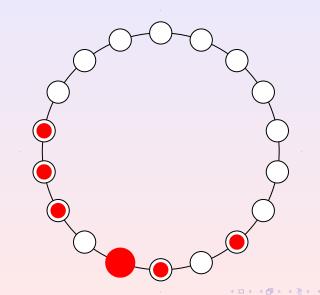


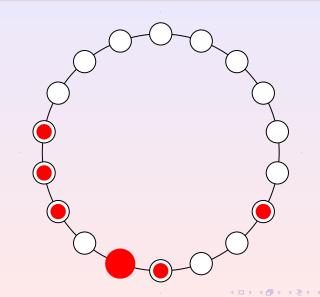


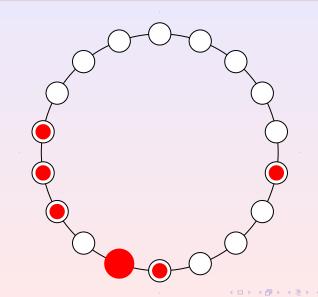


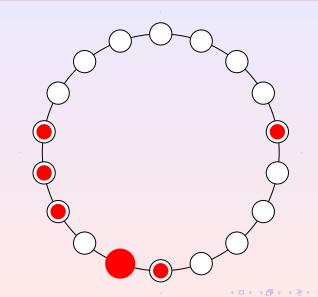


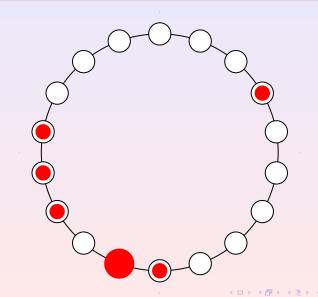


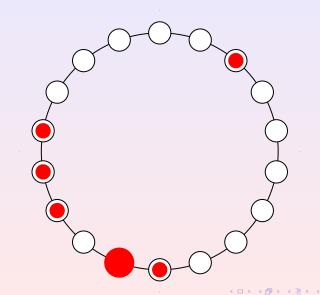


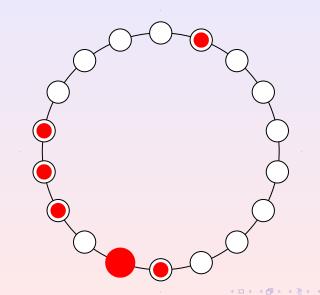


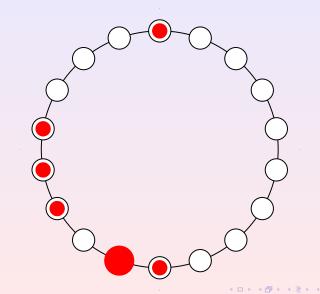


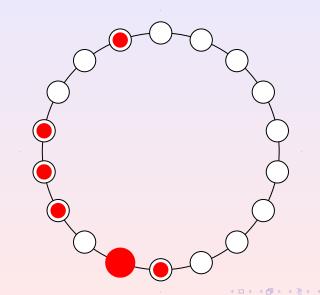


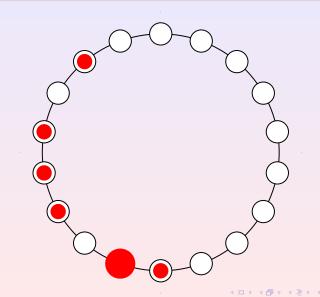


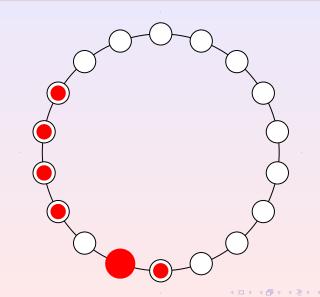












# Algorithmic techniques (exploration with stop)

#### Phase 1

Form a special configuration without multiplicities.

• Uses (co-)primality of n and k.

Phase 2

Create a multiplicity to mark beginning of exploration

Phase 3

Perform exploration by few (1-3) agents.

Uses multiplicity and other agents to

break symmetry

keep track of the progress

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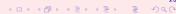
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#### Phase 3

Perform exploration by few (1–3) agents.

- Uses multiplicity and other agents to
  - break symmetry
  - keep track of the progress

But symmetry and asynchrony



## Models of synchrony

### **FSYNC** (fully synchronous)

All agents execute their Look-Compute-Move simultaneously.

### SSYNC (semi-synchronous)

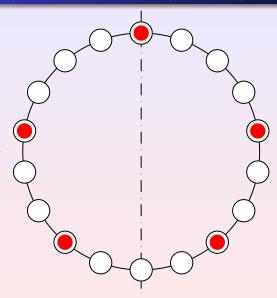
Look-Compute-Move is instantaneous.

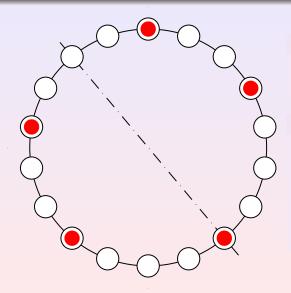
⇒ Difficult to move multiplicities

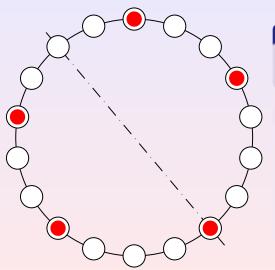
### ASYNC (asynchronous)

The time between Look, Compute, and Move operations is finite but unbounded.

⇒ Pending moves (moves based on outdated snapshots)





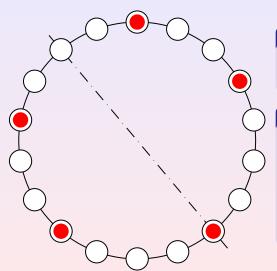


### Guideline 1

Move only one agent at a time if possible.

f the current

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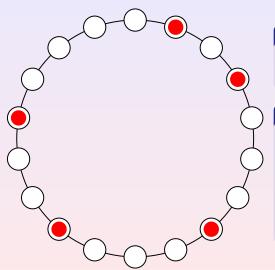


#### Guideline 1

Move only one agent at a time if possible.

#### Guideline 2

If the current configuration could come from a symmetric configuration, then move the "risky" agent first.



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## Known results (explo. with stop) 1/3

### In the rings

- [1] ASYNC det.:  $\rho^-(n)$  is in  $O(\log n)$
- [1] SSYNC det.:  $\rho^-(n)$  is in  $\Omega(\log n)$  for inf. many n
- [2] ASYNC det.:  $\rho^-(n) \le 5$  when n is not a multiple of 5
- [2] SSYNC det.:  $\rho^-(n) = 4$  when n is odd.
- [2] SSYNC det.:  $\rho^-(n) \ge 5$  if n is even
- [3] ASYNC prob.:  $\rho^-(n)$  is in  $\Omega(\log n)$  for inf. many n
- [3] SSYNC prob.:  $\rho^{-}(n) = 4$  for any n > 4
- [1] Flocchini, Ilcinkas, Pelc, Santoro. Algorithmica 2013
- [2] Lamani, Potop-Butucaru, Tixeuil. SIROCCO 2010
- [3] Devismes, Petit, Tixeuil. TCS 2013



# Known results (explo. with stop) 2/3

Limited visibility: snapshot up to distance  $\varphi$  only [Datta, Lamani, Larmore, Petit]

### In the rings, with arphi=1 (ICDCS 2013)

- SSYNC det.: impossible
- FSYNC det.:  $\rho^-(n) \geq 5$
- FSYNC det.: algo. k = 5 when initial config. is special

### In the rings, with arphi= 2 and 3 (IPDPS workshop 2015)

For  $\varphi = 2$ :

• ASYNC det.: algo. k = 7 when initial config. is special

For  $\varphi = 3$ :

• ASYNC det.:  $\rho^-(n) \le 7$  for any n > 21



## Known results (explo. with stop) 3/3

[Flocchini, Ilcinkas, Pelc, Santoro. TCS 2010]

- Trees max. deg. 3:  $\rho^-(n) \in \Theta(\log n / \log \log n)$
- Complete ternary trees:  $\rho^-(n) \in \Theta(n)$

[Flocchini, Ilcinkas, Pelc, Santoro. IPL 2011]

• Line: iff k = 3,  $k \ge 5$ , or k = 4 and n odd

[Devismes, Lamani, Petit, Raymond, Tixeuil. SSS 2012]

• Grid:  $\rho^-(n) = 3$  (SSYNC/ASYNC)

[Devismes, Lamani, Petit, Tixeuil. NETYS 2015]

• Torus: SSYNC  $\rho^-(n) = 4$  prob. ( $\geq 5$  det.)

[Chalopin, Flocchini, Mans, Santoro. WG 2010]

Arbitrary graphs with port numbers: rigid initial config.



## Focus on Exclusive Perpetual Exploration

### Exclusive perpetual exploration (reminder)

- Each node must be visited by every agent infinitely often.
- Agents must not collide (meet at a node or cross each other on an edge).
- Initial configurations are the configurations without multiplicities.

First studied family: Partial grids with NSEW

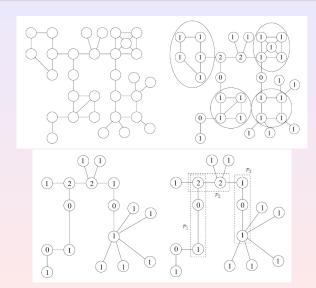
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First studied family: Partial grids with NSEW

## Mobility tree and parameter q



# Known results (exclusive perpetual explo.) 1/2

### In the partial grids (FSYNC)

[Baldoni, Bonnet, Milani, Raynal. IPL 2008]

- Definition of the parameter q
- Solvable iff  $k \le n q$

[Baldoni, Bonnet, Milani, Raynal. OPODIS 2008]

- Limited visibility  $\varphi = 1$
- Solvable iff  $k \le n q$ , except for  $q = 0 \Longrightarrow n 1$

### In the (complete) grids

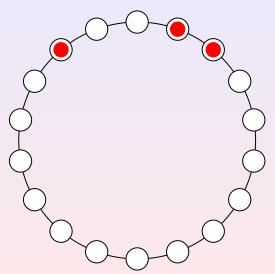
[Bonnet, Milani, Potop-Butucaru, Tixeuil. OPODIS 2011]

- ASYNC. no sense of direction
- Deterministically:  $\rho^-(n) = 3$

# Algorithmic techniques (exclusive perpetual explo.)

### Two phases

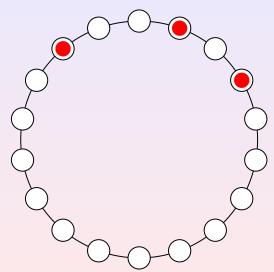
- Form a special asymmetric pattern
- Move the pattern through the graph



# Algorithmic techniques (exclusive perpetual explo.)

### Two phases

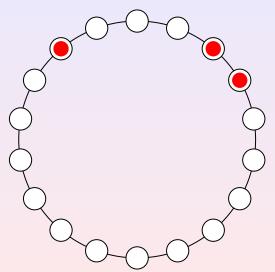
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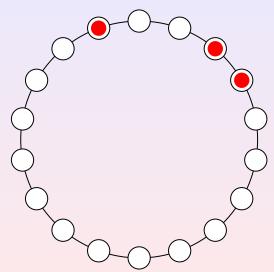
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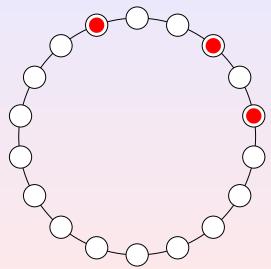
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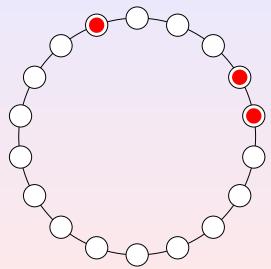
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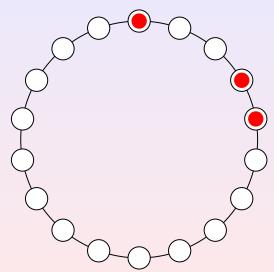
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## Known results (exclusive perpetual explo.) 2/2

## Case of the rings

[Blin, Milani, Potop-Butucaru, Tixeuil. DISC 2010]

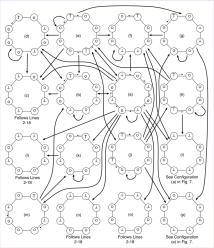
- FSYNC det.: Impossible when k is even
- ASYNC det.:  $\rho^-(n) = 3$  for  $n \ge 10$  and n and k coprime
- ASYNC det.:  $\rho^+(n) = n-5$  for k odd & n and k coprime [Bonnet, Défago, Petit, Potop-Butucaru, Tixeuil. SRDS Workshop 2014]
  - Generic method to list all protocols (SSYNC)
  - Application: impossible for k = 5 and n = 10

[D'Angelo, Di Stefano, Navarra, Nisse, Suchan. Algorithmica 2015]

- ASYNC but only for rigid initial configurations
- Det. algo for any k, n such that  $n \ge 10$  and  $5 \le k \le n 3$ , except the case k = 5 and n = 10

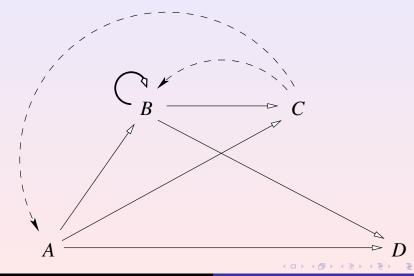
# Proofs of correctness (1/2)

## Case-by-case analysis



# Proofs of correctness (2/2)

DAG with potential functions



## Formal verification via model checking

[Bérard, Lafourcade, Millet, Potop-Butucaru, Thierry-Mieg, Tixeuil. Distributed Computing 2016

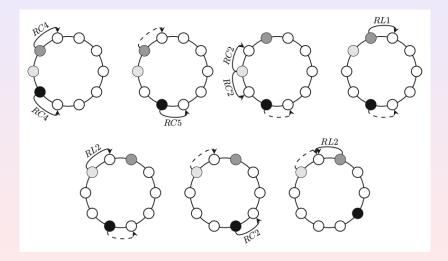
## Study of two algorithms

- Exploration with stop [FIPS Algorithmica 2013]
  - Correct for 17 < k < 21 and k < n < 22
  - Also correct for smaller values of k (5, 7 and 10)
- Exclusive perpetual exploration [BMPT DISC 2010]
  - Counter-example for k=3 and n=10

## [Doan, Bonnet, Ogata. SOFG+MSVL 2016]

- Also study of [BMPT DISC 2010]
- Same counter-example, with a different technique

## The counter-example



## Conclusion and perspectives

#### Conclusion

- Computability depends on the amount of symmetries
- Tricky to (correctly) handle symmetries and asynchrony

- Limited visibility
- Formal verification
  - (Parametrized) Model checking
- Proof assistant

## Conclusion and perspectives

#### Conclusion

- Computability depends on the amount of symmetries
- Tricky to (correctly) handle symmetries and asynchrony

#### Interesting perspectives

- Limited visibility
- Formal verification
  - (Parametrized) Model checking
  - Proof assistant

# Thank you for your attention