centre de recherche en informatique de lens	EXPERIMENTING WITH SMALL CHANGES IN CONFLICT-DRIVEN CLAUSE-LEARNING ALGORITHMS			
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-Conflict-Driven Clause Learning: -are believed to be well-known:

- * They embed dynamic heuristics [6, 3], learning [7], restarts [4, 1] and lazy data structures [6].
- * (efficient) SAT Solvers can be written from scratch in less than a thousand lines of code.
- -are not so well-known:
- * Who really understand the underlying mechanisms?
- * No real experimental studies: Progresses have been made with extensive tests "only"
- * What is a good learnt clause? When to forget? When to restart?



LISA SYNDROME REVISITED

-Principle

- Shuffling instances is admitted to be "bad" for industrial benchmarks, Is it really true?? How much it is bad?
- -We shuffled instances [5] and performed 50 runs by instance

-Observations

- -Curves "best" (resp. "besti") plot the result of virtual solverswhich would have the best CPU time on all shuffled runs (resp. the *ith* percentile)
- -MINISAT solves 55 problems, "best" 69.



-Crucial for future progresses:

- In-depth study rather than immediate improvements – Understanding each component of any CDCL solvers

-Questions adressed here:

- How small changes in Minisat [2] change results?
- Impact of data structures v.s. "good" ideas

-25% of chances to obtain better results by shuffling instances!

-Conclusion

- Shuffling instance may improve results. Many shuffling shows that MINISAT on original instance is not the best run. Improvements may be obtained here.
- Shuffling is not so bad. If it would have been so crucial to keep original benchmarks, we would have expected worst results.



-Principle

- -Solvers need parameters. How to select them? Is changing parameters as strong as changing solvers?
- -We took 10 magic values of MINISAT and built 126 solvers, each 1-parameter away from the original one.

-Observations

- -Released MINISAT is well-tuned. Taking 2 versions improves performances, but more than 3 don't. -Satelite is not always important (best of 2 contains a core version with fast restarts).
- Is MINISAT hard limit 74 or 85? What performances are awaited for a "better" solver than MINISAT? -Conclusion
- -High discrepency is observed when changing restart policies. Fast restart is efficient with MINISAT core.
- -Small changes can really change solvers. Ideas for multicore processors?

-Principle

-When new solvers are "better" than older ones? New ideas may hide noise effects on parameters! -Side-effect over parameters are simulated by adding only 10% noise each time a constant is requested.

-Observations

- Noisy MINISAT behaves like a real different solver. Hard to say if one is better than another.

- Conclusion

-When observing this kind of plot, nothing must be drawn from it, relatively to new ideas. This figure may only report a noisy side-effect of any data-structure.



CONCLUSION

-Experimenting?

- -Experimentation is the traditional companion section of papers. Competition are organized [5] to "rank" solvers.
- This is more "testing" than "experimenting".

- Conclusion

- -Large clauses are important for UNSAT proofs (already known).
- -Large clauses may have a local but important impact on the proof, when it is reduced in the search tree.
- The simple assumption that shorter clauses are better is false. We need a way to identify good clauses, and experimentally validate this.

-Poster Lessons

- -A lot of work is still pending to really understand the beahvior of CDCL solvers.
- Not easy to prove that a solver is "better" than another. What is "better"? Hard question, even on a fixed set of benchmarks.
- We need to build a real experimental science of CDLC solvers. This is possible, the SAT area is one of the pioneers in this direction.

References

[1] A. Biere. Adaptive restart strategies for conflict driven SAT solvers. SAT'08. [2] N. Een and N. Sorensson. An extensible sat-solver. SAT'03. [3] E. Goldberg and Y. Novikov. BerkMin: A fast and robust SAT-solver. DAM'07. [4] J. Huang. The effect of restarts on the efficiency of clause learning. *IJCAI'07*. [5] D. Le Berre and L. Simon. Essentials of the SAT'03 competition. SAT'03. [6] M. Moskewicz, C. Madigan, Y. Zhao, L. Zhang, and S. Malik. Chaff: Engineering an efficient SAT solver. DAC'01. [7] J. Marques-Silva and K. Sakallah. Grasp - a new search algorithm for satisfiability. ICCAD'96.