

Towards Wide-Coverage Semantics for French

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Bridge between statistical NLP and syntax/ semantics the way I (and many people here) like it!

Pon't worry, this will not be a talk of the style I improved on task X from Y% to Y+. 2%

There will be some percentages, but just to show we are up to the level of some of the statistical NLP guys.

- Many wide-coverage parsers for French exist (witness the participation of the Easy and Passage campaigns)
- My goal is not directly to compete with them, but to move towards a wide-coverage parser which produces structures which are more interesting (at least to me!) than shared forests

- I will talk about my current research on a wide-coverage categorial grammar for French.
- I will start by giving a short introduction to categorial grammars, showing how a categorial parse corresponds to a lambda-term in the simply typed lambda calculus.
- Since the work of Montague, we know that the simply typed lambda calculus forms a solid base for the semantic analysis of fragments of natural language.

Mowever, we are by no means limited to Montague semantics: Muskens (1994) and de Groote (2006) show that the semantics of categorial grammars are compatible with modern theories of dynamic semantics (DRT in the case of Muskens, and a continuation-based approach in the case of de Groote)

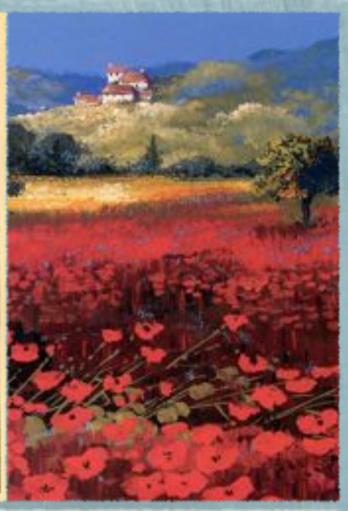
- In this talk I will present the Grail parser and the development of a wide-coverage grammar of French as well as the development of a prototype semantic lexicon producing DRSs, highlighting the treatment of presupposition.
- Wide-coverage semantics in this sense is a relatively new field, which was pioneered for English by Bos e.a. (2004)

Overview

- Categorial Grammars and Lambda Calculus Semantics
- Grammar Extraction
 - converting a corpus into categorial grammar
 - how to use this grammar for parsing
- Wide-Coverage Semantics







Lexicalized grammars and lambda calculus semantics

Formulas and corresponding expressions

- np
- n
- ₩ S
- mp\s
- mp/n
- $(np\s)/np$

- Jean, l'étudiant, ...
- étudiant, économie, ...
- Jean dort, Jean aime Marie
- dort, aime Marie
- w un, chaque, l'
- aime, étudie

Rules

Lambek categorial grammars have only four rules: an elimination and an introduction rule for both "\" and "/"

$$\begin{array}{c|cccc} A/B & B & B \setminus A \\ \hline A & & A \\ \end{array} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & B \setminus A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & A \\ & & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A & A \\ & & A \\ \end{bmatrix} \hspace{-0$$

$$[B]^{i} \qquad [B]^{i} \qquad \dots$$

$$\vdots \qquad \vdots \qquad \vdots$$

$$\frac{A}{A/B} [/I]^{i} \qquad \frac{A}{B \setminus A} [\backslash I]^{i}$$

Example

un étudiant dort np/n n np\s

$$\begin{array}{c|cccc} A/B & B & B \setminus A \\ \hline A & & A \\ \end{array} \hspace{-0.5cm} \begin{bmatrix} A & B \\ & A \\ \end{bmatrix} \hspace{-0.5cm} \begin{bmatrix} A$$

Example

$$\begin{array}{c|ccc} un & \text{\'etudiant} & \text{dort} \\ \hline np/n & n & np \backslash s \\ \hline & np & & \end{array}$$

Example

$$\frac{np/n}{\frac{np/n}{n}} \frac{n}{[/E]}$$

$$\frac{np}{np} \frac{np}{s}$$

$$S$$

Lamda calculus Beta reduction

$$(\lambda x.t) u = t[x:=u]$$

that is t, but with all occurrences of the variable x replaced by the term u (renaming variables when necessary)

- $(\lambda x. x+x) 2 = 2+2$
- $(\lambda y.(y x)) (\lambda z. (f z)) = (\lambda z. (f z)) x = f x$
- $(\lambda y.(y x)) (\lambda z. (f z)) = (\lambda z. (f z)) x = f x$

Lamda calculus Types

Inductive Definition

- Basic types e (for entity) and t (for truth value)
- ♠ If α and β are types, then α→β is a type

Lamda calculus Terms

- For each type α , there is a (countably infinite) number of variables x, y, z, ... which are terms of type α .
- For each type α , there is a (countably infinite) number of constants a, b, c, ... which are terms of type α .
- If x is term of type $\alpha \rightarrow \beta$ and y is a term of type α then (x y) is a term of type β .
- If y is a term of type β and x is a variable of type α , then $\lambda x.y$ is a term of type $\alpha \rightarrow \beta$

Lamda calculus Terms

Some useful constants

$\wedge_{\prime}\vee_{\prime}\Longrightarrow$	$t \rightarrow (t \rightarrow t)$
A'A	(e→t)→t
ı	(e→t)→e

Lamda calculus Notational conventions

- These notational conventions are useful for people familiar with predicate logic and Montague grammar.
- We will write $(x \land y)$ instead of $((\land y) x)$ and adopt a similar convention for the other boolean connectives
- We will write ιx .t instead of $\iota(\lambda x.t)$ and adopt a similar convention for \forall and \exists
- We will write p(x,y) instead of ((p y) x) for constants p (in general $p(x_1,...,x_n)$ for $(...(p x_n)...x_1)$

Lamda calculus Notational conventions

- $\lambda z.\lambda y.\lambda x.(((f z) y) x) \equiv \lambda z.\lambda y.\lambda x.f(x,y,z)$
- $\forall (\lambda x. \exists (\lambda y. ((love y) x)) \equiv \forall x. \exists y. love(x, y)$
- $\forall \forall (\lambda x.(\rightarrow (\text{man } x)) \text{ sleep } x)) \equiv \forall x.\text{man}(x) \rightarrow \text{sleep}(x)$

Lamda calculus Formulas as Types

This is the sassigning ty in categoria

In many of the examples, I will assign e to both np and vp, I will leave the changes required to turn these into terms of type (e->t)->t as an exercise

- ♠ Exception 1: np is lifted from e to (e→t)→t
- Exception 2: pp is assigned the same type as np

Formula	Type
type(np)	(e→t)→t
type(pp)	(e→t)→t
type(n)	e→t
type(s)	t
type(B/A)	$type(A) \rightarrow type(B)$
type(A\B)	$type(A) \rightarrow type(B)$

Lamda calculus Formulas as Types

♠ Exception 1: from e to (e=

In many of the examples, I will assign e to both np and vp, I will leave the changes required to turn these into terms of type (e->t)->t as an exercise

- Exception 2:the same type as np
- To simplify the examples which follow, I will use e as the type of np and pp in the examples (exercise: give the correct lambda terms)

Formula	Type
type(np)	e
type(pp)	e
type(n)	e→t
type(s)	t
type(B/A)	$type(A) \rightarrow type(B)$
type(A\B)	$type(A) \rightarrow type(B)$

- Proofs in categorial grammar correspond to lambda terms
- These lambda terms abstract away from the directions of the implications.

$$\frac{\text{t:A/B} \quad \text{u:B}}{\text{(t u):A}} \quad \frac{\text{u:B} \quad \text{t:B} \setminus \text{A}}{\text{(t u):A}}$$

Type lifting is now a simple consequence of the logical rules.

$$\frac{\text{t:A/B} \quad \text{u:B}}{\text{(t u):A}} \quad \frac{\text{u:B} \quad \text{t:B} \setminus \text{A}}{\text{(t u):A}}$$

x:np

Type lifting is now a simple consequence of the logical rules.

$$\frac{\text{t:A/B} \quad \text{u:B}}{\text{(t u):A}} \quad \frac{\text{u:B} \quad \text{t:B} \setminus \text{A}}{\text{(t u):A}}$$

$$\frac{x:np \quad y:(np \setminus s)}{(y \ x):s}$$

Type lifting is now a simple consequence of the logical rules.

$$\frac{\text{t:A/B} \quad \text{u:B}}{\text{(t u):A}} \quad \frac{\text{u:B} \quad \text{t:B} \setminus \text{A}}{\text{(t u):A}}$$

$$\frac{x:np \quad [y:(np\s)]}{(y \ x):s}$$

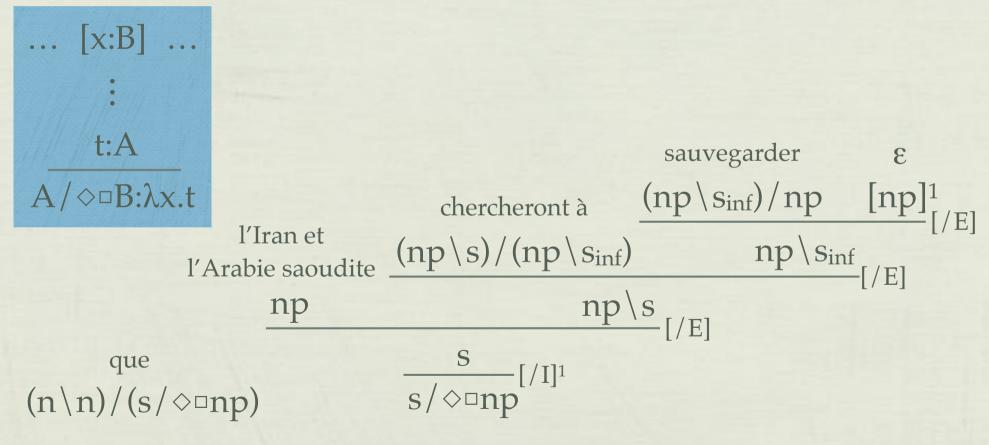
$$\frac{\lambda y.(y \ x):s/(np\s)}{(y \ x):s/(np\s)}$$

```
[x:B]
                                                                              sauvegarder
A/ΦΒ:λx.t
                                            \frac{(np \ s_{inf})/np \quad [np]}{(np \ s)/(np \ s_{inf})} \frac{(np \ s_{inf})/np \quad [np]}{(np \ s_{inf})} [/E]
                         l'Iran et
                   l'Arabie saoudite
                             np
          que
(n \ n)/(s/\diamond np)
                           n n
```

que
$$(n \ n)/(s/\diamond \square np)$$

que
$$(n \ n)/(s/\diamond \square np)$$

```
[x:B]
                                                                                sauvegarder
                                             \frac{(np \ s_{inf})/np \quad np}{(np \ s)/(np \ s_{inf})} \frac{(np \ s_{inf})/np \quad np}{(np \ s_{inf})}
A/ΦΒ:λx.t
                          l'Iran et
                    l'Arabie saoudite.
                                                                    np \setminus s [/E]
                              np
                                                      S
          que
(n \ n)/(s/\diamond np)
```



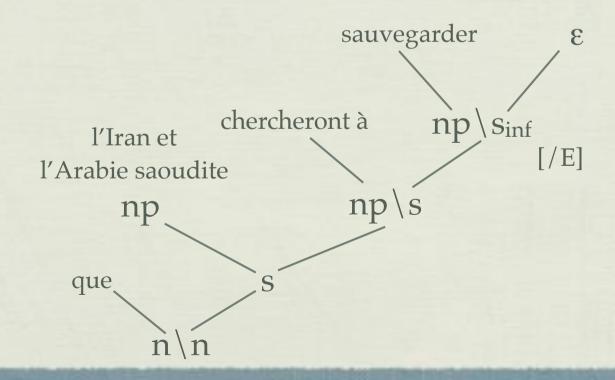
```
... [x:B]
                                                                             sauvegarder
A/ΦΒ:λx.t
                                            \frac{(np \ s_{inf})/np \quad [np]^1}{(np \ s)/(np \ s_{inf})} \frac{(np \ s_{inf})/np \quad [np]^1}{(np \ s_{inf})}
                         l'Iran et
                   l'Arabie saoudite
                             np
          que
(n \ n)/(s/\diamond np)
                           n \ n
```

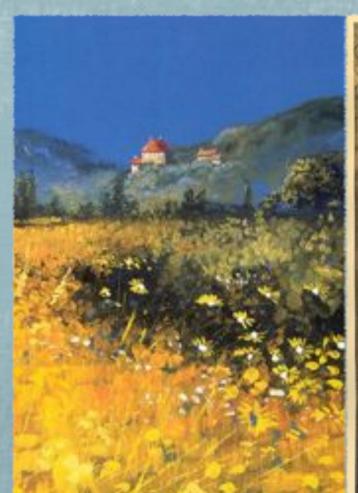
```
\lambda x. \acute{e}quilibre(x) \wedge d\acute{e}licat(x) \wedge
                               chercheront_à_sauvegarder(I&As,x)
                                                                                    sauvegarder
                                               \frac{(np \ s_{inf})/np \quad [np]^1}{(np \ s)/(np \ s_{inf})} \frac{(np \ s_{inf})/np \quad [np]^1}{(np \ s_{inf})}
A/\Diamond \Box B:\lambda x.t
                           l'Iran et
                    l'Arabie saoudite
                               np
          que
(n \ n)/(s/\lozenge np)
```

... cet équilibre délicate que l'Iran et l'Arabie saoudite chercheront à sauvegarder ε à Genève.

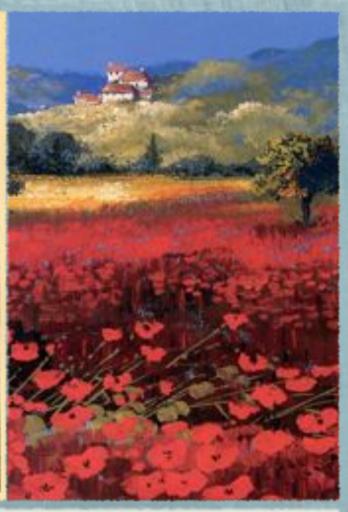


 $\lambda x. \acute{e}quilibre(x) \wedge d\acute{e}licat(x) \wedge$ chercheront_à_sauvegarder(I&As,x)









Grammar Extraction

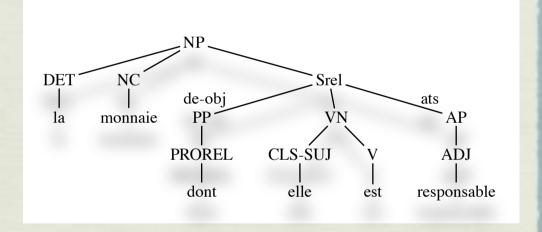
From the Paris VII corpus to a categorial lexicon, while developing several taggers

Grammar Extraction

Grammar extraction is the conversion of a linguistically annotated corpus (in our case, the Paris VII treebank) into a grammar into a grammar formalism the people doing the conversion really like (in our case, categorial grammar)

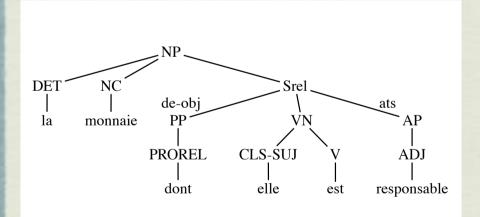
The Paris VII Corpus

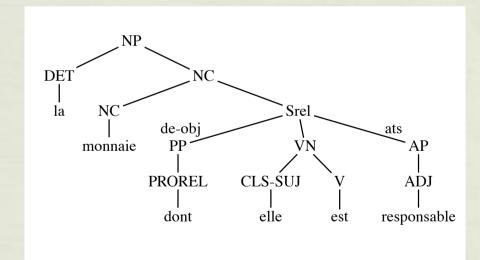
To the right is a small sentence fragment of the Paris VII corpus, which suffices to illustrate the extraction procedure



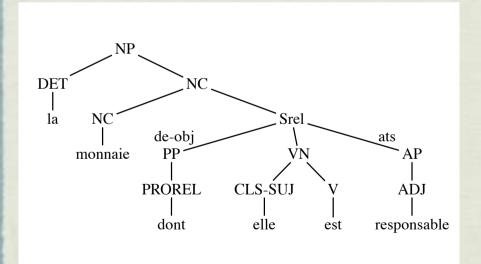
The extraction algorithm

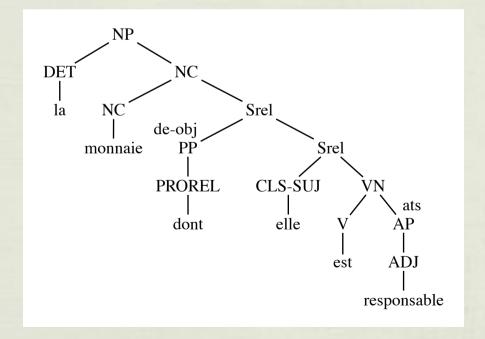
1. Binarize the annotation



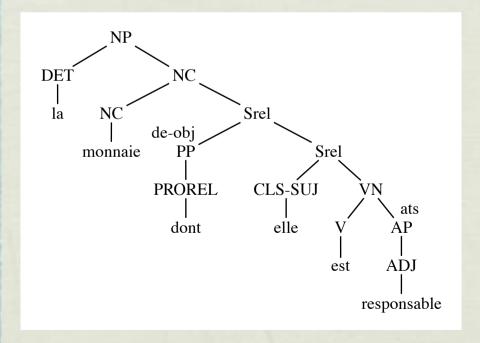


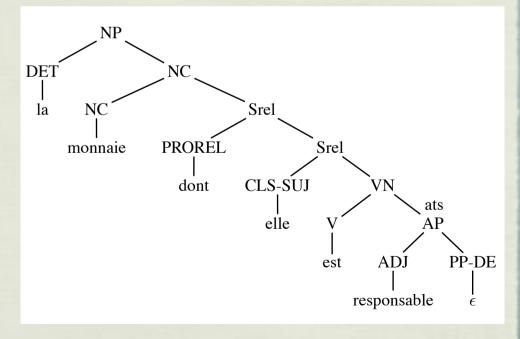
1. Binarize the annotation

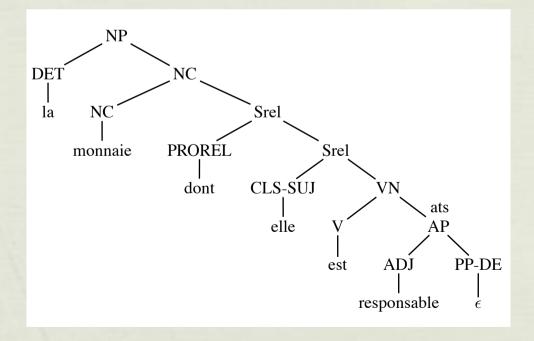


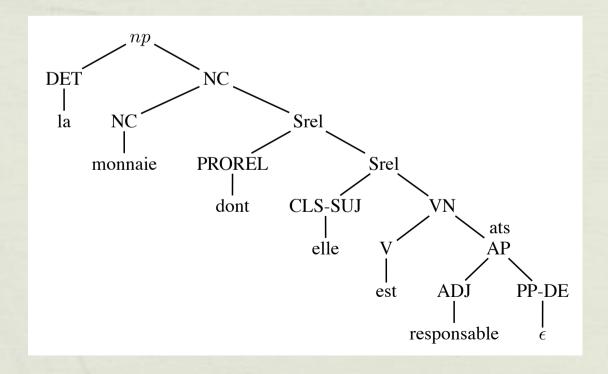


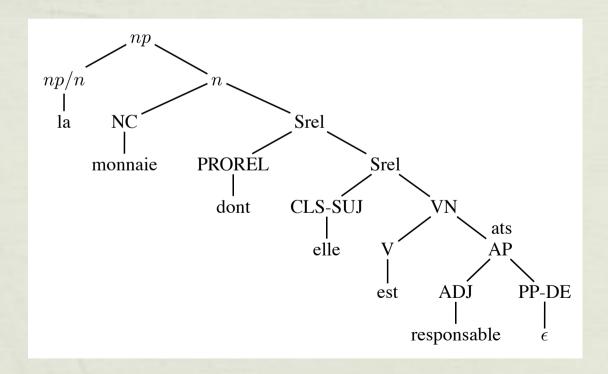
1. Binarize the annotation inserting traces for wh words

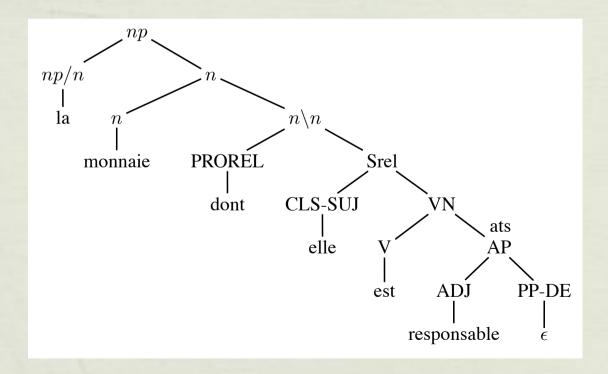


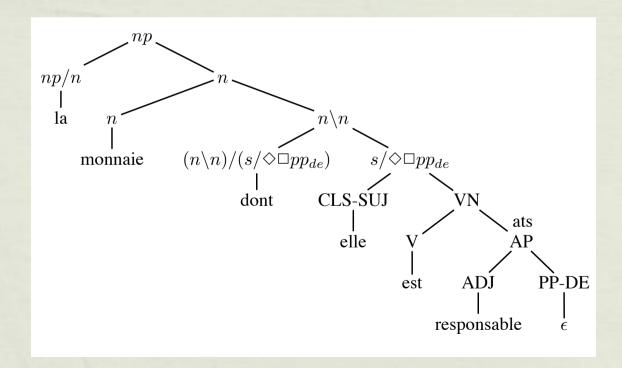


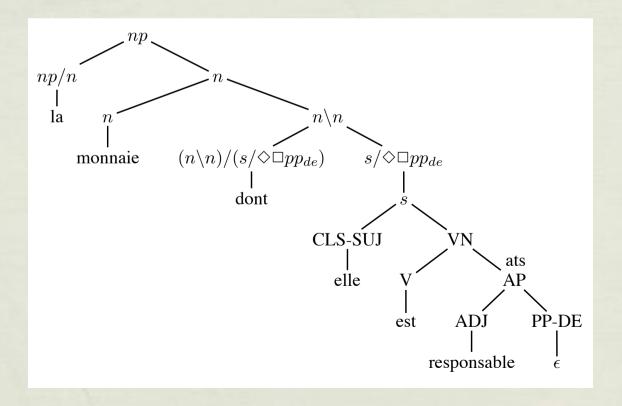


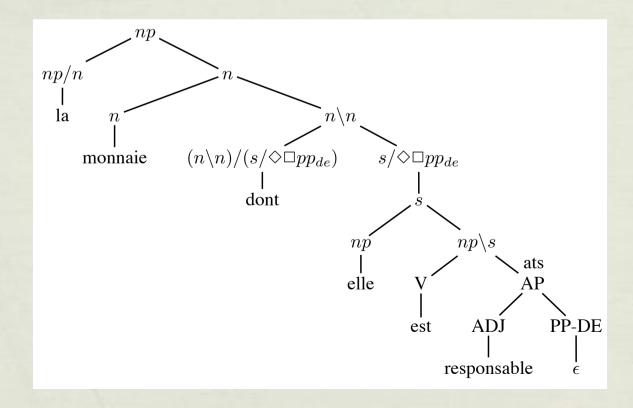


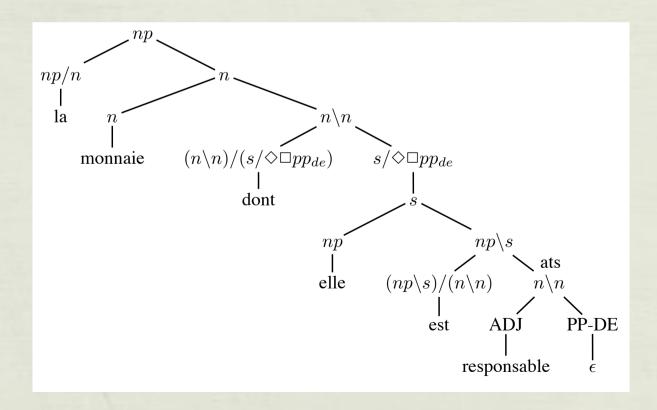


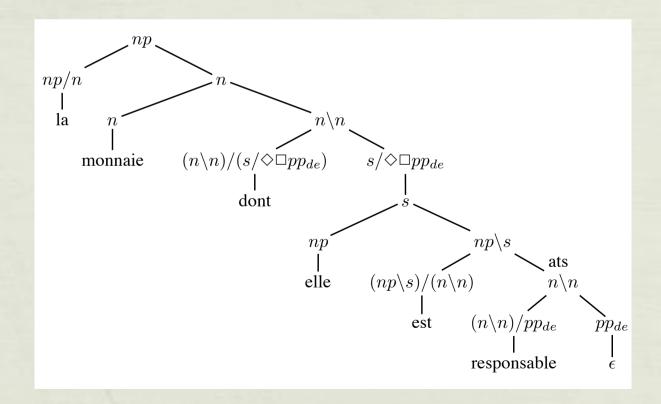












Grammar Extraction

- A lot of useful information (such as the position of "traces" of extracted elements) is not annotated but very useful for the grammar and needs to be added by hand.
- In addition, the extracted grammar has received a very significant amount of manual cleanup

- On the basis of the 382.145 words and 12.822 sentence of the treebank, the extraction algorithm extracts 883 different formulas, of which 664 occur more than once.
- Many frequent words are assigned <u>many</u> different formulas
- This is a significant bottleneck for parsing, however we can circumvent this problem using well-known NLP techniques (skip, skip to example)

Word	POS	#
et	conj	71
,	ponct	62
à	prp	55
plus	adv	44
ou	conj	42
est	verb	39
être	inf	36
en	prp	34
a	verb	31

POS	#
adv	206
conj	92
prp	149
ponct	89
verb	175

An illustration of some of the most ambiguous words and part-of-speech tags.

POS tagsets

MElt (Denis & Sagot) and Treetagger (Schmid) are the two main POS-taggers for French. They differ slightly in their tagsets: eg. Treetagger has a single tag "PRP:det" for "du" and "des" (harder for the supertagger), whereas the MElt tagger distinguishes between these words as determiner and as preposition (harder for the POS-tagger)

Word	MElt	Tt
Numerals	DET,ADJ,NC,NPP	NUM
du/des	DET, P+D	PRP:det
Verbs	V	VER:{simp,impf,}

- Formula assignments to the present tense form "fait"
- 124 occurrences in the corpus, with 19 different formulas assigned to it.

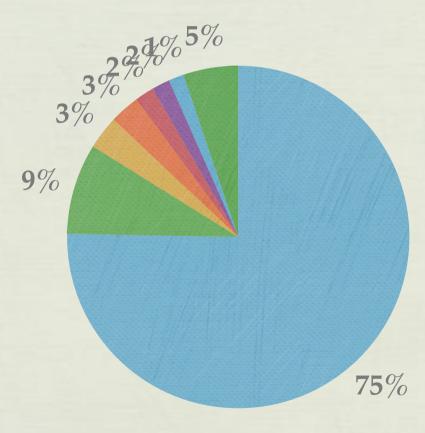
(np\s)/np
((np\s)/pp_de)/np)
(np\s)/(np/s_inf)
((np\s)/pp_a)/np
((np\s)/np)/(np\s_inf)
other

34

21

16

- Formula assignments to the comma ","
- 21,398 occurrences, 62 different formulas.
 - no formula
 - $(np \np)/np$
 - (n n)/n
 - $(np \np)/n$
 - $(s \setminus s)/s$
 - $((np\s)\(np\s))/(np\s))$
 - $((n \ n) \ (n \ n)) \ (n \ n))$
 - other



- The sum up, we have produced a categorial grammar for French, which is essentially a very big lexicon.
- The size of this lexicon, coupled with high lexical ambiguity, makes direct exploitation for parsing difficult.
- A fairly standard solution is to use a supertagger to estimate the most likely sequence of formulas for the given words.

- Supertagging is essentially part-of-speech tagging but with richer structure hence "super" tags.
- Like part-of-speech tagging,
 we use superficial
 contextual information and
 statistical estimation to
 decide the most likely tag.



- So what is the context for a supertagger?
- Typically, it consists of the current word, the surrounding words, the current and surrounding
 POS tags and the previous supertags.

Context for "de"

np/n	n	?		
DET	NC	P	NPP	NPP
la	voiture	de	Prince	Charles

- The basic procedure for finding the sequence of formulas then becomes
 - Find the correct POS tag sequence
 - Find the correct supertag sequence

Context for "de"

np/n	n	?		
DET	NC	P	NPP	NPP
la	voiture	de	Prince	Charles

- Estimation is done using maximum entropy models
- Wery standard and easy to modify (ie. we can add <u>any</u> information we think is useful and let the estimation algorithm decide which ones really are).

Context for "de"

np/n	n	?		
DET	NC	P	NPP	NPP
la	voiture	de	Prince	Charles

Any information which we can easily obtain, of course. If we think a word having an even number of letters is useful, we can add it.

Good performance and

POS/Supertagging

Note, that, though Partof-Speech tagging
helps, an incorrect POStag can actually hurt
the supertagger.

np/n	n	(np\s)/np	np/n	n
DET	NC	V	DET	N
la	petite	brise	la	glace

Errors in DET-N versusCLO-V POS-tags aredifficult for thesupertagger to recoverfrom.

np/n	n/n	n	(np\s)/((np\s)/np)	(np\s)/np
DET	ADJ	NC	CLO	V
la	petite	brise	la	glace

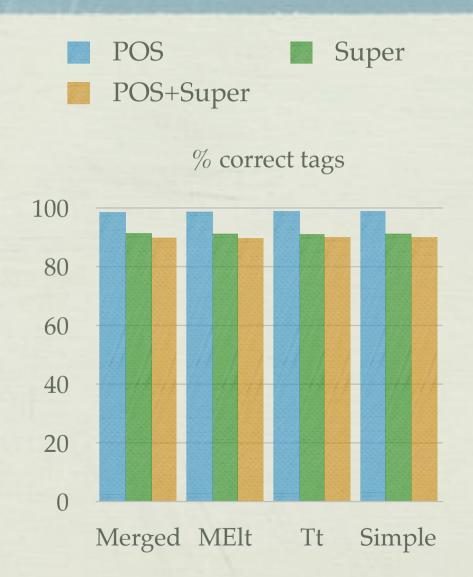
POS/Supertagging

Other difficult words for the POS-tagger include "que" (which can be a conjunction, an adverb or a relative pronoun)

np/n	n	(np\s)/np	np/n	n
DET	NC	V	DET	N
la	petite	brise	la	glace

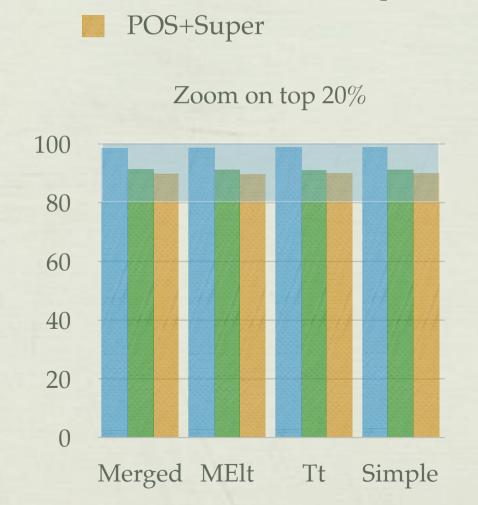
However, in general,the POS-taginformation helps (as we will see)

np/n	n/n	n	(np\s)/((np\s)/np)	(np\s)/np
DET	ADJ	NC	CLO	V
la	petite	brise	la	glace



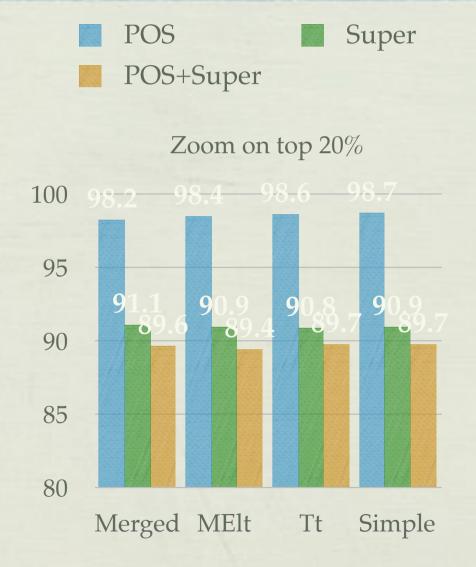
- A plot of POS/
 Supertagger results for the four different tagsets.
- POS+Super gives the % correct supertags given the POS-tag assigned by the tagger, Super is the correct supertag given

Super



POS

- A plot of POS/
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- A plot of POS/
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- POS+Super gives the % correct supertags given the POS-tag assigned by the model, Super is the correct supertag given

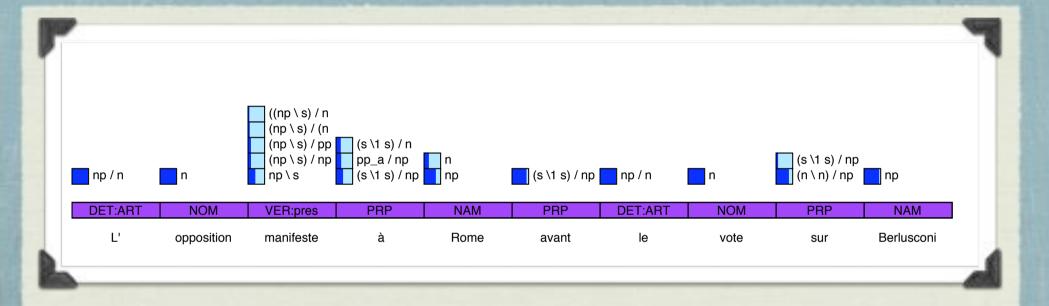
Multiple Solutions

- Though these results are comparable to the best supertaggers for English, in practice, even at around 91% correct supertags, we do not cover enough sentences of the corpus.
- A standard solution is to look at supertags within a range depending on the best supertag.
- This is called the β value.

Multiple Solutions

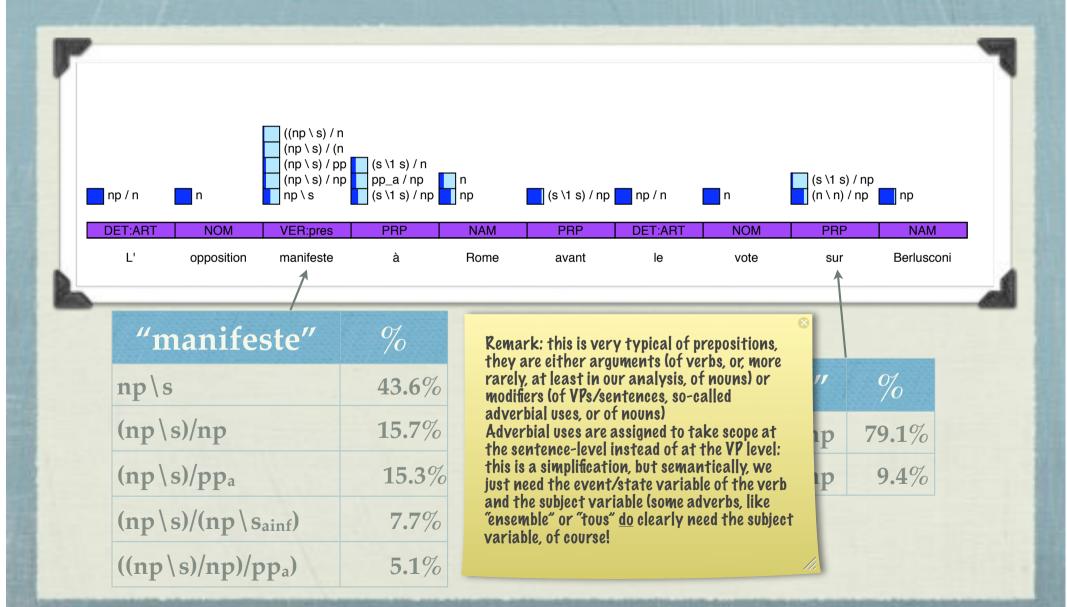
- Roughly speaking: if p is the probability of the best supertag, we will assign all supertags of probability > βp
- So, the less we are sure of our first supertag, the more alternatives we add.
- On average, a β of 0.1 gives 2.7 supertags per word,0.05 gives 3.1 and 0.01 gives 4.7

Example



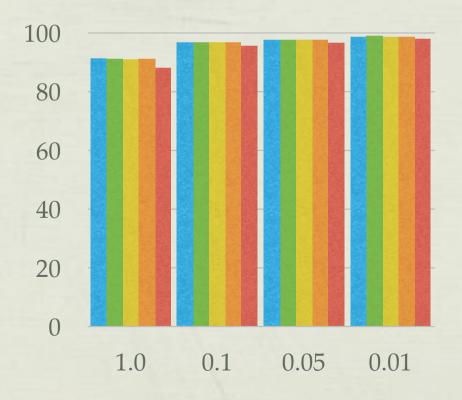
- Mere is an example with β=0.1
- We can see that many "easy" words get assigned a single supertag whereas difficult words (here: verbs and prepositions) get assigned many tags.

Example





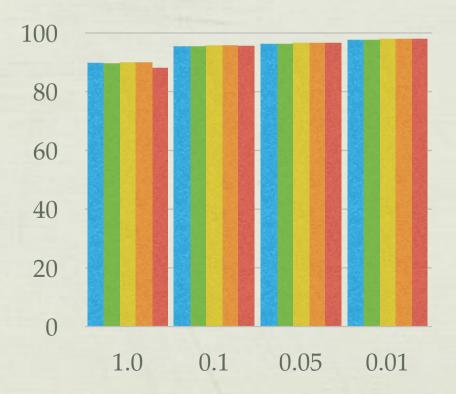
% correct supertags by model and β value



- Results with the use of different values of β.
- In a sense, the β value allows us to trade coverage for efficiency: at higher values of β , we parse more sentences, but we do so more slowly.



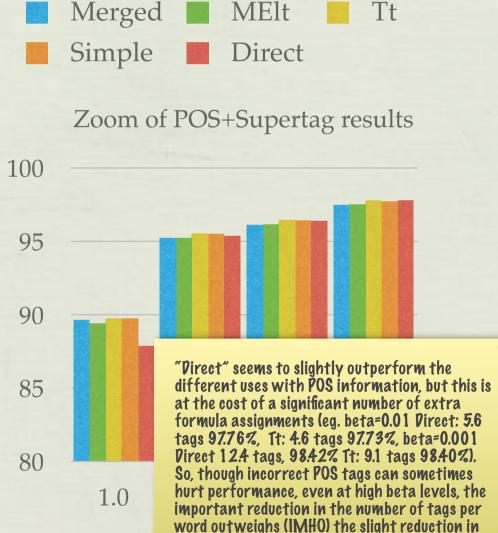
% correct supertags by model and β value



- As before, there is a slight decrease in performance once we switch from "gold" POS tag to tags assigned by the tagger.
- Eg. for the Treetagger tagset, it is -1.0% at β =0.1 and -0.5% at



- A comparison of the Supertagger and the combined POS/ Supertagger.
- Same results as the previous slides, but with a zoom on the top 20 percentile.
- Direct is the result of the

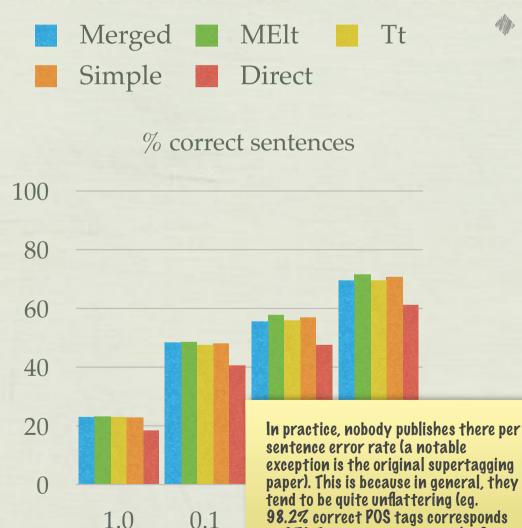


correct tags.

 A comparison of the Supertagger and the combined POS/ Supertagger.

Same results as the previous slides, but with a zoom on the top
20 percentile.

Direct is the result of the



to 65.1% correct sentences, the figures

for beta=0.01 indicate a similar

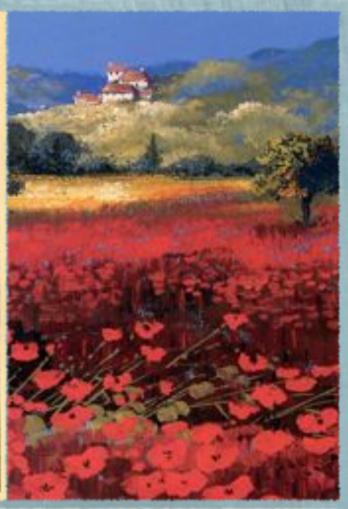
picture)

Finally, here is the percentage of sentences which are assigned the correct sequence of supertags for the different settings of β and the different POS models.

Note that we number of entences for which a arse is found is actually etter (around 85% at







On the development a wide-coverage semantic lexicon for the extracted categorial grammar

- As we have seen, formulas in categorial grammars correspond to types in the simply typed lambda calculus
- Proofs (parses) correspond to lambda terms.
- Thanks to the extracted grammar, we can obtain reasonable accurate parses for French sentences.
- So what is missing to obtain a Montague-style meaning of analysed sentences is a large enough lexicon.

- In order to move beyond a simple lexicon listing a limited number of words, it suffices to remark that many of the "open class" words (eg. names, nouns, verbs) follow a general schema to obtain their lexical semantics.
- For example, a noun "n" generally has λx.n(x) as its semantics.

- So the basic idea behind wide-coverage semantics is very simple:
 - the lexicon lists words which require special treatment (eg. conjunctions "et" and auxiliary verbs like "être" and "avoir")
 - other words are assigned a lambda term based on

So the general motto is: if you want to add more information to the semantic lexicon, there are two basic (non-exclusive) solutions: 1) you list the different cases 2) you train a (reliable) tagger Solution 1 would be an option for distinguishing subjet/object control verbs and Solution 2 would be an option for Named Entities (and their types: persons, places, enterprises), for more complicated semantic distinctions, like events versus states the solution is less clear.

Beyond Montague

- Montague grammar has some well-know limitations.
- I will talk briefly about two of them: anaphora and presuppositions
- I will sketch solutions to both of them, which stay within the framework of the simply typed lambda calculus.

Anaphora

A man walks in. He orders a beer.

- The meaning of this (tiny) discourse is clear: there exists a man, who enters and (presumably) this same man orders a beer.
- In first order logic, this is represented by the following formula:
- $\exists x. [enter(x) \land \exists y [beer(y) \land order(x,y)]]$

Anaphora

A man walks in. He orders a beer.

- \blacksquare $\exists x. [enter(x) \land \exists y [beer(y) \land order(x,y)]]$
- Now what are the formulas we can assign to the two sentences?
- \blacksquare $\exists x. [enter(x)]$
- \blacksquare $\exists y [beer(y) \land order(x,y)]$

** Kamp (1981) and Kamp & Reyle (1993) propose the following solution:

x

enter(x)

y,z

beer(y)

order(z,y)

z = ?

 $\exists x. [enter(x)] \exists y [beer(y) \land order(z,y)]$

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 $\exists x. [enter(x) \land \exists y [beer(y) \land order(x,y)]]$

- ** Kamp (1981) and Kamp & Reyle (1993) propose the following solution.
- This solution is compatible with the lambda calculus approach with have been adopting (Muskens 1994)

Example entries

marché : $n - \lambda x_0$. marché (x_0)

"dormir" is a state rather than an event, however, the current system does not distinguish between different types of eventualities.

Marie: $np - \lambda P_0.(y_0 \\ \mathbf{nomm\acute{e}}(y_0, \mathbf{mary}) \oplus P_0(y_0))$

chaque : $np/_0n - \lambda P_0Q_0.((z_0 \oplus P_0(z_0)) \to Q_0(z_0))$

Example entries

 "dormir" is a state rather than an event, however, the current system does not distinguish between different types of eventualities.

pousser: $((np \setminus_0 s)/_0 (np \setminus_0 s_{ainf}))/_0 np - \lambda x_0 y_0 z_0 x_1.x_0 (\lambda y_1.z_0 (\lambda z_1.$

 d_2

 $egin{aligned} \mathbf{pousser_\grave{a}}(x_1) \ \mathbf{agent}(x_1, z_1) \ \mathbf{patient}(x_1, y_1) \ \mathbf{theme}(x_1, y_2) \end{aligned}$

 $y_2:y_0(z_2,x_0)$

What is presupposition?

- Presupposition (like anaphora and their resolution) are a linguistic phenomenon on the semantics/pragmatics interface.
- Its particularity is the it *presupposes* something which may not be directly said.
- So we could say it is a sort of "window" through which we can observe aspects of the abstract notion of "common ground".

What is presupposition?

Some examples

- 1. Have you stopped beating your wife?
 - a. presupposes the listener has been beating his wife
- 2. George W. Bush would torture again.
 - b. presupposes Bush has tortured
- 3. Obama regrets intervening to save Terry Schiavo.
 - c. presupposes Obama intervened to save T. Schiavo

What is presupposition?

How can we decide if something is a presupposition (as opposed to an implicature or an entailment?

- 1. Presuppositions stay when embedded inside of a negation, a question, or a modal.
- 2. "Hey, wait a minute, I didn't know that..."
- Obama regrets intervening to save Terry Schiavo.
- Obama doesn't regret intervening to save Terry Schiavo.
- Obama may regret intervening to save Terry Schiavo.

What is presupposition?

How can we decide if something is a presupposition (as opposed to an implicature or an entailment?

- 1. Presuppositions stay when embedded inside of a negation, a question, or a modal.
- 2. "Hey, wait a minute, I didn't know that..."
- Hey, wait a minute, I didn't know he intervened to save Terry Schiavo.
- # Hey, wait a minute, I didn't know he regretted doing that.

From Karttunen (1973)

Question: what do the sentences presuppose about the guilt

of someone other than Nixon?

1. If Dean told the truth, Nixon is gu

2. If Haldeman is guilty, Nixon is gu

Just to show that things can get complicated and that inference is sometimes necessary: in (3), if destroying the tapes implies being guilty, then (3) as a whole doesn't imply someone besides Nixon is guilty, since this is part of the antecedent.

Note that with a lot of effort, we could do the same for (1): imagine a strange dictatorship, where it is against the law to speak the truth when it harms the president

3. If Miss Woods destroyed the missing tapes, Nixon is

guilty too.

Note that there is another possible lecture for "too", which would occur in cases like "Nixon was not just incompetent, he was quilty too".

- Presupposition is not an obscure phenomenon, which rarely occurs in a corpus: proper names and definite articles both presuppose existence of (at least an intension of) the name or definite description.
- So, with a quick and approximate calculation, we have an average of around three presuppositions per sentence!
- This means that in order to do wide-coverage semantics, we need to have at least some way of treating presupposition.

Presupposition in DRT

- DRT has been extended to handle presuppositions (Kamp 2001a, Kamp 2001b, Kamp & Reyle, to appear)
- The highly simplified version of it allows a DRS to be a pair of two DRSs, where the first contains the presuppositions of the second.

Presupposition in DRT

The man walks in. He orders a beer.

Presupposition in DRT

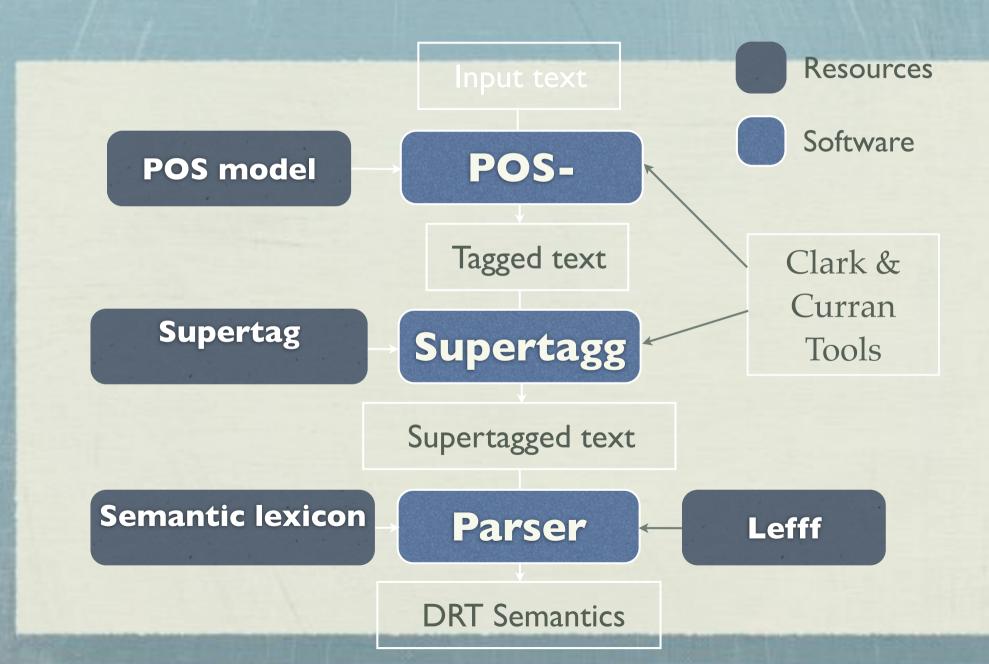
John walks in. He orders a beer.

name(x,John)

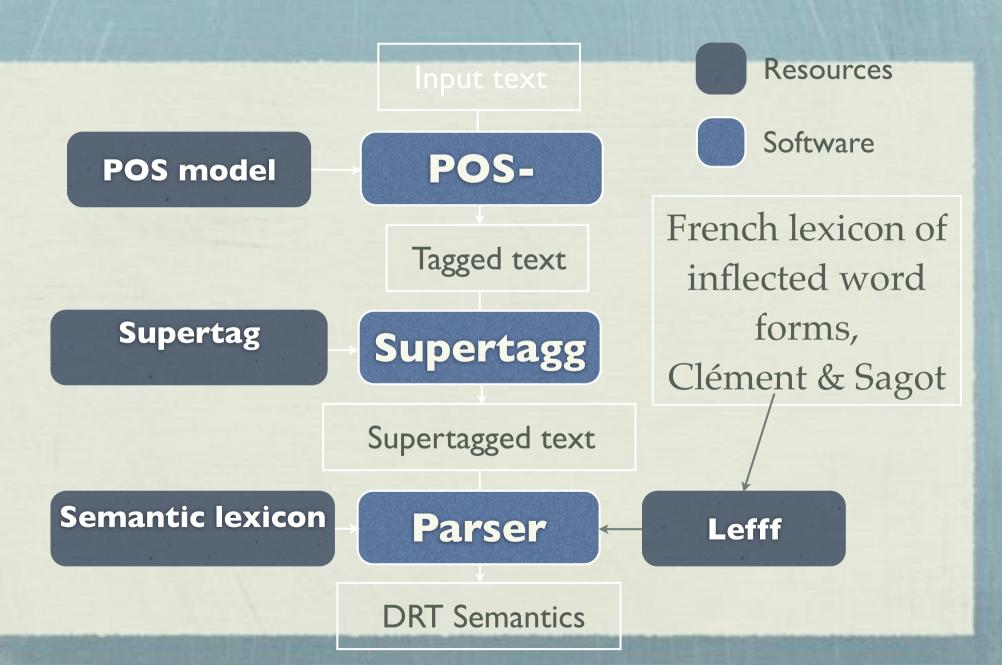
y,z

enter(x)
beer(y)
order(z,y)
z = ?

Grail & Friends



Grail & Friends



Demo

- All talk and no demo make Jack a dull boy.
- All talk and no demo make

Jac

Give a demo of the system with today's headlines from "Google Actualités"



Conclusion

- I have described the development of a wide-coverage categorial grammar for French and first steps towards using it for wide-coverage <u>semantics</u>
- All software and resources are available under LGPL (with the unfortunate exception of the annotated corpus, which is bound by the same conditions as the Paris VII treebank).

Future Work

A very long list, but I will mention some of the more important tasks.

Future Work - Parser

- Improve the accuracy of the extracted grammar and the parser
- Improve the efficiency of parser (eg. by using tree automata)

 (as in Noémie-Fleur's talk, of course!)
- Add a component for multi-word expressions.

Future Work - Semantics

- Incorporate a Named-Entity component.
- Incorporate a rudimentary analysis of tense/aspect and discourse structure using the French Timebank (Bittar 2010).
- Word sense disambiguation
- General problem: lack of annotated data

Future Work - Semantics

- Open questions:
 - how "deep" can we go with wide-coverage semantics?
 - what are appropriate evaluation measures?