Tutorial #5 - Inclusion polymorphism

Exercice 1: Inclusion polymorphism in Java

The Java type system follows a hierarchy that is described in the language specification (accessible at http://docs.oracle.com/javase/specs/jls/se8/html/jls-4.html) and that more or less corresponds to the following diagram :



- 1. Where are the existential types in this architecture? The universal types?
- 2. What are the subtyping relations between the primitive types? What are the relations between arrays and the other reference types?
- 3. Aside from the object type, give two examples of recursive types in the API (i.e types whose definition involves themselves somewhat). The Java classes are listed at https://docs.oracle.com/javase/8/docs/api/overview-tree.html.
- 4. What is the difference, from the point of view of subtyping, between the interface types and class types in Java?

The Scala programming language is a programming language developed at the EPFL (http://www.scala-lang.org), that can be compiled to Java bytecode and run on the Java virtual machine. A key property of the language is the possibility to mix Java and Scala .class files together. Its type system is described on this diagram¹.

5. What are the differences between the Scala types and the Java types?

^{1.} This diagram appears in chap. 11 of Programming in Scala, by Odersky et. al.

Exercice 2: Record polymorphism in OCaml

One of the specificities of the OCaml language is its support for objects (the "O" originally standing for "Objective"). The OCaml object layer shares a lot of similarities with the record types studied in the course. For reference, the documentation on objects is reachable at https://caml.inria.fr/pub/docs/manual-ocaml/objectexamples.html.

```
Objects can be defined independently from classes, as immediate objects :
    let o = object (self)
    val mutable x = 0
        (* Attribute *)
    method get = x
        (* Method *)
        (* Method *)
        (* det : int; move : int → unit >
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        (* det : int : int
```

The type of the object is inferred as a *record type*. Notice that the attributes do not appear in the type, and that the methods may have a constant type (such as get:int) even though they are not constant. The (self) part names the reference to the object inside the code.

Calling a method on an object is done the following way :

method move $d = x \leftarrow x + d$

end::

▶ A *class* is simply a constructor for an object :

```
class point = fun init → object (self)
  val mutable x = init
  method get = x
  method move d = x ← x + d
end;;
```

```
class point : int → object
  val mutable x : int
  method get_x : int
  method move : int → unit
end
```

Class definitions possess a nominal type (in this case point), whereas immediate objects possess a structural type, but in OCaml they are interchangeable. The creation of an object from a class is done the following way :

let o = new point 3;;

It is possible to *coerce* a value v into a supertype t with the expression v :> t.

Let us approximate the **Object** type to the following :

 $Object ::= {clone : Unit \rightarrow Object, equals : Object \rightarrow Bool, }$

1. Create the record type for **Object** in **OCaml**, and propose an implementation of the class. *Remark* : both constructions must be recursive.

Let us now consider a more utilitarian class :

- 2. Create a (recursive) record type olist to represent lists constructed with a head and a tail. The type should contain 3 methods : hd, t1, and is_empty. *Remark* : for the sake of this exercise, it is not necessary to create a polymorphic list.
- Write an immediate object that represents the empty list. Its non-trivial functions may return an error using failwith.
 Write a class that represents a non-empty list, with the usual head and tail elements passed to the constructor.
- 4. What would be an example of a subtype of olist?

Consider now writing an external function dealing with objects of type olist :

- 5. Write the length function for such lists, and analyze its type. Write this type, once with a universal variable and once with an existential variable.
- 6. Write the type for a dictionary containing pairs (key,value), that is compatible with the length function.

Exercice 3: Java selection method

Consider the selection algorithm for the Java 1.1 language that was described in the Java Language Reference by Mark Grand (1997)². Given an expression in Java containing a method call, this algorithm selects the appropriate method to apply when multiple choices are possible. It is described at http://web.deu.edu.tr/doc/oreily/java/langref/ch04_js.htm, corresponding in the book to the chapter 4, subsection « Method call expression ». The examples from this chapter are given in the sources of this tutorial.

^{2.} The older version of the Java language has a significatively simpler algorithm for method selection than the subsequent versions.

```
'/ First class hierarchy A \leftarrow B \leftarrow C \leftarrow D
class A {}
class B extends A {}
class C extends B {}
class D extends C {}
// Second class hierarchy W \leftarrow X \leftarrow Y \leftarrow Z
class W {
    void foo(D d) {System.out.println("W.D");}
class X extends W {
    void foo(A a) {System.out.println("X.A");}
    void foo(B b) {System.out.println("X.B");}
}
class Y extends X {
    void foo(B b) {System.out.println("Y.B");}
}
class Z extends Y {
    void foo(C c) {System.out.println("Z.C");}
}
// Main program
public class CallSelection {
    public static void main(String [] argv) {
         Z z = new Z();
         ((X) z).foo(new C());
    }
}
```

```
// Small class hierarchy A ← B
class A {}
class B extends A {}
// Examples of ambiguous calls
class AmbiguousCall {
   void foo(B b, double x){}
   void foo(A a, int i){}
   void doit() {
     foo(new A(), 8); // Matches foo(A, int)
     foo(new A(), 8.0); // Error: doesn't match anything
     foo(new B(), 8); // Matches foo(B, double)
   }
}
```

- 1. Read the article and describe the three (3) steps of the selection algorithm.
- 2. Determine in the article the two (2) potential sources of errors that can occur during the selection algorithm.
- 3. Determine which part of this algorithm deals with overloading (adhoc polymorphism) and which part deals with overriding (inclusion polymorphism).
- 4. Determine which part of this algorithm is done statically and which part is done dynamically (*Remark* : the answer is different from the previous question)