### Graph Exploration by Oblivious Agents

### David ILCINKAS

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David ILCINKAS Graph Exploration by Oblivious Agents

### Model/context

- Team of mobile entities
  - sensing the environment by taking a snapshot of it
  - that do not communicate
  - that are anonymous and oblivious

In a discrete environment: anonymous graphs

In this talk: exploration

Each node must be visited.

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#### Look

The agent takes an egocentric instantaneous snapshot of the network and its agents.

#### Compute

Based on this observation, it decides to stay idle or to move to some neighboring node.

#### Move

In the latter case, it <u>instantaneously</u> moves towards its destination.

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### Identical oblivious asynchronous agents

#### Identical

Agents have no IDs. They execute the same program.

#### Oblivious

The agents have no memory of observations, computations and moves made in previous cycles.

#### Possibly asynchronous

The time between Look, Compute, and Move operations is finite but unbounded.

Reminder:

Non-communicating

No communication mechanisms between agents, even locally.

### Anonymous setting

- Agents and nodes have no IDs
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### In case of symmetry

- Compute : choice of an equivalence class of neighbors
- Actual choice: made by the adversary (i.e. worst case)



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Weak multiplicity detection "zero", "one", or "more than one" agents		
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#### Exploration with stop

- Each node must be visited by at least one agent.
- All agents must stop after finite time.

#### Exclusive perpetual exploration

- Each node must be visited by every agent infinitely often.
- Agents must not collide (meet at a node or cross each other on an edge).

In both cases, the initial configurations are the configurations without multiplicities.

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#### *n*: number of nodes

#### <u>Smallest</u> exploring team

 $\rho^{-}(n) =$  Minimum number of agents that can explore any *n*-node graph of a given family.

#### Largest exploring team

 $\rho^+(n) = Maximum$  number of agents that can explore any *n*-node graph of a given family.

Families: Rings, (Partial) Grids, Trees, Tori, etc.

# A simple impossibility result

#### Observation (in rings)

Impossible for both types of exploration when k divides n.

n: number of nodes
k: number of agents



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### Focus on Exploration with stop

#### Exploration with stop (reminder)

- Each node must be visited by at least one agent.
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### Family: Rings

#### Additional assumption

n and k are coprime.

- *n*: number of nodes
- k: number of agents

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### Some additional definitions

#### Interdistance

Minimum distance taken over all pairs of distinct agents.

Here interdistance=2.

#### Block

Maximal set of agents, of size at least 2, forming a line with an agent every *d* nodes. (*d*=interdistance)



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#### Block

Maximal set of agents, of size at least 2, forming a line with an agent every d nodes. (d=interdistance)



# The algorithm

### Set-Up Phase

Goal: to transform the (arbitrary) initial configuration into a configuration of interdistance 1 where there is a single block or two blocks of the same size.

Method: decrease the number of blocks whenever possible. Otherwise, decrease the interdistance.

### Tower-Creation Phase

Goal: to create one or two multiplicities (towers) inside each block; furthermore a number of agents become uniquely identified as explorers.

### **Exploration Phase**

Goal: to perform exploration thanks to the explorers until reaching an identified final configuration.






























































# Algorithmic techniques (exploration with stop)

#### Phase 1

Form a special configuration without multiplicities.

• Uses (co-)primality of *n* and *k*.

#### Phase 2

Create a multiplicity to mark beginning of exploration.

#### Phase 3

Perform exploration by few (1-3) agents.

- Uses multiplicity and other agents to
  - break symmetry
  - keep track of the progress

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#### But symmetry and asynchrony

### Models of synchrony

#### **FSYNC** (fully synchronous)

All agents execute their Look-Compute-Move simultaneously.

#### <u>SSYNC</u> (semi-synchronous)

Look-Compute-Move is instantaneous.

 $\implies$  Difficult to move multiplicities

#### <u>ASYNC</u> (asynchronous)

The time between Look, Compute, and Move operations is finite but unbounded.

 $\implies$  Pending moves (moves based on outdated snapshots)

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#### Guideline 1

Move only one agent at a time if possible.

If the current configuration could come from a symmetric configuration, then move the "risky" agent first.



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# Known results (explo. with stop) 1/3

#### In the rings

- [1] ASYNC det.:  $\rho^-(n)$  is in  $O(\log n)$
- [1] SSYNC det.:  $\rho^-(n)$  is in  $\Omega(\log n)$  for inf. many n
- [2] ASYNC det.:  $\rho^{-}(n) \leq 5$  when *n* is not a multiple of 5
- [2] SSYNC det.:  $\rho^-(n) = 4$  when n is odd.
- [2] SSYNC det.:  $\rho^-(n) \ge 5$  if *n* is even
- [3] ASYNC prob.:  $\rho^-(n)$  is in  $\Omega(\log n)$  for inf. many n
- [3] SSYNC prob.:  $\rho^-(n) = 4$  for any n > 4

Flocchini, Ilcinkas, Pelc, Santoro. Algorithmica 2013
Lamani, Potop-Butucaru, Tixeuil. SIROCCO 2010
Devismes, Petit, Tixeuil. TCS 2013

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### Known results (explo. with stop) 2/3

Limited visibility: snapshot up to distance  $\varphi$  only [Datta, Lamani, Larmore, Petit]

#### In the rings, with $\varphi = 1$ (ICDCS 2013)

- SSYNC det.: impossible
- FSYNC det.:  $\rho^-(n) \ge 5$
- FSYNC det.: algo. k = 5 when initial config. is special

#### In the rings, with $\varphi = 2$ and 3 (IPDPS workshop 2015)

For  $\varphi = 2$ :

• ASYNC det.: algo. k = 7 when initial config. is special For  $\varphi = 3$ :

• ASYNC det.:  $\rho^-(n) \leq 7$  for any n > 21

### Known results (explo. with stop) 3/3

[Flocchini, Ilcinkas, Pelc, Santoro. TCS 2010]

- Trees max. deg. 3:  $\rho^-(n) \in \Theta(\log n / \log \log n)$
- Complete ternary trees:  $\rho^-(n) \in \Theta(n)$

[Flocchini, Ilcinkas, Pelc, Santoro. IPL 2011]

• Line: iff k = 3,  $k \ge 5$ , or k = 4 and n odd

[Devismes, Lamani, Petit, Raymond, Tixeuil. SSS 2012]

• Grid:  $\rho^-(n) = 3$  (SSYNC/ASYNC)

[Devismes, Lamani, Petit, Tixeuil. NETYS 2015]

• Torus: SSYNC  $\rho^{-}(n) = 4$  prob. ( $\geq 5$  det.)

[Chalopin, Flocchini, Mans, Santoro. WG 2010]

• Arbitrary graphs with port numbers: rigid initial config.

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### Focus on Exclusive Perpetual Exploration

#### Exclusive perpetual exploration (reminder)

- Each node must be visited by every agent infinitely often.
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- Initial configurations are the configurations without multiplicities.

#### First studied family: Partial grids with NSEW

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#### Mobility tree and parameter q





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# Known results (exclusive perpetual explo.) 1/2

#### In the partial grids (FSYNC)

[Baldoni, Bonnet, Milani, Raynal. IPL 2008]

- Definition of the parameter q
- Solvable iff  $k \leq n q$

[Baldoni, Bonnet, Milani, Raynal. OPODIS 2008]

- Limited visibility  $\varphi = 1$
- Solvable iff  $k \leq n q$ , except for  $q = 0 \Longrightarrow n 1$

#### In the (complete) grids

[Bonnet, Milani, Potop-Butucaru, Tixeuil. OPODIS 2011]

- ASYNC, no sense of direction
- Deterministically:  $\rho^-(n) = 3$

# Algorithmic techniques (exclusive perpetual explo.)

#### Two phases

- Form a special asymmetric pattern
- Move the pattern through the graph



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# Known results (exclusive perpetual explo.) 2/2

Case of the rings

### [Blin, Milani, Potop-Butucaru, Tixeuil. DISC 2010]

- FSYNC det.: Impossible when k is even
- ASYNC det.:  $\rho^-(n) = 3$  for  $n \ge 10$  and n and k coprime

• ASYNC det.:  $\rho^+(n) = n - 5$  for k odd & n and k coprime [Bonnet, Défago, Petit, Potop-Butucaru, Tixeuil. SRDS Workshop 2014]

• Generic method to list all protocols (SSYNC)

• Application: impossible for k = 5 and n = 10

[D'Angelo, Di Stefano, Navarra, Nisse, Suchan. Algorithmica 2015]

- ASYNC but only for rigid initial configurations
- Det. algo for any k, n such that  $n \ge 10$  and

$$5 \le k \le n-3$$
, except the case  $k = 5$  and  $n = 10$ 

# Proofs of correctness (1/2)

Case-by-case analysis



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# Proofs of correctness (2/2)

DAG with potential functions



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### Formal verification via model checking

[Bérard, Lafourcade, Millet, Potop-Butucaru, Thierry-Mieg, Tixeuil. Distributed Computing 2016]

Study of two algorithms

- Exploration with stop [FIPS Algorithmica 2013]
  - Correct for  $17 \le k \le 21$  and  $k < n \le 22$
  - Also correct for smaller values of k (5, 7 and 10)
- Exclusive perpetual exploration [BMPT DISC 2010]

• Counter-example for k = 3 and n = 10

### [Doan, Bonnet, Ogata. SOFG+MSVL 2016]

- Also study of [BMPT DISC 2010]
- Same counter-example, with a different technique

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### The counter-example



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### Conclusion and perspectives

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- Computability depends on the amount of symmetries
- Tricky to (correctly) handle symmetries and asynchrony

#### Interesting directions of research

- Limited visibility
- Formal verification
  - (Parametrized) Model checking
  - Proof assistant

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# Thank you for your attention

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