

Finding Short Synchronizing Words for Prefix Codes

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Introduction Prefix codes are a simple and powerful class of variable-length codes that are widely used in information compression and transmission. A famous example of prefix codes are Huffman’s codes [Huf52]. In general, variable length codes are not resistant to errors, since one deletion, insertion or change of a symbol can desynchronize the decoder causing incorrect decoding of the whole remaining part of the message. However, in a large class of codes called synchronizing codes resynchronization of the decoder is possible in such situations. It is known that almost all maximal finite prefix codes are synchronizing [FJTZ03]. Synchronization of finite prefix codes has been investigated a lot [Ba16, Bis08, BP09, CDSGV92, Sch64, Sch67], see also the book [BPR10] and references therein. For efficiency reasons it is important to use as short words resynchronizing the decoder as possible to decrease synchronization time. However, despite the interest to synchronizing prefix codes, the computational complexity of finding short synchronizing words for them has not been studied so far. We provide a systematic investigation of this topic.

Each recognizable (by a finite automaton) maximal prefix code can be represented by an automaton decoding the star of this code. For a finite code, this automaton can be exponentially smaller than the representation of the code by listing all its words (consider, for example, the code of all words of some fixed length). This can of course happen even if the code is synchronizing. In different applications the first or the second way of representing the code can be useful. In some cases large codes having a short description may be represented by a minimized decoder, while in other applications the code can be described by simply providing the list of all codewords. We study the complexity of problems for both arbitrary and literal decoders of finite prefix codes.

Huffman decoders There is a strong relation between partial automata and prefix codes [BPR10]. A set X of words is called a *prefix code* if no word in X is a prefix of another word. The class of recognizable (by an automaton) prefix codes can be described as follows. Take a strongly connected partial automaton A and pick a state r in it. Then the set of all *first return words* of r (that is, words mapping r to itself such that each non-empty prefix does not map r to itself) is a recognizable prefix code. Moreover, each recognizable prefix code can be obtained this way. A prefix code is called *maximal* if it is not a subset of another prefix code. The class of maximal recognizable prefix codes corresponds to the class of complete automata. If a state r can be picked in an automaton

in such a way that the set of all first return words is a finite prefix code, we call the automaton a *partial Huffman decoder*. If such automaton is complete (and thus the finite prefix code is maximal), we call it simply a *Huffman decoder*.

For the mentioned classes of decoders we obtain the following inapproximability result.

Theorem 1. *Unless $P = NP$, the problem of finding a shortest synchronizing words cannot be approximated in polynomial time within a factor of*

- (i) $n^{1-\varepsilon}$ for any $\varepsilon > 0$ for n -state binary strongly connected automata;
- (ii) $c \log n$ for some $c > 0$ for binary n -states Huffman decoders;
- (iii) $n^{\frac{1}{2}-\varepsilon}$ for any $\varepsilon > 0$ for binary n -state partial Huffman decoders.

We remark that for strongly connected automata the bound is optimal because there exists a $O(n)$ -approximation algorithm [GH11]. We also conjecture that for binary n -states Huffman decoders the $c \log n$ -inapproximability bound is optimal.

Literal decoders The literal automaton of a maximal finite prefix code X is defined as follows. The set of its states is the set of proper prefixes of the words in X and the transitions are naturally defined to concatenate letters to the prefixes (or to map to the empty prefix if the resulting word is in X). The number of states of a literal automaton is polynomially equivalent to the total length of all words in the corresponding finite prefix code. Thus, it is a natural model for the problems where the code is provided by simply enumerating all its codewords.

Theorem 2. *For literal n -state decoders there exist*

- (i) a polynomial time $O(\log n)$ -approximation algorithm
- and
- (ii) for any $\varepsilon > 0$ a quasi-polynomial time $(1 + \varepsilon)$ -approximation algorithm for the problem of finding a shortest synchronizing word.

We conjecture that there exists a polynomial time exact algorithm for this problem.

Mortal words A word is called *mortal* for a partial automaton A if its mapping is undefined for all the states of A . The techniques that we develop can be easily adapted to get the same inapproximability for the problem of finding a shortest mortal word. This problem is connected for instance to the famous Restivo's conjecture [Res81].

Theorem 3. *Unless $P = NP$, the problem of finding a shortest mortal word cannot be approximated in polynomial time within a factor of*

- (i) $n^{1-\varepsilon}$ for any $\varepsilon > 0$ for n -state binary strongly connected partial automata;
- (ii) $c \log n$ for some $c > 0$ for n -state binary partial Huffman decoders.

We also propose a simple polynomial time $O(\log n)$ -approximation algorithm for this problem in partial literal decoders.

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FIXED POINTS OF STURMIAN MORPHISMS AND THEIR DERIVATED WORDS

ABSTRACT

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1. INTRODUCTION

Sturmian words are probably the most studied object in combinatorics on words. They are aperiodic words over a binary alphabet having the least factor complexity possible, in other words, their factor complexity satisfies $\mathcal{C}_{\mathbf{u}}(n) = n + 1$ for each $n \in \mathbb{N}$. Many properties, characterizations and generalizations are known, see for instance [4, 3, 2].

One of their characterizations is in terms of return words to their factors. Let $\mathbf{u} = u_0 u_1 u_2 \cdots$ be a binary infinite word with $u_i \in \{0, 1\}$. Let $w = u_i u_{i+1} \cdots u_{i+n-1}$ be its factor. The integer i is called an *occurrence* of the factor w . A return word to a factor w is a word $u_i u_{i+1} \cdots u_{j-1}$ with i and j being two consecutive occurrences of w such that $i < j$. In [13], Vuillon showed that an infinite word \mathbf{u} is Sturmian if and only if each nonempty factor w has exactly two distinct return words. A straightforward consequence of this characterization is that if w is a prefix of \mathbf{u} , we may write

$$\mathbf{u} = r_{s_0} r_{s_1} r_{s_2} r_{s_3} \cdots$$

with $s_i \in \{0, 1\}$ and r_0 and r_1 being the two return words to w . The coding of these return words, the word $d_{\mathbf{u}}(w) = s_0 s_1 s_2 \cdots$ is called the *derived word of \mathbf{u} with respect to w* , introduced in [6]. A simple corollary of the characterization by return words and a result of [6] is that the derived word $d_{\mathbf{u}}(w)$ is also a Sturmian word. This simple corollary follows also from other results. For instance, it follows from [1], where the authors investigate the derived word of a standard Sturmian word and give its precise description. It also follows from the investigation of a more general setting in [5], which may in fact be used to describe derived words of any episturmian word — generalized Sturmian words [7].

By the main result of [6], if \mathbf{u} is a fixed point of a primitive morphism, the set of all derived words of \mathbf{u} is finite (the result also follows from [8]). In this case, again by [6], a derived word itself is a fixed point of a primitive morphism.

In this article we study derived words of fixed points of primitive Sturmian morphisms. By the results of [10], any primitive Sturmian morphism may be decomposed using elementary Sturmian morphisms — generators of the Sturmian monoid. We use the elementary Sturmian morphisms to describe the relation between the derived words of a Sturmian sequence. The main result of our article is an exact description of the morphisms fixing the derived words $d_{\mathbf{u}}(w)$ of \mathbf{u} , where \mathbf{u} is fixed by a Sturmian morphism ψ and w is its prefix. For this purpose, we introduce an operation Δ acting on the set of Sturmian morphisms with unique fixed point, see Definition 4. Iterating this operation we create the desired list of the morphisms as stated in Theorem 5. The Sturmian morphisms with two fixed points are treated separately, see Proposition 6.

We continue our study by counting the number of derived words, in particular by counting the distinct elements in the sequence $(\Delta^k(\psi))_{k \geq 1}$. This number depends on the decomposition of ψ into the generators of the special Sturmian monoid, see below in Section 2. Using this decomposition, Propositions 8 and 9 provide the exact number of derived words for two specific classes of Sturmian morphisms.

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For a general Sturmian morphism ψ , Proposition 7 gives a sharp upper bound on their number. The upper bound depends on the number of the elementary morphisms in the decomposition of ψ .

For our purposes, we do not fix the alphabet of a derivated word; two derivated words which differ only by a permutation of letters are identified one with another. Moreover, in the sequel, we work only with infinite words which are *uniformly recurrent*, i.e. each prefix w of \mathbf{u} occurs in \mathbf{u} infinitely many times and the set of all return words to w is finite (and, thus, the alphabet of the derivated word of \mathbf{u} with respect to w is finite). Our aim is to describe the set

$$\text{Der}(\mathbf{u}) = \{d_{\mathbf{u}}(w) : w \text{ is a prefix of } \mathbf{u}\}.$$

2. STURMIAN MORPHISMS

Let \mathcal{A} be a finite alphabet. A *morphism* over \mathcal{A}^* is a mapping $\psi : \mathcal{A}^* \mapsto \mathcal{A}^*$ such that $\psi(vw) = \psi(v)\psi(w)$ for all $v, w \in \mathcal{A}^*$. The domain of the morphism ψ can be naturally extended to $\mathcal{A}^{\mathbb{N}}$ by

$$\psi(u_0u_1u_2\cdots) = \psi(u_0)\psi(u_1)\psi(u_2)\cdots.$$

A morphism ψ is *primitive* if there exists a positive integer k such that the letter a occurs in the word $\psi^k(b)$ for each pair of letters $a, b \in \mathcal{A}$. A *fixed point* of a morphism ψ is an infinite word \mathbf{u} such that $\psi(\mathbf{u}) = \mathbf{u}$.

A morphism ψ is a *Sturmian morphism* if $\psi(\mathbf{u})$ is a Sturmian word for any Sturmian word \mathbf{u} . The set of Sturmian morphisms together with composition forms the so-called *Sturmian monoid* usually denoted St . We work with these four elementary Sturmian morphisms:

$$\varphi_a : \begin{cases} 0 \rightarrow 0 \\ 1 \rightarrow 10 \end{cases} \quad \varphi_b : \begin{cases} 0 \rightarrow 0 \\ 1 \rightarrow 01 \end{cases} \quad \varphi_\alpha : \begin{cases} 0 \rightarrow 01 \\ 1 \rightarrow 1 \end{cases} \quad \varphi_\beta : \begin{cases} 0 \rightarrow 10 \\ 1 \rightarrow 1 \end{cases}$$

and with the monoid \mathcal{M} generated by them, i.e. $\mathcal{M} = \langle \varphi_a, \varphi_b, \varphi_\alpha, \varphi_\beta \rangle$. The monoid \mathcal{M} is also called *special Sturmian monoid*. For a nonempty word $u = u_0 \cdots u_{n-1}$ over the alphabet $\{a, b, \alpha, \beta\}$ we put

$$\varphi_u = \varphi_{u_0} \circ \varphi_{u_1} \circ \cdots \circ \varphi_{u_{n-1}}.$$

The monoid \mathcal{M} is not free. It is easy to show that for any $k \in \mathbb{N}$ we have

$$\varphi_{\alpha a^k \beta} = \varphi_{\beta b^k \alpha} \quad \text{and} \quad \varphi_{a \alpha^k b} = \varphi_{b \beta^k a}.$$

We can equivalently say that the following rewriting rules hold on the set of words from $\{a, b, \alpha, \beta\}^*$:

$$(1) \quad \alpha a^k \beta = \beta b^k \alpha \quad \text{and} \quad a \alpha^k b = b \beta^k a \quad \text{for any } k \in \mathbb{N}.$$

In [12], the author reveals a presentation of the Sturmian monoid which includes the special Sturmian monoid $\mathcal{M} = \langle \varphi_a, \varphi_b, \varphi_\alpha, \varphi_\beta \rangle$. A presentation of the special Sturmian monoid follows from this result. It is also given explicitly in [9]:

Theorem 1. *Let $w, v \in \{a, b, \alpha, \beta\}^*$. The morphism φ_w is equal to φ_v if and only if the word v can be obtained from w by applying the rewriting rules (1).*

Note that the presentation of a generalization of the Sturmian monoid, the so-called *episturmian monoid*, is also known, see [11]. The next lemma summarizes several simple and well-known properties of Sturmian morphisms we exploit in the sequel.

Lemma 2. *Let $w \in \{a, b, \alpha, \beta\}^+$.*

- (i) *The morphism φ_w is primitive if and only if w contains at least one Greek letter α or β and at least one Latin letter a or b .*
- (ii) *If φ_w is primitive, then each of its fixed points is aperiodic and uniformly recurrent.*
- (iii) *If φ_w is primitive, then it has two fixed points if and only if w belongs to $\{a, \alpha\}^*$.*

For $w \in \{a, b, \alpha, \beta\}^*$ the rules (1) preserve positions in w occupied by Latin letters $\{a, b\}$ and positions occupied by Greek letters $\{\alpha, \beta\}$. We define that $a < b$ and $\alpha < \beta$ which allows the following definition.

Definition 3. Let $w \in \{a, b, \alpha, \beta\}^*$. The lexicographically greatest word in $\{a, b, \alpha, \beta\}^*$ which can be obtained from w by application of rewriting rules (1) is denoted $N(w)$. If $\psi = \varphi_w$, then the word $N(w)$ is the normalized name of the morphism ψ and it is also denoted by $N(\psi) = N(w)$.

3. DERIVATED WORDS OF FIXED POINTS OF STURMIAN MORPHISMS

Let $\psi \in \langle \varphi_a, \varphi_b, \varphi_\alpha, \varphi_\beta \rangle$ and $N(\psi) = w \in \{a, b, \alpha, \beta\}^* \setminus \{a, \alpha\}^*$ be the normalized name of the morphism ψ . The word w has a prefix $a^k\beta$ or $\alpha^k b$ for some $k \in \mathbb{N}$. This property enables us to define a transformation on the set of morphisms from $\mathcal{M} \setminus \langle \varphi_a, \varphi_\alpha \rangle$. This transformation is in fact the desired algorithm returning the morphisms $\psi_1, \psi_2, \dots, \psi_\ell$ mentioned above.

Definition 4. Let $w \in \{a, b, \alpha, \beta\}^* \setminus \{a, \alpha\}^*$ be the normalized name of a morphism ψ , i.e., $\psi = \varphi_w$. We put

$$\Delta(w) = \begin{cases} N(w'a^k\beta) & \text{if } w = a^k\beta w', \\ N(w'\alpha^k b) & \text{if } w = \alpha^k b w' \end{cases}$$

and, moreover, $\Delta(\psi) = \varphi_{\Delta(w)}$.

Theorem 5. Let $\psi \in \langle \varphi_a, \varphi_b, \varphi_\alpha, \varphi_\beta \rangle$ be a primitive morphism and $N(\psi) = w \in \{a, b, \alpha, \beta\}^* \setminus \{a, \alpha\}^*$ be its normalized name. Denote \mathbf{u} the fixed point of ψ . The word \mathbf{x} is (up to a permutation of letters) a derived word of \mathbf{u} with respect to one of its prefixes if and only if \mathbf{x} is the fixed point of the morphism $\Delta^j(\psi)$ for some $j \geq 1$.

Given a finite word u , we define the *cyclic shift* of $u = u_0 u_1 \dots u_{n-1}$ to be the word

$$\text{cyc}(u) = u_1 u_2 \dots u_{n-1} u_0.$$

Proposition 6. Let $w \in \{a, \alpha\}^*$ be the normalized name of a primitive morphism ψ and let a be its first letter.

- (i) Let \mathbf{u} be the fixed point of ψ starting with 0. Denote $v = b^{-1}N(wb) \in \{a, \beta\}^*$. We have $\text{Der}(\mathbf{u}) = \{\mathbf{v}\} \cup \text{Der}(\mathbf{v})$, where \mathbf{v} is the unique fixed point of the morphism φ_v .
- (ii) Let \mathbf{u} be the fixed point of ψ starting with 1. Put $v = \text{cyc}(w)$. We have $\text{Der}(\mathbf{u}) = \text{Der}(\mathbf{v})$, where \mathbf{v} is the fixed point of the morphism φ_v .

4. THE NUMBER OF DERIVATED WORDS

Proposition 7. If $w \in \{a, b, \alpha, \beta\}^* \setminus \{a, \alpha\}^*$ is the normalized name of a primitive Sturmian morphism $\psi = \varphi_w$ and \mathbf{u} is a fixed point of ψ , then

$$(2) \quad 1 \leq \#\text{Der}(\mathbf{u}) \leq 3|w| - 4.$$

Moreover, for any length $n \geq 2$ there exist normalized names $w', w'' \in \{a, b, \alpha, \beta\}^* \setminus \{a, \alpha\}^*$ of length n such that

- (i) $\varphi_{w'}$ and $\varphi_{w''}$ are not powers of other Sturmian morphisms,
- (ii) for the fixed points \mathbf{u}' and \mathbf{u}'' of the morphism $\varphi_{w'}$ and $\varphi_{w''}$, the lower resp. the upper bound in (2) is attained.

We also provide precise numbers of distinct derived words for these three types of morphisms:

- (1) ψ is a standard morphism from \mathcal{M} , i.e. $\psi \in \langle \varphi_b, \varphi_\beta \rangle$,
- (2) ψ is a standard morphism from $\mathcal{M} \circ E$, i.e. $\psi \in \langle \varphi_b, \varphi_\beta \rangle \circ E$,
- (3) ψ is a morphism from $\langle \varphi_a, \varphi_\alpha \rangle$.

To describe these numbers, we introduce the following morphism $F : \{a, b, \alpha, \beta\}^* \mapsto \{a, b, \alpha, \beta\}^*$ determined by

$$F(a) = \alpha, \quad F(\alpha) = a, \quad F(b) = \beta, \quad F(\beta) = b,$$

and we set

$$\text{cyc}_F(w_1 w_2 w_3 \dots w_n) = w_2 w_3 \dots w_n F(w_1)$$

for a finite word $w_1 w_2 w_3 \dots w_n$.

Proposition 8. *Let \mathbf{u} be a fixed point of a standard Sturmian morphism ψ which is not a power of any other Sturmian morphism.*

- (i) *If $\psi = \varphi_w$, then \mathbf{u} has $|w|$ distinct derivated words, each of them (up to a permutation of letters) is fixed by one of the morphisms*

$$\varphi_{v_0}, \varphi_{v_1}, \varphi_{v_2}, \dots, \varphi_{v_{|w|-1}}, \quad \text{where } v_k = \text{cyc}^k(w) \text{ for } k = 0, 1, \dots, |w| - 1.$$

- (ii) *If $\psi = \varphi_w \circ E$, then \mathbf{u} has $|w|$ distinct derivated words, each of them (up to a permutation of letters) is fixed by one of the morphisms*

$$\varphi_{v_0} \circ E, \varphi_{v_1} \circ E, \varphi_{v_2} \circ E, \dots, \varphi_{v_{|w|-1}} \circ E, \quad \text{where } v_k = \text{cyc}_F^k(w) \text{ for } k = 0, 1, \dots, |w| - 1.$$

Proposition 9. *Let $w \in \{\alpha, a\}^*$ be the normalized name of a primitive morphism ψ such that the letter a is a prefix of w . Moreover, assume that ψ is not a power of any other Sturmian morphism.*

- (i) *The fixed point of ψ starting with 0 has exactly $1 + |w|_\alpha$ distinct derivated words.*
(ii) *The fixed point of ψ starting with 1 has exactly $1 + |w|_a$ distinct derivated words.*

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k -Spectra of Strictly Balanced Words

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Abstract. A word u is a scattered factor of w if there exist u_1, u_2, \dots, u_n , and v_0, v_1, \dots, v_n such that $u = u_1 u_2 \dots u_n$ and $w = v_0 u_1 v_1 u_2 v_2 \dots u_n v_n$. We consider the set of length- k scattered factors (k -spectrum) of a given word w , denoted $\text{ScatFact}_k(w)$. We prove several properties of the sets $\text{ScatFact}_k(w)$ in the case of words w over a binary alphabet of length $2k$ for which the number of occurrences of each letter is equal. Such words are called strictly balanced. In particular, motivated by the task of recognising whether a set of words is a k -spectrum of some word w , we consider the question of which cardinalities $n = |\text{ScatFact}_k(w)|$ are obtainable for each k . We also consider the task of reconstructing words from their strictly balanced scattered factors.

1 Introduction

A scattered factor of w can be thought of as a representation of w in which some parts are missing. As such, there is considerable interest in the relationship of a word and its scattered factors from both a theoretical and practical point of view. For an introduction, see [3]. On the one hand, it is easy to imagine how, in any situation where discrete, linear data is read from an imperfect input – such as when sequencing DNA or during the transmission of a digital signal – scattered factors form a natural model, as multiple parts of the input may be missed, but the rest will remain unaffected and in-sequence. On the other hand, from a more theoretical perspective, there have been efforts to bridge the gap between the non-commutative field of combinatorics on words with traditional commutative mathematics via Parikh matrices (cf. e.g., [5, 6]) which are closely related to, and influenced by the topic of scattered factors.

One of the most fundamental questions about scattered factors of words and sets of scattered factors in general, is: given a set S of words (of length k), is S the set of scattered factors (or a k -spectrum) of some word w . In general, it remains a long standing goal of the theory to give a “nice” descriptive characterisation of scattered factor sets (and similarly, k -spectra), and to better understand their structure [3]. Another fundamental question concerning k -spectra, and one well motivated in several applications, is the question of reconstruction: given a word w of length n , for what values k does the k -spectrum of w uniquely determine w ? This question has generally had more success with definitive answers in a variety of cases. In particular, in [1], the exact bound of $\frac{n}{2} + 1$ is given in the general case.

Other variations, including for the definition of k -spectra where multiplicities are also taken into account, are considered in [4], while [2] considers the question of reconstructing words from their palindromic scattered factors.

In the current work, we consider the restricted setting of strictly balanced words: words over a binary alphabet $\{\mathbf{a}, \mathbf{b}\}$ with equal numbers of \mathbf{a} s and \mathbf{b} s. We show that the cardinality of their scattered factor sets ranges between $k + 1$ and 2^k and we prove for every $k + 1 \leq i \leq 3k - 2$ whether a k -spectrum of cardinality i exists. Moreover some results between $3k - 1$ and 2^k are given. In Section 4 we approach the question of reconstructing strictly balanced words from k -spectra in the specific case that the spectra are also limited to strictly balanced words only. While we are not able to resolve the question completely, we conjecture that the situation is similar to the general case; we show that this bound holds in the case that w contains at most two blocks of \mathbf{b} s.

Before we are able to present our results, we need to define the setting of strictly balanced words. We consider words w over an alphabet $\Sigma = \{\mathbf{a}, \mathbf{b}\}$. The number of occurrences of a letter $\mathbf{a} \in \Sigma$ in a word $w \in \Sigma^*$ is denoted by $|w|_{\mathbf{a}}$. The subset of Σ^* which contains only words with equal numbers of occurrences of letters is defined by $\Sigma_{sb}^* = \{w \in \Sigma^* \mid \forall x, y \in \Sigma : |w|_x = |w|_y\}$ and these words are called *strictly balanced*. For example, \mathbf{abaa} is not strictly balanced, while \mathbf{abbaba} is.

Definition 1. A word $u = a_1 \dots a_n \in \Sigma^n$, for $n \in \mathbb{N}$, is a scattered factor of a word $w \in \Sigma^+$ if there exists $v_0, \dots, v_n \in \Sigma^*$ with $w = v_0 a_1 v_1 \dots v_{n-1} a_n v_n$. Let $\text{ScatFact}(w)$ denote the set of w 's scattered factors and consider additionally $\text{ScatFact}_k(w)$ (full k -spectrum) and $\text{ScatFact}_{\leq k}(w)$ (k -spectrum) as the two subsets of $\text{ScatFact}(w)$ which contain only the scattered factors of length $k \in \mathbb{N}$ or the ones up to length $k \in \mathbb{N}$.

We note two obvious, but important symmetries regarding k -spectra: for $w \in \Sigma^*$, $\text{ScatFact}(w^R) = \{u^R \mid u \in \text{ScatFact}(w)\}$ and $\text{ScatFact}(\bar{w}) = \{\bar{u} \mid u \in \text{ScatFact}(w)\}$ hold with the renaming morphism $\bar{\cdot}$. Thus, from a structural point of view, it is sufficient to consider only one representative (here the lexicographically smallest with $\mathbf{a} < \mathbf{b}$) from the equivalence classes.

2 Cardinalities of k -Spectra of Strictly Balanced Words

In the current section, we are interested in the cardinalities of the k spectra, and in the question: which cardinalities are not possible? It is a straightforward observation that not every subset of Σ^k is a k -spectrum of some word w . For example \mathbf{aa} and \mathbf{bb} can only be scattered factors of a word containing both \mathbf{a} s and \mathbf{b} s, and therefore having either \mathbf{ab} or \mathbf{ba} as a scattered factor. In fact, for $k = 2$, the sets $\{\mathbf{aa}, \mathbf{ab}, \mathbf{bb}\}$ and $\{\mathbf{aa}, \mathbf{ba}, \mathbf{bb}\}$ are the smallest possible k -spectra of words of length $2k$ in both the general case, and when restricted to strictly balanced words only. Moreover these sets are equivalent in the sense that one is a renaming (or a reversal) of the other. Note that the largest possible set in this case is $\{\mathbf{aa}, \mathbf{ab}, \mathbf{ba}, \mathbf{bb}\}$ which has size $4 = 2k = 2^k$. Our first result generalises the previous observation about minimal-size and maximal-size k -spectra.

Lemma 1. *For all $k \in \mathbb{N}$, the smallest reachable cardinality for any $w \in \Sigma_{sb}^{2k}$ is $|\text{ScatFact}_k(w)| = k + 1$, reached exactly for $w = \mathbf{a}^k \mathbf{b}^k$ (up to renaming and reversal), and $\text{ScatFact}_k(\mathbf{a}^k \mathbf{b}^k) = \{\mathbf{a}^r \mathbf{b}^s \mid r + s = k, r, s \in [k]_0\}$ holds.*

Lemma 2. *Let $k \in \mathbb{N}$. Then $w \in \{ab, ba\}^k$ if and only if $\text{ScatFact}_k(w) = \Sigma^k$.*

By the Lemmas 1 and 2, the characterisation for the smallest and the largest closure w.r.t. cardinality of the given set S are given. Now the gap in between will be investigated. Since there does not exist a gap for $k = 2$, assume $k \in \mathbb{N}_{\geq 3}$. The following two statements show that $2^k - 1$ and $2k$ are always reachable and thus the possible cardinalities for $k = 3$ are fully characterised.

1. $|\text{ScatFact}_k(w)| = 2^k - 1$ iff $w \in \{(\mathbf{ab})^i \mathbf{a}^2 \mathbf{b}^2 (\mathbf{ab})^{k-i-2} \mid i \in [k-2]_0\}$ (in particular $\text{ScatFact}_k(w) = \Sigma^k \setminus \{\mathbf{b}^{i+1} \mathbf{a}^{k-i-1}\}$),
2. $|\text{ScatFact}_k(w)| = 2k$ iff $w \in \{\mathbf{a}^{k-1} \mathbf{bab}^{k-1}, \mathbf{a}^{k-1} \mathbf{b}^k \mathbf{a}\}$

The cardinality of $2k$ is important since there is a gap between $k + 1$ and $2k$, i.e. $\forall w \in \Sigma_{sb}^{2k} : |\text{ScatFact}_k(w)| \notin \{k + 2, \dots, 2k - 1\}$. This shows that with increasing k the number of possible cardinalities at the *beginning* of the scala from $k + 1$ to 2^k decreases: the larger k is the more unlikely it is somehow to find a k -spectrum of a small cardinality. To investigate the second gap we have $|\text{ScatFact}_k(\mathbf{a}^{k-i} \mathbf{b}^k \mathbf{a}^i)| = k(i+1) - i^2 + 1$ for $i \in [\lfloor \frac{k}{2} \rfloor]$. It is worth noting that this includes all square numbers being at least four: $|\text{ScatFact}_k(\mathbf{a}^{\frac{k}{2}} \mathbf{b}^k \mathbf{a}^{\frac{k}{2}})| = (\frac{k}{2} + 1)^2$ holds for k even. Moreover $|\text{ScatFact}_k(\mathbf{a}^{k-2} \mathbf{b}^k \mathbf{a}^2)| = 3k - 3$ holds. This result is important since it will be shown in the following that the cardinalities $2k + 1$ up to $3k - 4$ are not reachable. In other words $\mathbf{a}^{k-2} \mathbf{b}^k \mathbf{a}^2$ delivers the third smallest cardinality after $k + 1$ and $2k$. Contrarily the cardinality $3k - 2$ belongs to the word $\mathbf{a}^{k-1} \mathbf{b}^2 \mathbf{ab}^{k-2}$.

Proposition 1. *For $k \geq 5$, no word $w \in \Sigma_{sb}^{2k}$ has k -spectrum of cardinality $2k + i$ for $i \in [k - 4]$, i.e. between $2k + 1$ and $3k - 4$ is a cardinality-gap.*

We will end this analysis with the conjecture that in contrast to the first gap, the last gap ends earlier the larger k is. More precisely, if for $k \in \mathbb{N}_{\geq 4}$ and $i \in [k - 2]_0$, $w = \mathbf{a}^2 \mathbf{b}^2 (\mathbf{ab})^{k-3-i} \mathbf{ba} (\mathbf{ab})^i$ holds then $|\text{ScatFact}_k(w)| = 2^k - 2 - i$ follows. Notice that this conjecture implies that indeed similar to the second gap here $4k - 4$ is always reached. On the other hand, in contrast to the second gap, the third gap is not of the form $4k - 4 - i$ for $i \in [k - 4]$.

3 Reconstructing Strictly Balanced Words from their k -Spectra

As with the general case, it is easy to see that strictly balanced words of length $2k$ are not uniquely determined by their scattered factors of length k . In the current section we discuss the question of when a strictly balanced word w of length $2k$ is uniquely identified by the set $\text{ScatFact}_{k'}(w) \cap \Sigma_{sb}^{k'}$ for $2k > k' > k$.

Of course if k' is odd then $\text{ScatFact}_{k'}(w) \cap \Sigma_{sb}^{k'} = \emptyset$ for all words w , so in these cases the answer is trivially negative. In the general case, Dress and Erdős [1] showed, that if $\text{ScatFact}_{k+1}(w) = \text{ScatFact}_{k+1}(w')$ holds for $w, w' \in \Sigma^{2k}$ then $w = w'$ follows. If w is strictly balanced we found a straightforward proof for their proposition. However, in both proofs, there is a necessity in some cases to consider scattered factors u consisting mostly of **as** or mostly of **bs** – i.e., that do not belong to Σ_{sb}^* . Thus it remains an open problem whether the same bound of $k+1$ (or in the case that k is even, $k+2$) is sufficient. While we do not resolve the question completely, we conjecture that these bounds do still hold.

Conjecture 1. Let $k \in \mathbb{N}$. Let $k' = k+1$ if k is odd, and $k' = k+2$ if k is even. Let $w, w' \in \Sigma_{sb}^{2k}$ such that $\text{ScatFact}_{k'}(w) = \text{ScatFact}_{k'}(w')$. Then $w = w'$.

It is possible to show that the conjecture holds when there are at most two blocks of **bs** (by symmetry at most two blocks of **as**), i.e. $w \in \mathbf{a}^*\mathbf{b}^*\mathbf{a}^*\mathbf{b}^*\mathbf{a}^* \cap \Sigma_{sb}^{2k}$:

- for k odd, w is uniquely determined by $\text{ScatFact}_{k+1}(w) \cap \Sigma_{sb}^k$,
- for k even, w is uniquely determined by $\text{ScatFact}_{k+2}(w) \cap \Sigma_{sb}^k$.

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Multidimensional subshifts of finite type are discrete dynamical systems as a set of colorings of an infinite regular grid with elements of a finite set \mathcal{A} together with the shift action. The set of colorings is defined by forbidding a finite set of patterns all over the grid (also called local rules). The most simple and most considered grids of this type are \mathbb{Z}^2 and more generally \mathbb{Z}^d for $d \geq 1$. In this case, one can consider a coloring as a bi-dimensional and infinite word on the alphabet \mathcal{A} .

They are notably involved in statistical physics in the study of so-called lattice models. These models are often simple to describe : for instance, the hard square model is defined on alphabet $\mathcal{A} = \{0, 1\}$ and by forbidding two 1 to appear on horizontally or vertically adjacent positions of the lattice. However, these models and their physical constants, such as the entropy are difficult to apprehend with general methods, and involve specific properties of the considered model.

Although it is known that it is possible to compute the entropy of one-dimensional version of these models by computing the greatest eigenvalue of a matrix which derives from the description of the subshift, this is not possible for multidimensional subshifts. This is the consequence of a result by M. Hochman and T. Meyerovitch in 2010, which states that the possibles values of the entropy for multidimensional subshifts of finite type are the Π_1 -computable numbers, including in particular non-computable numbers.

The method developed for this purpose originates in the work of R. Berger and R. Robinson. It has been developed further in order to characterize other dynamical aspects of SFT with computability conditions, with similar constructions. It consists in the implementation of Turing machines in hierarchical structures that emerge from the local rules.

However, models studied in statistical physics obey to strong dynamical constraints and there is still hope to include them into a sub-class of subshifts of finite type for which the entropy is uniformly computable (this means that there is an algorithm which can provide arbitrarily precise approximations of the entropy, provided the precision and the local rules of the subshift). An example of a constraint defining a class where this is verified is the block gluing : this was proved by R. Pavlov and M. Schraudner. This property means that two square blocks can be viewed in any relative positions in some element of the subshift provided that the distance between the two blocks is sufficiently large, with minimal distance not depending on the size of the blocks) is a computable real number. Although they provided a construction to realize some class numbers as entropy of block gluing SFT, they did not prove a characterization, and this problem seems difficult. However, it could be possible to find broader class for which the entropy is still computable.

A strategy to understand the limit between the general regime where Hochman and Meyerovitch's result holds and this restricted block gluing class is to quantify this property. This means imposing that two patterns can be glued in any two positions in a configuration of the subshift, provided that the distance is great enough, where the minimal distance is a linear function of the size of these patterns.

In a work with Mathieu Sablik, we made a step towards the limit, proving that the result of Hochman and Meyerovitch is robust under the linear version of this property (where the minimal distance function is $O(n)$ where n is the size of the two square blocks). The aim of this talk would be, after a presentation of the problem, to give an insight on the obstacles to this property in the initial construction of Hochman and Meyerovitch, using a construction slightly simpler to present, and on the methods used to overcome the obstacles. These methods involves in particular a modification of the Turing machine model and an operator on subshifts that acts by distortion.

Unary Patterns of Size Four with Morphic Permutations

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Abstract. We investigate the avoidability of unary patterns of size of four with morphic permutations. More precisely, we show how to identify precisely, given the positive integers i, j, k , the alphabets over which a pattern $x\pi^i(x)\pi^j(x)\pi^k(x)$ is avoidable, where x is a word variable and π is a function variable with values in the set of all morphic permutations of the respective alphabets. This continues the work of [Manea et al., 2015], where a complete characterisation of the avoidability of cubic patterns with permutations was given.

1 Introduction

The avoidability of patterns in infinite words is an old area of interest with a first systematic study going back to Thue. In these initial papers it was shown that there exist a binary infinite morphic word and a ternary infinite morphic word that avoid cubes and squares, respectively. That is, these infinite words do not contain instances of the patterns xxx and xx , respectively.

In this article, we are studying the avoidability of repetitions in a generalised setting. Namely, we are interested in the avoidability of unary patterns with functional dependencies between variables. We are considering patterns like $x\pi^i(x)\pi^j(x)\pi^k(x)$, where x is a word variable while π is function variable, which can be replaced by bijective morphisms only. The instances of such patterns over an alphabet Σ are obtained by replacing x with a concrete word, and π by a morphic permutation of Σ . For example, an instance of the pattern $x\pi(x)x\pi(x)$ over $\Sigma = \{a, b\}$ is the word $uvuv$ such that $|u| = |v|$, and v is the image of u under any permutation on the alphabet. Considering the permutation $a \rightarrow b$, and $b \rightarrow a$, then $aba|bab|aba|bab$ is an instance of $x\pi(x)x\pi(x)$.

In this setting, we continue the work of [Manea et al., 2015] as follows. In that paper, a complete characterisation of the avoidability of cubic patterns with permutations $x\pi^i(x)\pi^j(x)$ was given. Furthermore, it was shown that there exists a ternary word that avoids all patterns $\pi_{i_1}(x) \dots \pi_{i_r}(x)$ where $r \geq 4$, x a word variable over some alphabet Σ , with $|x| \geq 2$ and $|\Sigma| \geq 3$, and the π_{i_j} function variables that may be replaced by anti-/morphic permutations of Σ . However, this result only holds when the length of x is restricted to be at least 2. Also, it was shown that all patterns $\pi^{i_1}(x) \dots \pi^{i_n}(x)$ with $n \geq 4$ under morphic permutations are avoidable in alphabets of size 2, 3, and 4, but there

exist patterns which are unavoidable in alphabets of size 5. We extend these results by showing how to determine exactly, for a given unary pattern \mathcal{P} of size four with permutations, which are the alphabets in which it is avoidable.

The main result of our paper is that given i, j, k , we show how to compute the value m such that the pattern $x\pi^i(x)\pi^j(x)\pi^k(x)$, with $i, j, k \geq 0$, is unavoidable in alphabets of size at least m and avoidable in alphabets of size $2, 3, 4, \dots, m-1$. To achieve this, we define a series of parameters that allow us to characterize, for each alphabet, what form the instances of the pattern may have over a certain alphabet. Then, we show that for each pattern there exists an interval (whose left end is 2 and right end is defined based on the respective parameters) such that over each alphabet whose size is in the respective interval, there exists an infinite word that does not contain instances of the given pattern. The structure of the paper is as follows: we first give a series of basic definitions and preliminary results. Then we define the aforementioned parameters, and show how to use them to compute, for a given pattern, the minimum size σ of an alphabet over which the respective pattern is unavoidable. Finally, we show the correctness of the computation done in the previous step: for alphabets with less than σ symbols the pattern is avoidable.

2 Preliminaries

We define $\Sigma_k = \{0, \dots, k-1\}$ to be an alphabet with k letters; the empty word is denoted by ε . For words u and w , we say that u is a prefix (resp. suffix) of w , if there exists a word v such that $w = uv$ (resp. $w = vu$). If $f : \Sigma_k \rightarrow \Sigma_k$ is a permutation, we say that the order of f , denoted $\mathbf{ord}(f)$, is the minimum value $m > 0$ such that f^m is the identity. If $a \in \Sigma_k$ is a letter, the order of a with respect to f , denoted $\mathbf{ord}_f(a)$, is the minimum number m such that $f^m(a) = a$.

In this paper, we consider only unary patterns (i.e., containing only one variable) with morphic permutations, that is, all function variables are unary and are substituted by morphic permutations only.

The infinite Hall word h is defined as $h = \lim_{n \rightarrow \infty} \phi_h^n(0)$, for the morphism $\phi_h : \Sigma_3^* \rightarrow \Sigma_3^*$ where $\phi_h(0) = 012$, $\phi_h(1) = 02$ and $\phi_h(2) = 1$. The infinite word h avoids the pattern xx (squares).

3 Avoidability of patterns under permutations

In this section we try to identify an upper bound on the size of the alphabets Σ_m in which a patterns $x\pi^i(x)\pi^j(x)\pi^k(x)$, with $i, j, k \geq 0$ is unavoidable, when π is substituted by a morphic permutation.

In the pattern $x\pi^i(x)\pi^j(x)\pi^k(x)$, the factors x , $\pi^i(x)$, π^j , or $\pi^k(x)$ are called x -items in the following. Our analysis is based on the relation between the possible images of the four x -items occurring in a pattern, following the ideas of [?]. For instance, we want to check whether in a possible image of our pattern, all four x -items can be mapped to a different word, or whether the second and the last x -items can be mapped to the same word, etc.

To achieve this, we define in Table 1 the parameters α_i , with $1 \leq i \leq 14$.

| | |
|---|------|
| $\alpha_1 = \inf\{t : t \nmid i, t \nmid j, t \nmid k, t \nmid i-j , t \nmid i-k , t \nmid j-k \}$ | 0123 |
| $\alpha_2 = \inf\{t : t \mid i, t \nmid j, t \nmid k, t \nmid j-k \}$ | 0012 |
| $\alpha_3 = \inf\{t : t \nmid i, t \mid j, t \nmid k, t \nmid i-k \}$ | 0102 |
| $\alpha_4 = \inf\{t : t \nmid i, t \nmid j, t \mid i-k \}$ | 0121 |
| $\alpha_5 = \inf\{t : t \nmid i, t \nmid j, t \nmid i-j , t \nmid i-k , t \mid j-k \}$ | 0122 |
| $\alpha_6 = \inf\{t : t \mid i, t \mid j, t \nmid k\}$ | 0001 |
| $\alpha_7 = \inf\{t : t \mid i, t \nmid j, t \mid k\}$ | 0010 |
| $\alpha_8 = \inf\{t : t \nmid i, t \mid j, t \mid k\}$ | 0100 |
| $\alpha_9 = \inf\{t : t \nmid i, t \mid i-j , t \mid i-k \}$ | 0111 |
| $\alpha_{10} = \inf\{t : t \mid i, t \nmid j, t \mid j-k \}$ | 0011 |
| $\alpha_{11} = \inf\{t : t \nmid i, t \mid j, t \mid i-k \}$ | 0101 |
| $\alpha_{12} = \inf\{t : t \nmid i, t \mid k, t \mid i-j \}$ | 0110 |
| $\alpha_{13} = \inf\{t : t \nmid i, t \nmid k, t \mid i-j \}$ | 0112 |
| $\alpha_{14} = \inf\{t : t \nmid i, t \nmid j, t \mid i-j \}$ | 0120 |

Table 1. Definition of the values α_i , with $1 \leq i \leq 14$.

Now based on combinatorial relations, we define the some collections of sets. The idea behind all these collections is to generate sets of parameters α_i s that cannot be avoided and have a minimal cardinality. No matter what will be added to these sets will preserve their unavoidability, while erasing something from them will make them avoidable. To obtain these collections we used a computer program and randomly generated some unavoidable sets of parameters of size five. Using the similarities between the instances modelled by these sets, defined in terms of (gapped) squares and cubes occurring in their digit representation, we developed an algorithm to generate more sets of patterns. Based on these relations we constructed fourteen sets \mathcal{S}_i s . We just define one of them as an example.

Let \mathcal{S}_1 be the collection of sets (each with five elements) that contain α_1 and:

- one of the α_i s whose representation has a prefix or a suffix square, but no gapped cube. That is: α_2 or α_5 .
- one of the α_i s that has a gapped square, but does not have two gapped squares. These are α_3 or α_4 .
- one of the α_i s that contain cubes or two squares: α_6 or α_9 or α_{10} .
- one of the α_i s that contain gapped cubes: α_7 or α_8 .

For example, one possible set from \mathcal{S}_1 is $\{\alpha_1, \alpha_2, \alpha_4, \alpha_6, \alpha_7\}$.

Lemma 1 *Let $K' \subset K$ be any subset of size at most 4 of K . There exists an infinite word w such that w does not contain 4-powers and if w contains an instance of the pattern $x\pi^i(x)\pi^j(x)\pi^k(x)$ then it can not be modelled by any tuples of the set of patterns K' .*

Theorem 1. *Let i, j, k be positive integers such that $i \neq j \neq k \neq i$, and consider the pattern $p = x\pi^i(x)\pi^j(x)\pi^k(x)$. Let $\sigma = \min\{\max(S) \mid S \in \cup_{1 \leq i \leq 13} \mathcal{S}_i\}$. Then $\sigma \geq 5$ and p is unavoidable in Σ_m , for $m \geq \sigma$.*

Proof. We briefly prove this Theorem. The complete proof is in the main paper. We checked with the aid of a computer, by a straightforward backtracking algorithm, that if $m \geq \max(S)$, for some $S \in \cup_{1 \leq i \leq 13} S_i$, then p is unavoidable in Σ_m . Our computer program tries construct a word as long as possible by always adding a letter to the current word it constructed by backtracking; this letter is chosen in all possible ways from the letters contained in the word already, or it may also be a new letter. \square

4 Algorithm to generate avoidable cases

Algorithm 1 Algorithm to generate avoidable cases

- 1: Let $n = 13$. Using the sets S_i , ($1 \leq i \leq 13$), generate all sets of α_i s of cardinality n , that have no unavoidable sets of patterns as subset; show that they are avoidable;
 - 2: For all n from 12 downto 4, generate all sets of cardinality n that have no unavoidable sets of patterns as subset; these sets should not be subsets of the avoidable sets of α_i s of cardinality $n+1$ (to avoid generating repetitive avoidable sets of cases generated in the past step); show that they are avoidable.
-

Theorem 2. *Given a pattern $p = x\pi^i(x)\pi^j(x)\pi^k(x)$ we can determine effectively the values m , such that the pattern is avoidable in Σ_m .*

Proof. We briefly prove this Theorem. The complete proof is in the main paper. In Theorem 1, we proved that given the pattern $p = x\pi^i(x)\pi^j(x)\pi^k(x)$, for each i, j , and k , we can compute the cardinality of an alphabet over which the pattern is unavoidable. Now to show that this is the minimum cardinality over which the pattern of size four is unavoidable, we proceed as follows. We will show that the subsets and the complements of all the sets S_i , ($1 \leq i \leq 13$) are avoidable. By complement we mean here all the sets of parameters α_i of which the sets Patterns of Size Four with Morphic Permutations are not subsets of. The reason to define it this way is that if a set of parameters is unavoidable, whatever we add to it remains unavoidable, so this set should not be subset of any avoidable set of patterns. Furthermore, to show that the value $\alpha = \alpha_i$, $1 \leq i \leq 14$ in the set S_j , $1 \leq j \leq 13$, is the minimum cardinality of an alphabet over which the pattern of size four is unavoidable, we should prove that α_i is the minimum value such that if we add it to the set $S_j \setminus \alpha_i$, makes it unavoidable set of patterns. To reach this, we proved that all proper subsets of the sets S_j , $1 \leq j \leq 13$ are avoidable sets of patterns. \square

On closed and open factors of Arnoux-Rauzy words *

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Given a finite non-empty set \mathbb{A} , let $\mathbb{A}^{\mathbb{N}}$ denote the set of (right) infinite words $x = x_1x_2x_3 \cdots$ with $x_i \in \mathbb{A}$. For each infinite word $x = x_1x_2x_3 \cdots \in \mathbb{A}^{\mathbb{N}}$, the factor complexity $p_x(n)$ counts the number of distinct blocks (or factors) $x_ix_{i+1} \cdots x_{i+n-1}$ of length n occurring in x . First introduced by Hedlund and Morse in their seminal 1938 paper [13] under the name of *block growth*, the factor complexity provides a useful measure of the extent of randomness of x . Periodic words have bounded factor complexity while digit expansions of normal numbers have maximal complexity. A celebrated theorem of Morse and Hedlund in [13] states that every aperiodic (meaning not ultimately periodic) word contains at least $n + 1$ distinct factors of each length n . Sturmian words are those aperiodic words of minimal factor complexity: $p_x(n) = n + 1$ for each $n \geq 1$.

Other notions of complexity have been successfully used in the study of infinite words and their combinatorial properties [1, 5, 6, 7, 15, 16]. In this note, we introduce and study two new complexity functions based on the notions of open and closed words [8]. We recall that a word $u \in \mathbb{A}^+$ is said to be *closed* if either $u \in \mathbb{A}$ or if u is a complete first return to some proper factor $v \in \mathbb{A}^+$, meaning u has precisely two occurrences of v , one as a prefix and one as a suffix. Otherwise, if u is not closed then u is *open*. For example, *abbbab* and *aabaaabaa* are both closed words while *ab* and *abaabbababbaaba* are both open. It is easily seen that all privileged words [15] are closed and hence so are all palindromic factors of rich words [9]. The terminology open and closed was first introduced by the authors in [3] although the notion of a closed word had already been introduced earlier by A. Carpi and A. de Luca in [4]. For a nice overview of open and closed words we refer the reader to the recent survey article by G. Fici [8].

To each infinite word $x \in \mathbb{A}^{\mathbb{N}}$ we consider the functions $f_x^c, f_x^o : \mathbb{N} \rightarrow \mathbb{N}$ which count the number of closed and open factors of x of each length $n \in \mathbb{N}$. We study the behaviour of these complexity functions for Arnoux-Rauzy words [2]. Recall an infinite word $x \in \mathbb{A}^{\mathbb{N}}$ is called an *Arnoux-Rauzy word* if it is recurrent and if x contains, for each $n \geq 0$, precisely one right special factor of length n which is a prefix of $|\mathbb{A}|-$ many factors of x of length $n + 1$ and precisely one left special factor of length n which is a suffix of $|\mathbb{A}|-$ many factors of x .

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of length $n + 1$. In particular one has $p_x(n) = (|\mathbb{A}| - 1)n + 1$ and each factor u of x has precisely $|\mathbb{A}|$ distinct complete first returns. Arnoux-Rauzy words were first introduced in [2] in the special case of a 3-letter alphabet. Let us note that in case $|\mathbb{A}| = 2$, then x is Sturmian. Since for any word $x \in \mathbb{A}^{\mathbb{N}}$ we have that $f_x^c(n) + f_x^o(n) = p_x(n)$, it suffices to understand the behaviour of $f_x^c(n)$.

Our main result in Theorem 1 below provides an explicit formula for the closed complexity function $f_x^c(n)$ for an Arnoux-Rauzy word x on a t -letter alphabet \mathbb{A} . The formula is expressed in terms of two related sequences associated to x . The first is the sequence $(b_k)_{k \geq 0}$ of the lengths of the bispecial factors $\varepsilon = B_0, B_1, B_2, \dots$ of x , ordered in increasing length. The second is the sequence $(p_a^{(k)})_{a \in \mathbb{A}}^{k \geq 0}$ where for each $k \geq 0$, the t coordinates of $(p_a^{(k)})_{a \in \mathbb{A}}$ are the lengths of the t first returns in x to B_k . More precisely, $p_a^{(k)} = |R_a^{(k)}| - b_k$ where $R_a^{(k)}$ is the complete first return to B_k in x beginning in $B_k a$. Both sequences have already been extensively studied in the literature. In particular, following [11] one has that

$$b_k = \frac{\sum_{a \in \mathbb{A}} p_a^{(k)} - t}{t - 1}.$$

Furthermore, for each $k \in \mathbb{N}$, the coordinates of $(p_a^{(k)})_{a \in \mathbb{A}}$ are coprime and each is a period of the word B_k . Moreover, B_k is an extremal Fine and Wilf word i.e., any word u having periods $(p_a^{(k)})_{a \in \mathbb{A}}$ and of length greater than b_k is a constant word, i.e., $u = a^n$ for some n (see [17]). The sequence $(p_a^{(k)})_{a \in \mathbb{A}}^{k \geq 0}$ is computed recursively as follows : $p_a^{(0)} = 1$ for each $a \in \mathbb{A}$. For $k \geq 1$, let $a \in \mathbb{A}$ be the unique letter such that aB_{k-1} is a right special factor of x . Then $p_a^{(k)} = p_a^{(k-1)}$, and $p_b^{(k)} = p_b^{(k-1)} + p_a^{(k-1)}$ for $b \in \mathbb{A} \setminus \{a\}$.

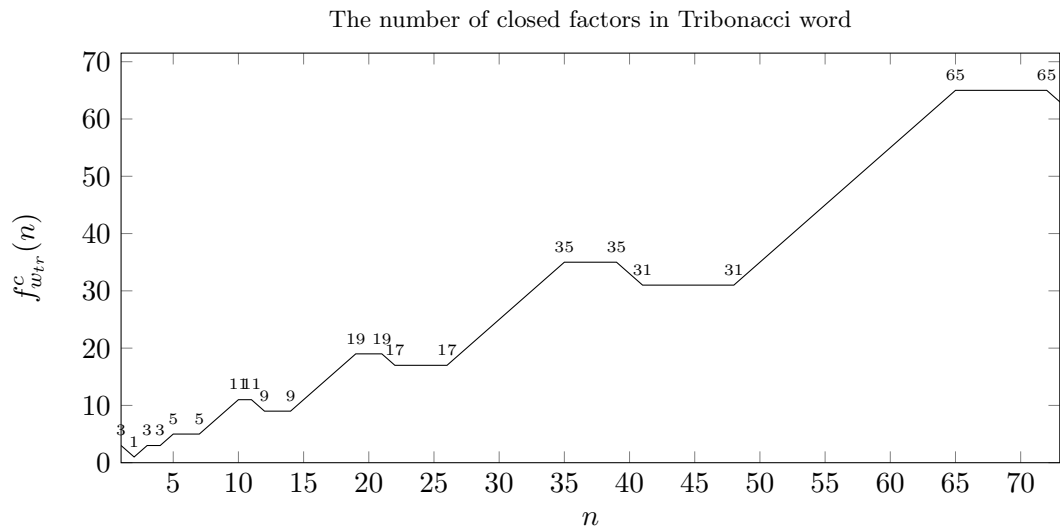
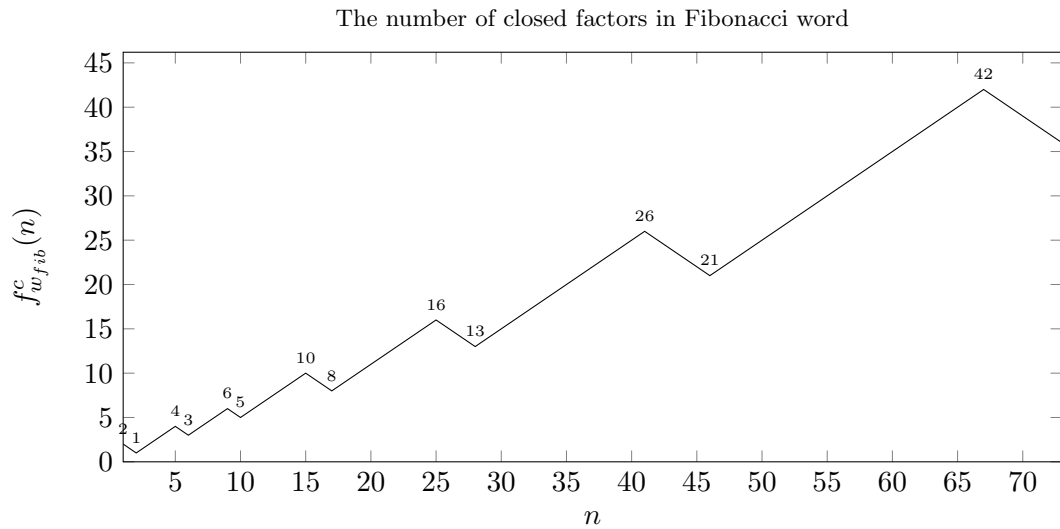
Theorem 1. *Let $x \in \mathbb{A}^{\mathbb{N}}$ be an Arnoux-Rauzy word. For each $k \in \mathbb{N}$ and $a \in \mathbb{A}$ set $I_{k,a} = [b_{k-1} - p_k + p_a^{(k)} + 2, b_k + p_a^{(k)}]$ where $p_k = \min_{b \in \mathbb{A}} \{p_b^{(k)}\}$. Let*

$$F(a, n) = \sum_{\substack{k \in \mathbb{N} \\ n \in I_{k,a}}} (d(n, I_{k,a}) + 1) \tag{1}$$

where for $n \in I_{k,a}$, the quantity $d(n, I_{k,a})$ denotes the minimal distance from n to the endpoints of the interval $I_{k,a}$. Then the number of closed factors of x for each length n is $f_x^c(n) = \sum_{a \in \mathbb{A}} F(a, n)$.

It is easily checked that the length of each interval $I_{k,a}$ is $2p_k - 2$ and that for each fixed n the sum in (1) is actually a finite sum.

The following figures illustrate the behaviour of the closed complexity function f_x^c in the case of the Fibonacci word and the Tribonacci word.



It is evident that in general f_x^c is not monotone. However as a consequence to Theorem 1 we are able to show :

Corollary 2. *Let x be an Arnoux-Rauzy word. Then $\liminf f_x^c(n) = +\infty$.*

In contrast, it is shown in [14] that for any paperfolding word x , $\liminf f_x^c(n) = 0$, in other words, for infinitely many n , x has no closed factors of length n .

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ON THE GROUP OF A RATIONAL MAXIMAL BIFIX CODE

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AND DOMINIQUE PERRIN

ABSTRACT. We give necessary and sufficient conditions for the group of a rational maximal bifix code Z to be isomorphic with the F -group of $Z \cap F$, when F is recurrent and $Z \cap F$ is rational. The case where F is uniformly recurrent receives special attention.

1. INTRODUCTION

In the past few years, special attention has been given to bifix codes which may not be maximal but are maximal within some language, which is usually chosen to be recurrent or uniformly recurrent. This line of research has produced new and strong connections between bifix codes, subgroups of free groups and symbolic dynamical systems (cf. [4] and the sequels [5, 6, 7, 8]).

If Z is a thin maximal bifix code and F is a recurrent set, then $X = Z \cap F$ is an F -maximal bifix code, that is, a maximal bifix code within F (with X finite if F is uniformly recurrent) [4]. This leads to a process of “relativization” of several previously known definitions for maximal codes. An important example is the group $G(Z)$ of a rational code Z , i.e., the Schützenberger group of the minimum ideal $J(Z)$ of the syntactic monoid $M(Z^*)$ of Z^* . In this case, the relativization consists in taking the intersection $X = Z \cap F$ and the Schützenberger group of the minimum \mathcal{J} -class $J_F(X)$ that intersects the image of F in $M(X^*)$, when X is rational. This group, denoted by $G_F(X)$, is the F -group of X . How are $G(Z)$ and $G_F(X)$ related? They are not always isomorphic, even if Z is a group code (i.e., Z is a code with $M(Z^*)$ a finite group) and F is uniformly recurrent. In [4] it is shown that if Z is a group code and F is Sturmian, then $G(Z)$ and $G_F(X)$ are isomorphic. This was extended to tree sets in the manuscript [10], thanks to a novel approach consisting in exploring links between $G(Z)$, $G_F(X)$ and the Schützenberger (profinite) group $G(F)$ of the minimum \mathcal{J} -class $J(F)$ of the topological closure of F within the free profinite monoid generated by the alphabet of F , and, with the help of these links, taking advantage of results on $G(F)$ from [2, 3]. Building on this approach, we get new results about when $G(Z) \simeq G_F(X)$ holds, recovering previous results in the process.

2. PRELIMINARIES ON FREE PROFINITE MONOIDS

Here A is always a finite alphabet. Take $u, v \in A^*$. If $u \neq v$, there is a homomorphism $\varphi: A^* \rightarrow M$ onto a finite monoid such that $\varphi(u) \neq \varphi(v)$.

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Let $r(u, v)$ be the minimum for $|M|$. We may consider the metric d on A^* such that $d(u, v) = 2^{-r(u, v)}$ if $u \neq v$. The *free profinite monoid* $\widehat{A^*}$ is the compact monoid resulting from the completion of A^* under d , a terminology justified as $\widehat{A^*}$ is the free object in the class of A -generated profinite monoids. For an extended introduction, see [11]. The elements of $\widehat{A^*}$ are called *pseudowords* over A . Words of A^* are topologically isolated in $\widehat{A^*}$. Pseudowords generalize words, but the structure of $\widehat{A^*}$ is richer than that of A^* . Next is a glimpse of that. If F is a factorial subset of A^* , then the topological closure \overline{F} is itself factorial in $\widehat{A^*}$, and when F satisfies the stronger property of being recurrent, there is a minimum \mathcal{J} -class $J(F)$ contained in \overline{F} , which moreover is regular. Maximal subgroups of $J(F)$ were identified in [2, 3]. We mention that the proper factors (that is, strictly \mathcal{J} -above) of $u \in \widehat{A^*} \setminus A^*$ belong to A^* if and only if $u \in J(F)$ for some uniformly recurrent set F [1].

3. PREPARATORY TECHNICAL RESULTS

In this section we prepare the main results of the next section.

Recall that a *parse* of a word w with respect to a subset X of A^* is a triple (v, x, u) such that $w = vxu$ with $v \in A^* \setminus A^*X$, $x \in X^*$ and $u \in A^* \setminus XA^*$. The number of parses of w with respect to X is denoted by $\delta_X(w)$. The F -degree of X is $d_F(X) = \sup\{\delta_X(w) \mid w \in F\}$. The *degree* of X is $d(X) = d_{A^*}(X)$. For F recurrent containing a bifix code X , one has $d_F(X)$ finite if and only if X is F -thin and an F -maximal bifix code [4].

The notion of parse was generalized to pseudowords in [10], and so we may extend δ_X to $\widehat{A^*}$. Since then, we obtained the following useful tool.

Proposition 3.1. *Consider a factorial set F of A^* . Let X be a rational subset of F with finite F -degree d . Then $\delta_X(w) \leq d$ for every $w \in \overline{F}$, and the mapping $\delta_X: \overline{F} \rightarrow \{1, \dots, d\}$ thus defined is continuous, if we endow $\{1, \dots, d\}$ with the discrete topology.*

A pseudoword u is *forbidden* in $Y \subseteq \widehat{A^*}$ if u is not a factor of an element of Y . The next proposition was deduced with the help of Proposition 3.1. We explain the notation used there. Let L be a rational language of A^* . By the universal property of $\widehat{A^*}$, the syntactic homomorphism $\eta_L: A^* \rightarrow M(L)$ admits a unique extension to a continuous homomorphism $\hat{\eta}_L: \widehat{A^*} \rightarrow M(L)$.

Proposition 3.2. *Let Z be a rational maximal bifix code of A^* . Suppose that F is a recurrent subset of A^* and that the intersection $X = Z \cap F$ is rational. The equality $d_F(X) = d(Z)$ holds if and only if the elements of $J(F)$ are forbidden in \overline{Z} . Moreover, if $d_F(X) = d(Z)$ then $\hat{\eta}_{Z^*}(J(F)) \subseteq J(Z)$.*

Consider a language L of A^* . Let $u, v \in A^*$. By definition, $\eta_L(u) \leq \eta_L(v)$ if and only if the context of u is contained in the context of v .

Proposition 3.3. *Let Z and F be subsets of A^* , with F factorial, and let $X = Z \cap F$. Suppose Z^* and X^* are rational. Let $e, f \in \widehat{A^*} \setminus \{1\}$ be idempotents, and let $u, v \in \widehat{A^*}$ with $u = euf$, $v = evf$ and $u \in \overline{F}$. Then:*

- (1) $\hat{\eta}_{X^*}(u) \leq \hat{\eta}_{X^*}(v) \Rightarrow \hat{\eta}_{Z^*}(u) \leq \hat{\eta}_{Z^*}(v)$, if e and f are forbidden in \overline{Z} ;
- (2) $\hat{\eta}_{Z^*}(v) \leq \hat{\eta}_{Z^*}(u) \Rightarrow \hat{\eta}_{X^*}(v) \leq \hat{\eta}_{X^*}(u)$, if e and f are forbidden in \overline{X} .

Applying the preceeding tools, we deduce relationships between the maximal subgroups of $J(F)$, $J(Z)$ and $J_F(Z \cap F)$, for suitable Z and F .

Theorem 3.4. *Let F be a factorial subset of A^* , take a rational prefix code Z of A^* , and suppose $X = Z \cap F$ is a rational F -maximal prefix code. Let H be a maximal subgroup of $\widehat{A^*}$ with $H \subseteq \overline{F}$ and the elements of H being forbidden in \overline{Z} . Consider the maximal subgroup H_X of $M(X^*)$ containing $\hat{\eta}_{X^*}(H)$ and the maximal subgroup H_Z of $M(Z^*)$ containing $\hat{\eta}_{Z^*}(H)$. There is an injective homomorphism $\alpha: H_X \rightarrow H_Z$ such that the diagram*

$$(3.1) \quad \begin{array}{ccc} H & \xrightarrow{\hat{\eta}_{Z^*}} & H_Z \\ \hat{\eta}_{X^*} \downarrow & \nearrow \alpha & \\ H_X & & \end{array}$$

commutes.

4. MAIN RESULTS

Based on the technical results of the previous section, namely Theorem 3.4, we deduce our main results.

Let F be a recurrent set of A^* . Say that a rational code Z of A^* is *F-charged* if $\hat{\eta}_{Z^*}$ maps maximal subgroups of $J(F)$ onto maximal subgroups of $J(Z)$. A rational code X contained in F is *weakly F-charged* if $\hat{\eta}_{X^*}$ maps maximal subgroups of $J(F)$ onto maximal subgroups of $J_F(X)$.

Theorem 4.1. *Consider a recurrent subset F of A^* and a rational bifix code Z of A^* with finite degree such that $X = Z \cap F$ is rational. Let H be a maximal subgroup of $J(F)$. The following conditions are equivalent:*

- (1) Z is F -charged;
- (2) $d_F(X) = d(Z)$, $G_F(X) \simeq G(Z)$ and X is weakly F -charged;
- (3) $d_F(X) = d(Z)$, $|G_F(X)| = |G(Z)|$ and X is weakly F -charged.

Recall that if $F \subseteq A^*$ is (uniformly) recurrent and Z is a maximal bifix code of A^* , then $Z \cap F$ is an F -maximal bifix (finite) code [4].

We show that a group code of A^* is F -charged when F is an uniformly recurrent *connected* set (that is, with only connected extension graphs, see [5]) with alphabet A . Therefore, we get the following corollary.

Corollary 4.2. *If Z is a group code of A^* and F is a uniformly recurrent connected set of alphabet A , then $d(Z) = d_F(Z \cap F)$ and $G(Z) \simeq G_F(Z \cap F)$.*

We say that a rational code Z is *nil-simple* if all idempotents of $M(Z^*)$ are in $J(Z)$. Group codes and finite codes are nil-simple. If F is uniformly recurrent and Z is nil-simple, the equality $d_F(X) = d(Z)$ in Theorem 4.1 becomes redundant, as seen next.

Theorem 4.3. *Let Z be a uniformly recurrent subset of A^* , and let Z be a nil-simple rational maximal bifix code Z of A^* . The following are equivalent:*

- (1) Z is F -charged;
- (2) $G_F(Z \cap F) \simeq G(Z)$ and $Z \cap F$ is weakly F -charged;
- (3) $|G_F(Z \cap F)| = |G(Z)|$ and $Z \cap F$ is weakly F -charged.

Moreover, if Z is F -charged, then $d(Z) = d_F(Z \cap F)$.

The special case of Corollary 4.2 in which Z is a group code and F is Sturmian was first proved in [4].

We also studied F -groups as permutation groups acting in a natural manner. In what follows, Q_Y is the set of vertices of the minimal automaton of Y^* , and i_Y is the corresponding initial state.

Theorem 4.4. *Let F be a recurrent subset of A^* . Suppose that Z is a rational bifix code of finite degree d . Let $X = Z \cap F$ and suppose that X is rational. Let H be a maximal subgroup of $J(F)$ such that $H_Z = \hat{\eta}_Z^*(H)$ is a maximal subgroup of $J(Z)$, and let $H_X = \hat{\eta}_X^*(H)$. Take the map $f: Q_X \cdot H_X \rightarrow Q_Z \cdot H_Z$ given by $f(i_X \cdot u) = i_Z \cdot u$, for $u \in H$, and take the unique group isomorphism $\alpha: H_X \rightarrow H_Z$ such that Diagram (3.1) commutes. Then the pair (f, α) is an equivalence of permutation groups with degree d .*

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Complexity of Robinson tiling

Abstract

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Raphael Robinson in his work on undecidability of the domino problem [1] introduced the set of six tiles depicted in Figure 1. The tiles can be rotated and reflected, one tile can fit to the other only in such a way that the arrow head matches arrow tail and each 2×2 block must contain exactly one *bumpy corner*, the leftmost in Figure 1.

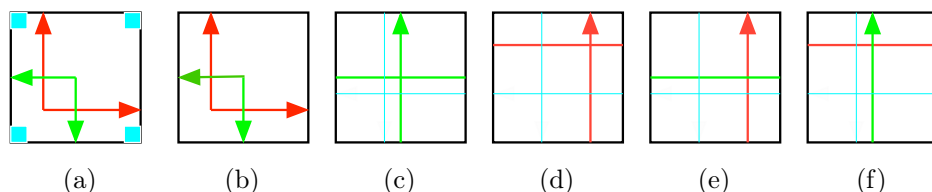


Figure 1: Tiles of type (a) are called *bumpy corners*, tiles of type (b) are called *corners*, all the other tiles are called *arms*.

It is possible to tile the Euclidean plane with copies of this six tiles, but only in an *aperiodic* way. The key to this result is that any Robinson tiling has a hierarchical structure (see Figure 2). For all n , it is possible to define the *supertiles* of rank n inductively. Bumpy corner tiles are said to be supertiles of the first rank, supertiles of second and third rank are shown in Figure 2. An increasing union of supertiles of rank n is called an infinite order supertile. The Robinson tiling can be either made of only *one* infinite order supertile or contain *two* or *four* infinite order supertiles.

We will prove that the number of distinct blocks of size $n \times n$ (with $n \geq 2$) that could appear in a Robinson tiling made of one infinite order supertile is defined by the formula

$$A(n) = 32n^2 + 72n \cdot 2^{\lfloor \log_2 n \rfloor} - 48 \cdot 2^{2\lfloor \log_2 n \rfloor}.$$

Similar method for counting the number of factors in the *paper folding sequence* was used by Jean-Paul Allouche in [2].

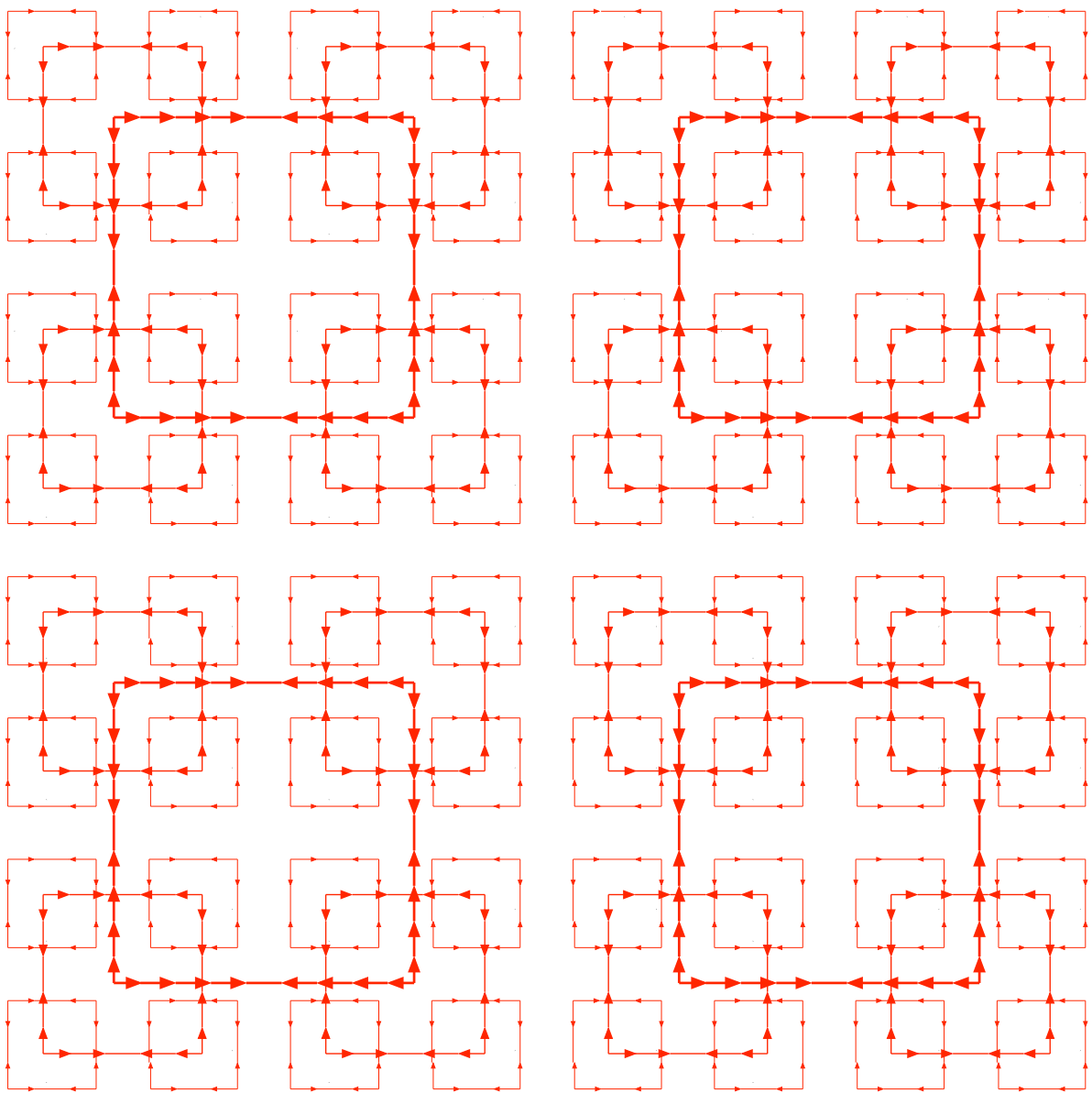


Figure 2: Hierarchy.

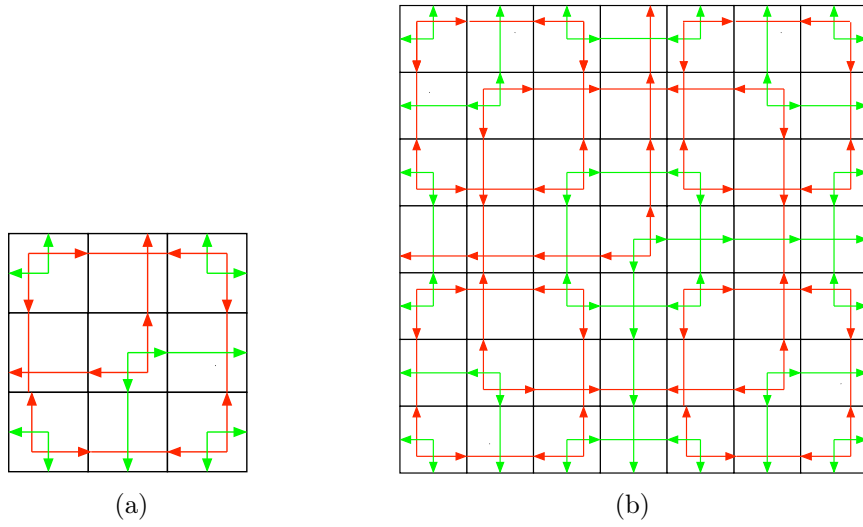


Figure 3: Supertiles of second and third rank.

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Abelian Anti-Powers in Infinite Words

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Abstract

We introduce and study the notion of an abelian anti-power in the context of combinatorics on words. An abelian anti-power of order k (or simply an abelian k -anti-power) is a concatenation of k consecutive words of the same length having pairwise distinct Parikh vectors. This definition therefore generalizes to the abelian setting the notion of a k -anti-power, as introduced in [G. Fici et al., *Anti-powers in infinite words*, J. Comb. Theory, Ser. A, 2018], that is a concatenation of k pairwise distinct words of the same length. In particular, we deal with the question to determine whether a word contains abelian k -anti-powers for arbitrarily large k . A word with bounded abelian complexity clearly cannot contain abelian anti-powers of arbitrary order. We show that the Sierpiński word (whose abelian complexity grows logarithmically) does not contain abelian 11-anti-powers. Another question is to find words with low factor complexity that contain both abelian powers and abelian anti-powers of arbitrary order. We show that all paperfolding words have this property.

1 Introduction

Many of the classical definitions in combinatorics on words (e.g., period, run, power, factor complexity, etc.) have a counterpart in the abelian setting, though they may not enjoy the same properties.

Recall that the Parikh vector $P(w)$ of a word w over a finite ordered alphabet $\mathbb{A} = \{a_1, a_2, \dots, a_{|\mathbb{A}|}\}$ is the vector whose i -th component is equal to the number of occurrences of the letter a_i in w , $1 \leq i \leq |\mathbb{A}|$. For example, the Parikh vector of $w = abbca$ over $\mathbb{A} = \{a, b, c\}$ is $P(w) = (2, 2, 1)$. This notion is at the basis of the abelian combinatorics on words, where two words are considered equivalent if and only if they have the same Parikh vector.

The fundamental result of Morse and Hedlund [4] (an infinite word is aperiodic if and only if its factor complexity is unbounded) does not hold anymore in the case of the abelian complexity (the function that counts the number of distinct Parikh vectors of factors of length n for each n), as there exist aperiodic words with bounded abelian complexity. In fact, Richomme et al. [5] have observed that if a word has bounded abelian complexity, then it contains abelian powers of any order — an abelian power of order k is a concatenation of k words having the same Parikh vector. However, this is not a characterization of words with bounded abelian complexity. Madill and Rampersad proved that the regular paperfolding word has unbounded abelian complexity [3], and Štěpán Holub proved that it contains abelian powers of any order [2].

In a recent paper [1], the first and the third author, together with Antonio Restivo and Luca Zamboni, introduced the notion of an anti-power. An *anti-power of order k* , or simply a *k -anti-power*, is a concatenation of k consecutive pairwise distinct words of the same length. E.g., *aabaaabbbaba* is a 4-anti-power.

In [1], it is proved that the existence of powers of any order or anti-powers of any order is an unavoidable regularity for infinite words:

Theorem 1. [1] Every infinite word contains powers of any order or anti-powers of any order.

Note that in the previous statement there is no hypothesis on the alphabet size.

In this paper, we extend the notion of an anti-power to the abelian setting.

Definition 2. An *abelian anti-power of order k* , or simply an *abelian k -anti-power*, is a concatenation of k consecutive words of the same length having pairwise distinct Parikh vectors.

For example, *aabaaabbbabb* is an abelian 4-anti-power. Notice that an abelian k -anti-power is a k -anti-power but the converse does not necessarily holds (which is dual to the fact that a k -power is an abelian k -power but the converse does not necessarily holds).

We think that an analogous of Theorem 1 may still hold in the case of abelian anti-powers, but unfortunately the proof of Theorem 1 does not seem to be generalizable to the abelian setting.

Problem 1. Does every infinite word contain abelian powers of any order or abelian anti-powers of any order?

Clearly, if a word has bounded abelian complexity, then it cannot contain abelian anti-powers of arbitrary order. However, we show in this paper that the converse is not true. Indeed, we prove that the Sierpiński word does not contain abelian 11-anti-powers. The Sierpiński word has logarithmic abelian complexity (by construction) and contains abelian powers of any order (since it contains arbitrarily long blocks of *bs*).

An infinite word can contain both abelian powers of any order and abelian anti-powers of any order. This is the case, for example, of any word with full factor complexity. However, finding a class of words with low factor complexity satisfying this property seems a more difficult task. Indeed, most of the well-known examples of aperiodic words (Thue-Morse, Sturmian words, etc.) have bounded abelian complexity, hence they cannot contain abelian anti-powers of any order — whereas, by the aforementioned remark of Richomme et al. [5], they contain abelian powers of any order. Building upon the theory that Štěpán Holub developed to prove that all paperfolding words contain abelian powers of any order [2], we prove that all paperfolding words contains also abelian anti-powers of any order.

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Inflation of digitally convex polyominoes

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1. Introduction

In the literature, many papers introduced and studied different convexity notions. For example, Kim and Rosenfeld in [1] investigated different notions of discrete convex sets, where a set in Euclidean geometry is convex if and only if for any pair of points p_1, p_2 in a region R , the line segment joining them is completely included in R . In discrete geometry on square grids, this notion refers to the digitally convex convexity. We recall that a *polyomino* is a finite 4-connected set of unit squares in the lattice \mathbb{Z}^2 . If P is a polyomino and if for all p_1, p_2 inside P and such that the discrete segment joining them is completely included in P then P is a digitally convex polyomino. Digitally convex polyominoes are also the discretization of convexes of \mathbb{R}^2 , except when this discretization is not 4-connected. It implies that the intersection of two such polyominoes is also digitally convex, as soon as it is 4-connected.

In this talk, we would like to study discrete geometrical constructions to deflate or inflate digitally convex polyominoes.

2. How to deflate a digitally convex polyomino?

We first study how to deflate a digitally convex polyomino P . This polyomino P can be decomposed into four paths. Indeed, a polyomino P is a finite set thus we define the minimal bounded box which is a rectangle such that P touches the four sides of the rectangle. Those paths start at the first unit square of intersection with each side of the rectangle. We denote the unit squares of intersection by W (the lowest unit square on the leftmost side), N (the leftmost unit square in the top side), E (the highest unit square on the rightmost side) and S (the rightmost unit square on the bottom side). The contour of a convex polyomino is then the union of the four (clockwise) paths WN , NE , ES and SW .

We use the result given by Brlek et al. [2] where they introduced a link between convexity and combinatorics on words by encoding the contour of the convex polyomino. Their result was based on *Christoffel* and *Lyndon words*. They considered each path of the convex polyomino and used the alphabet $A = \{0, 1, \bar{0}, \bar{1}\}$ to code the boundary of each polyomino where $0, 1, \bar{0}, \bar{1}$ encode $\rightarrow, \uparrow, \leftarrow, \downarrow$ respectively.

The main result in [2] states that a convex polyomino is characterized by the fact that the WN path admits a unique *Lyndon factorization* $\ell_1^{n_1} \dots \ell_k^{n_k}$ where the ℓ_i 's have decreasing slopes and they are primitive Christoffel words, and similar results for the three other paths.

It is not difficult to deflate the digitally convex polyomino. Indeed, we look at the set $\{s_1, s_2, \dots, s_{k+1}\}$ of unit squares (called corners) in the boundary of the polyomino such that each s_i are exactly the unit square corresponding to the extremities of each factor $\ell_j^{n_j}$ in the factorization given before (see Figure 1), or equivalently have a common unit square with the convex hull of the polyomino. Remove any square s_i , then the new polyomino is still digitally convex, and this process can be iterated.

This property has two consequences.

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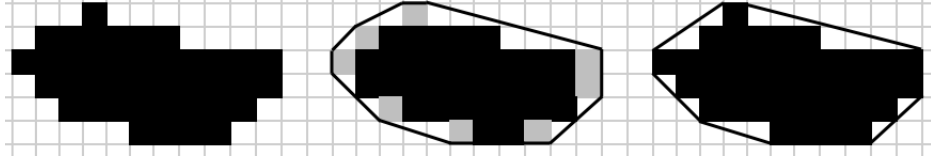


Figure 1: A digitally convex polyomino, its convex hull and its corners (grey unit squares) and a smaller digitally convex polyomino.

1. Given two digitally convex polyomino $C_1 \subset C_2$, C_2 can be deflated to C_1 in such a way that at each step one unit square only is canceled and such that at each step we have a digitally convex polyomino: as the polyomino C_2 is bigger than C_1 it contains a corner which is not in C_1 , then cancel this corner.
2. It is possible to inflate step by step from C_1 to C_2 , by the reverse process. However, it does not give a practical way to choose the unit square that we must add at each step.

3. How to inflate a digitally convex polyomino?

It is more difficult to get effective methods to inflate a convex polyomino C_1 to a convex polyomino C_2 where $C_1 \subset C_2$ with the constraint that we must add a single unit square at each step and maintain the digitally convex property at each step until reaching C_2 .

3.1. The spiral and strate constructions

First of all we take C_1 as a unit square anywhere inside C_2 .

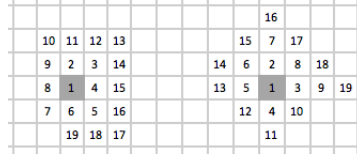


Figure 2: The spiral construction

We first investigate the spiral construction by adding a corner around the polyomino from a unit square polyomino by adding at each step a corner in a clockwise order (see Figure 2). This construction leads to an octogonal shape digitally convex polyomino. However, by keeping only those unit squares which are contained in C_2 (see Figure 3 left and center), we get a global construction of C_2 such that we maintain the convexity property at each step.

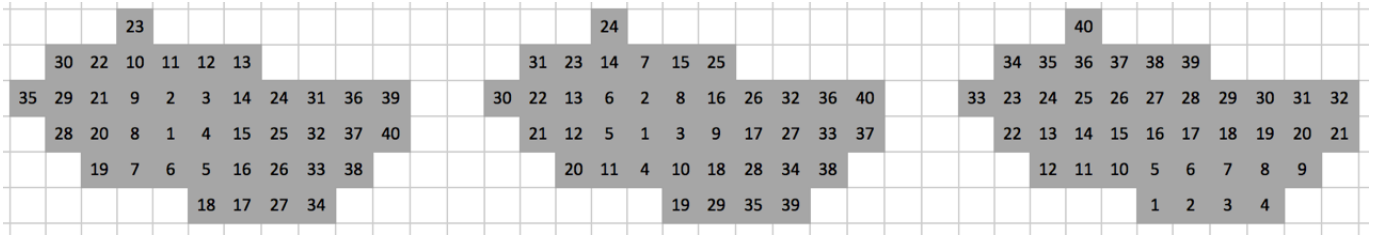


Figure 3: Two spiral and a strate construction of a digitally convex polyomino

In fact, we have many variants of this construction, by considering 4-connected spiral or 8-connected spiral or strate construction (see Figure 3). The strate construction consists in taking at first the lowest unit squares from the left to the right, and to continue to the second row in a correct order, and so on (see Figure 3, right). It has many variants.

If we try to inflate some digitally convex polyomino by the spiral method, it works only when we take special octogones (see Figure 4, left). In the general case, the convexity property disappears at some steps (see Figure 4, right, unit square number 8).

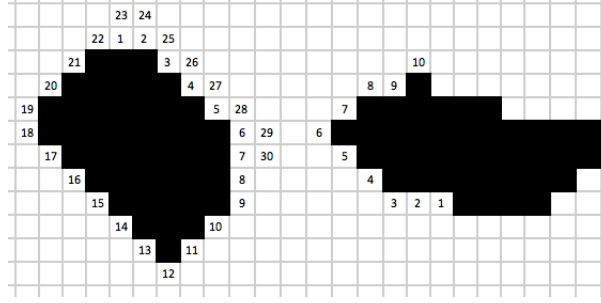


Figure 4: The spiral inflation

Thus now we must add unit squares using one construction of the paper of Dulio et al.[4], the authors introduces the *split* operator (based on the Borel Laubie standard factorization of Christoffel words [3]) in order to add unit square to the border of a digitally convex polyomino. By this operation the Christoffel property of words along the four paths is maintained and thus the difficult part of the inflation remains to maintain the monotony of the slopes of each Christoffel word along the four paths.

3.2. The split operator

From now on, we are working on the WN path. The split operator, as mentioned in [4], considers the furthest point of the Christoffel word ℓ_i with respect to its line segment and switches the factor 01 to 10. Using this operator and replacing the factors is equivalent to adding a unit square to the polyomino. We obtain in this case two new Christoffel words ℓ_i^+ and ℓ_i^- , where $\ell_i = \ell_i^- \ell_i^+$ is the standard factorization of the Christoffel word, and $\ell_i^- < \ell_i < \ell_i^+$ both for slopes and lexical order. The split operator consists in replacing ℓ_i by $\ell_i^+ \ell_i^-$ in the coding word of the path. Then the convexity property remains true in the simple case :

$$\ell_{i+1} \leq \ell_i^- < \ell_i^+ \leq \ell_{i-1}.$$

but these inequalities can fail.

So we might have different cases and situations, that we will present in the next section depending on the property of each unit square to add.

3.3. Adding one unit square:

We start by considering the case, where we will add one unit square to the convex $C_1(j)$ for a given j ($C_1(j)$ represents the j th step of the construction that is C_1 with j added unit squares). Which means that with this case we reach the step $C_1(j+1)$. In fact, three different cases can occur:

1. The first case, where we add a unit square on the Christoffel word ℓ_i with $n_i = 1$ of a given path using the split operator and no problems are faced. Which means the monotone order of slopes is maintained, i.e., as the split operator consists in replacing ℓ_i by $\ell_i^+ \ell_i^-$ in the coding word of the path:

$$\ell_{i+1} \leq \ell_i^- < \ell_i^+ \leq \ell_{i-1}.$$

and the convexity is conserved.

2. The second case also needs $n_i = 1$. It is if we add a unit square to ℓ_i of a given path we keep the Lyndon factorization property, but this factorization is not as in the first case

$$\ell_1^{n_1} \dots (\ell_{i-1}^{n_{i-1}} \ell_i^+) \ell_i^- \ell_{i+1}^{n_{i+1}} \dots \ell_k^{n_k}.$$

The convexity is conserved in this second case.

3. The third case that we can face, is if we add a unit square to ℓ_i of a given path we loose the Lyndon factorization property; which means the monotone order of slopes is no longer correct, and the new polyomino is not digitally convex.

However, we can get a characterization of the first case. We give this characterization in the simplest case, which corresponds to $n_i = 1$. Going back to the factorization $\ell_1^{n_1} \dots \ell_k^{n_k}$ of the WN path, cases 1 or 2 occur when the following conditions are satisfied.

1. $\ell_{i+1} \leq \ell_i^-$ or $\ell_{i+1} = \ell_i^{-k} \ell$ for some positive integer k and some Christoffel word ℓ ;
2. $\ell_{i-1} \geq \ell_i^+$ or $\ell_{i-1} = \ell \ell_i^{+k'}$ for some positive integer k' and some Christoffel word ℓ .

3.4. An inflation method

We consider as before two distinct digitally convex polyominoes $C_1 \subset C_2$. Then there always exists in each of the four paths describing C_1 at least one Christoffel word corresponding to the cases 1 or 2 before. More precisely, we can choose the longest Christoffel word of the path, among those which does not correspond to some side of C_2 . It gives an effective method to inflate C_1 to C_2 and conserving the convex property at each step.

Some other methods based on completely different considerations could be investigated. For example we can use a continuous and increasing deformation of the border of the convex hull of C_1 to the border of the convex hull of C_2 , such that this deformation contains at any time at most one integer point. An effective construction of such a deformation can be made, using simple arguments but encountering some technical difficulties.

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Séries \mathbb{Q} -Hadamard et \mathbb{Q} -automates circulaires

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On étudie ici les séries formelles sur un monoïde libre A^* finiment engendré avec coefficients dans un corps commutatif \mathbb{K} , avec \mathbb{K} égal à \mathbb{Q} , \mathbb{R} ou \mathbb{C} . Dans ce cadre, les séries rationnelles et les séries reconnaissables coïncident et le *produit d'Hadamard* (ou produit terme à terme, noté \odot) de deux séries rationnelles est rationnel [8].

On considère deux opérations naturelles liées au produit d'Hadamard. Soient s et t deux séries formelles dans $\mathbb{K}\langle\langle A^* \rangle\rangle$.

- si pour tout w dans A^* , $\langle t, w \rangle \neq 0$, alors le quotient d'Hadamard de s et t est défini par :

$$\frac{s}{t}^{\odot} = \sum_{w \in A^*} \frac{\langle s, w \rangle}{\langle t, w \rangle} w;$$

- si pour tout w dans A^* , $\langle s, w \rangle^* = \sum_{k \geq 0} \langle s, w \rangle^k$ existe, alors l'itération d'Hadamard de s est défini par :

$$s^{\oplus} = \sum_{w \in A^*} \langle s, w \rangle^* w.$$

Proposition 1 *L'ensemble des séries \mathbb{K} -rationnelles n'est clos ni par quotient d'Hadamard (cf. [7]), ni par itération d'Hadamard.*

On définit donc l'ensemble des séries d'Hadamard comme la plus petite famille de séries contenant les séries rationnelles et close pour ces opérations.

Théorème 2 *L'ensemble des séries \mathbb{K} -Hadamard sur A^* est définie de manière équivalente comme*

- a) *l'ensemble des quotients d'Hadamard de séries \mathbb{K} -rationnelles ;*
- b) *la clôture des séries \mathbb{K} -rationnelles par somme, produit d'Hadamard et itération d'Hadamard.*

Remarquons que l'ensemble des séries \mathbb{K} -Hadamard n'est pas clos pour les opérations rationnelles (il n'est même pas clos par produit de Cauchy).

On considère par ailleurs des *automates circulaires* à multiplicité dans un corps. Comme un automate pondéré classique, lors de tout calcul, un tel automate lit son entrée de gauche à droite en calculant une valeur comme produit des valeurs associées aux transitions empruntées ; mais en plus, arrivé à la fin de son entrée l'automate peut soit stopper s'il se trouve dans un état final, soit emprunter une transition spéciale qui remplace la tête de lecture au début du mot. La valeur associée à un mot est la somme (potentiellement infinie) des valeurs

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des calculs acceptant ce mot. Si tout mot a une valeur bien définie, l'automate est dit *valide* et son *comportement* est la série dans laquelle le coefficient de chaque mot est la valeur calculée par l'automate.

Ce modèle naturel d'automate – souvent considéré comme une restriction des automates bi-directionnels [3, 5] – permet, dans le cas non pondéré, de reconnaître l'intersection de deux langages avec une machine qui compte un nombre linéaire d'états par rapport aux automates d'entrée et, dans le cas pondéré, de réaliser avec la même complexité le produit d'Hadamard de deux séries (même dans le cas de poids non commutatifs).

Comme on ne peut pas décider si la série réalisée par un \mathbb{Q} -automate probabiliste a tous ses coefficients strictement inférieurs à $1/2$ [4], si la série réalisée par un \mathbb{Q} -automate a tous ses coefficients strictement inférieurs à 1, en faisant boucler un tel automate, on ne peut pas décider si l'automate circulaire obtenu est valide.

Proposition 3 *Il n'est pas décidable si un \mathbb{Q} -automate circulaire est valide.*

Ceci est à mettre en lien avec le fait qu'on ne peut pas décider si le support d'une série \mathbb{Q} -rationnelle est A^* tout entier ; on ne peut donc pas décider si l'inverse d'Hadamard d'une telle série est défini.

Par contre, si les séries sont bien définies, on a l'équivalence suivante.

Théorème 4 *L'ensemble des séries réalisables par des \mathbb{K} -automates circulaires valides est exactement l'ensemble des séries \mathbb{K} -Hadamard.*

Corollaire 5 *L'équivalence des \mathbb{Q} -automates circulaires valides est décidable.*

De fait, le comportement s de tout \mathbb{Q} -automate circulaire peut être représenté par une paire de \mathbb{Q} -automates classiques réalisant des séries t_1 et t_2 telles que $s = \odot_{t_2}^{t_1}$. Savoir si deux comportements $s = \odot_{t_2}^{t_1}$ et $s' = \odot_{t'_2}^{t'_1}$ sont égaux revient alors à savoir si $t_1 \odot t'_2$ et $t'_1 \odot t_2$, qui sont des séries \mathbb{Q} -rationnelles, sont égales. Or, l'équivalence des séries \mathbb{Q} -rationnelles est décidable [1].

Comme nous l'avons dit plus haut, les automates circulaires peuvent être vus comme restriction des automates bi-directionnels (ou boustrophédons) ; ces derniers sont donc au moins aussi puissants. Dans le cas non pondéré, tous ces modèles sont équivalents aux automates classiques [9, 6], alors que dans le cas pondéré les automates bi-directionnels peuvent être strictement plus puissants que les automates circulaires, eux-mêmes plus puissants que les automates classiques (voir par exemple [2] pour le cas des transducteurs, équivalents à des automates sur le semi-anneau $\text{Rat}B^*$).

De manière assez surprenante, dans le cas des corps, les automates bi-directionnels ne sont pas plus puissants que les automates circulaires.

Théorème 6 *Toute série réalisée par un \mathbb{K} -automate bi-directionnel peut être réalisée par un \mathbb{K} -automate circulaire.*

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A definition and counting of biperiodic recurrent configurations in the sandpile model on \mathbf{Z}^2

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Abstract

For the sandpile model on the usual two dimensional grid, we propose a weaker version of Dhar criterion to define recurrent configurations among stable biperiodic configurations. We check this new criterion via an algorithm which auto-stabilises to a canonical ultimately periodic behaviour independent of details in its not fully specified initialisation. This leads to ultimately periodic edge/vertex traversals similar to those of Cori-Le Borgne [2] in the case of finite graphs and then to a bijection with *some* cycle-rooted forests on the torus describing the period. A determinantal formula [5] counts all those forests and the refinement with some monodromy parameters allows to identify in some coefficients the number of recurrent configurations.

The Abelian sandpile model was introduced by physicists Bak, Tang and Wiesenfeld in [1] as a model of self-organized criticality. Given a simple, undirected graph $(V \cup \{s\}, E)$ where we distinguish s as the sink of the graph, we consider *configurations* in this model which are an assignments $\eta : V \mapsto \mathbf{Z}$ of some grains of sand on each vertex. We say that η is stable at $x \in V$ if $\eta(x) < \deg(x)$, and η is stable if it is stable at all $x \in V$. If η is unstable at x , then x is allowed to topple which means that the vertex x sends one grain along each incident edge. This toppling is said *legal*. A toppling is *forced* when it is not necessarily legal. Grains arriving at the sink are lost. Given a configuration η , we define a stabilization as a sequence of allowed topplings until a stable configuration is reached. The result of all stabilizations is unique due to commutations of topplings of unstable vertices and is noted $stab(\eta)$.

Let $P(\eta)$ be the result of a stabilization of $\eta + \mathbf{1}_{s\sim}$, that is η with an extra grain on each neighbour of s , which may be interpreted as a forced toppling of the sink. $P(\eta)$ is also called the Dhar criterion since the set of recurrent configurations is a subset of the stable configurations characterized by Dhar [4] as the fixed points of P . For such a fixed point, each vertex topples exactly once in this process. The notion of recurrence is related to a natural Markov chain in this model not discussed here [3], and it is well studied for its connection with spanning trees [4], uniform spanning tree, the Tutte polynomial on the underlying graph [2]. Also on finite graphs, the set of recurrent configurations equipped with the operation $(\eta, \mu) \mapsto stab(\eta + \mu)$ is an abelian group [3]. When the graph is the grid \mathbf{Z}^2 , the existence of such a group is open.

One of our motivations is the search of finite groups on a subset of recurrent configurations on \mathbf{Z}^2 , which may be subgroups of the hypothetical (infinite) group. We focus on the subset of biperiodic configurations on the grid defined as follows. Let $\vec{P}_1, \vec{P}_2 \in \mathbf{Z}^2$ two non collinear vectors. A configuration η of \mathbf{Z}^2 is biperiodic of period (\vec{P}_1, \vec{P}_2) if for all $x \in \mathbf{Z}^2$, $\eta(x + \vec{P}_1) = \eta(x + \vec{P}_2) = \eta(x)$.

Several approaches are suggested by litterature to define the notion of recurrence: for example, adding one edge per period to an extra vertex called the sink (dissipative sandpiles [6]), another example merges in one sink vertex all vertices outside a finite polygon and then scale this polygon [7]. Our approach relies on a weaker form of Dhar criterion and leads to a degenerate polygon which is an half-plane. Indeed we place the sink at infinity in a direction by analogy to the projective plane. The sink s is a point at infinity and will be describe by an euclidean vector $\vec{s} = (s_x, s_y) \in \mathbf{Z}^2$ where $\gcd(s_x, s_y) = 1$. Note that \vec{s} and $-\vec{s}$ refer to different sinks.

The sink being sent to infinity, the difficulty of toppling it appears. We replace this by a forced toppling of an half-plane of line boundary orthogonal to the sink \vec{s} .

Definition 1 (Weak Dhar criterion in a rational direction). *A configuration η is said recurrent in the direction of the sink \vec{s} if and only if for any $k \in \mathbf{Z}$ the forced toppling of the vertices of the half-plane $\{(x, y) \in \mathbf{Z}^2 \mid s_x x + s_y y \geq k\}$ leads to the legal toppling of all other vertices.*

Proposition 1. *There exists execution of the weak Dhar criterion on biperiodic configurations that is auto-stabilizing to an ultimately periodical behaviour which does not depend on the position of the half-plane defined by a line colinear to \vec{s}^\perp and can be simulated in finite time.*

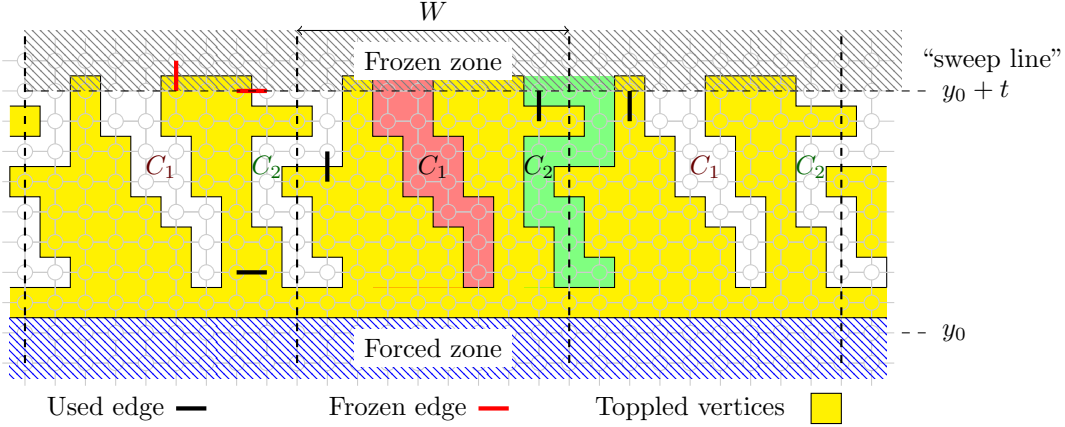


Figure 1: Weak Dhar criterion after step t , C_1 and C_2 are next connected components

Sketch of the proof when $\vec{s} = (0, -1)$ (that can be generalized for all \vec{s}). Let η be a recurrent configuration in direction $(0, -1)$ of period (\vec{P}_1, \vec{P}_2) . Without losing generality, we can assume that $\vec{P}_1 = (W, 0)$ and $\vec{P}_2 = (0, H)$ where $W, H > 0$. We equip the set of edges with the following order $\prec_{\vec{s}}$. Let e_1 (resp. e_2) be an edge of middle m_1 (resp. m_2) in the usual embedded of \mathbf{Z}^2 , $e_1 \prec_{\vec{s}} e_2$ if and only if $\vec{s} \cdot \vec{m}_1 < \vec{s} \cdot \vec{m}_2$ or $(\vec{s} \cdot \vec{m}_1 = \vec{s} \cdot \vec{m}_2 \text{ and } \vec{s}^\perp \cdot \vec{m}_1 < \vec{s}^\perp \cdot \vec{m}_2)$ where $\vec{s} \cdot \vec{m}_1$ is the usual scalar product between \vec{s} and \vec{m}_1 and \vec{s}^\perp is the vector $(-s_y, s_x)$. For $\prec_{(0, -1)}$, edges are ordered increasingly from top to bottom in priority, ties broken from left to right. When a vertex become unstable, it topples and its grains are sent along incident edges which become pending. Such a crossing grain is received at opposite endpoint when this pending edge is activated. We control the process of stabilisation by activating the maximal pending edge according to $\prec_{\vec{s}}$ (see [2] for details).

We start by a forced toppling of the half-plane $(x, y) \cdot (-\vec{s}) = y \leq y_0 \in \mathbf{Z}$, denoted $H_{\leq y_0}$. The remaining legal topplings, more precisely edges allowed to be activated, are enclosed between this half-plane and a “sweep line” $y = y_0 + t$, where $t \geq 0$ is called a step of execution for Dhar criterion. More precisely at each step $t > 0$, we only consider edges which middle has ordinate $y \in [y_0, y_0 + t]$. Thus there is a forced toppled zone $y \leq y_0$, a frozen zone $y > y_0 + t$ and a working zone $y_0 < y \leq y_0 + t$, see Figure 1. Since η is recurrent, at step t there is at least one vertex v_t that topples on line $y = y_0 + t$. The order $\prec_{\vec{s}}$ guarantees that the set of connected components of untopped vertices in the working zone is periodic of period $(W, 0)$ (C_1 and C_2 on Figure 1) and that these components are finite as long as η is recurrent since enclosed between the sweep line

and the sequences of topplings leading to toppled vertices $(v_t + k\vec{P}_1)_{k \in \mathbf{Z}}$. Thus the toppling of a vertex v is independent of the toppling of each $v + k(W, 0)$ with $k \in \mathbf{Z}^*$. This observation allows to simulate the criterion on a cylinder $[1, W] \times \mathbf{Z}$ with a periodic configuration of period $(0, H)$ for some H .

The more we force topplings, the more we topple vertices. This and translation symmetry \vec{P}_2 implies that if a vertex v topples at step $t + H$, then $v - (0, H) = v - \vec{P}_2$ topples no later than step t . As a corollary, at most one of the vertices $(v + k\vec{P}_2)_{k \in \mathbf{Z}}$ can topple legally in any H consecutive steps. Thus for any H consecutive steps, there is at most WH vertices that topple.

We can show by contradiction that each vertex v topples at some step. For any vertex v , either all $(v + k\vec{P}_2)_{k \in \mathbf{Z}}$ topples or there exists k_v such that exactly $(v + k\vec{P}_2)_{k \geq k_v}$ do not topple. If not all vertices topples at some step, we can define $k_{WH} = \max_v k_v$ where v runs over the subset S of the vertices of $[1, W] \times [1, H]$ for which k_v is defined and $\bar{S} := [1, W] \times [1, H] \setminus S$ the complement subset. By definition, in the half-plane $H_{>y_0+k_{WH}H}$, the subsets $(S + k\vec{P}_2)_{k \geq k_{WH}}$ never topples and all other vertices in $(\bar{S} + k\vec{P}_2)_{k \geq k_{WH}}$ topples. From this, we also deduce the existence of a step $t_{W,H}$ such that no vertex in the half plane $H_{\leq y_0+k_{WH}H}$ can legally topple at steps $t \geq t_{W,H}$. For steps $t \geq t_{W,H}$, only vertices in $(\bar{S} + k\vec{P}_2)_{k \geq k_{WH}}$ periodically topples. This describes an (infinite) sequence of legal topplings toward a stable configuration where $(S + k\vec{P}_2)_{k \geq k_{WH}}$ did not topple, which is a contradiction with the recurrence of η (so $S = \emptyset$).

Hence, let T be the first step when all vertices of $[1, W] \times [1, H]$ has been toppled. Since there is at most WH vertices that topple at step T , the sequence of toppling that destabilizes the last untoppled vertex of $[1, W] \times [1, H]$ starts on line $y = y_0 + T$, is at most of length WH and ends to a line below $y = y_0 + H$ so $T \leq H + WH$. So the criterion is finite and effective.

Moreover, from steps $T + 1$ to $T + H$, we observe the ultimately periodic behaviour of the Dhar criterion: WH vertices topple, all having distinct copies in the fundamental domain $[1, W] \times [1, H]$. In addition, for $k \leq H$, the half-plane $H_{\leq y_0+k}$ has toppled at step T so forcing toppling of this half-plane, instead of $H_{\leq y_0}$, results after $T - k$ steps in the same set of toppled vertices. So the ultimately periodic behaviour of Dhar criterion starting either by $H_{\leq y_0}$ or $H_{\leq y_0+k}$ are the same. \square

The proof induces a bijection with a subset of cycle rooted spanning forests [5] on the toroidal grid: from step $T + 1$ to step $T + H$, we attach to each vertex and its repetitions the edge that destabilizes it. In order to respect the order $\prec_{\vec{s}}$ on the cylinder, it is enough to process the toppling in only one repetition of each connected component at a given step. The result is a cycle rooted spanning forest of the toroidal grid with non contractible cycles. These cycles correspond to infinite periodic branches in the plane whose slopes are not orthogonal to \vec{s} .

Theorem 1. *Let \vec{s} be a sink. The set of recurrent biperiodic configuration of pattern size $W \times H$ on \mathbf{Z}^2 is in bijection with the set of cycle rooted spanning forests of the toroidal grid $W \times H$ whose slope (a, b) is such that $a \cdot s_x + b \cdot s_y \neq 0$.*

We assume in the next part that $W, H \geq 2$. We call *NCRSF* a non-contractible cycle rooted spanning forests. We want to count the recurrent configurations in a direction \vec{s} . By definition of homology, the image of one copy of an (oriented) cycle of homology class $(i, j) \in \mathbf{Z}^2$ starting at (x, y) into the plane ends at $(x + jW, y + iH)$. Kenyon [5] gives the following determinantal formula $\sum_{\text{NCRSFs } \gamma} (2 - z^i w^j - z^{-i} w^{-j})^k = \det \Delta$ where (i, j) is the homology class of the cycles of γ and k their number, z is the monodromy of cycles with homology class $(1, 0)$, w is the monodromy of cycles with homology class $(0, 1)$, Δ is the Laplacian on the line bundle with connection 1 over all

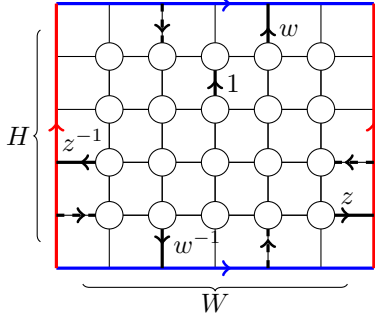


Figure 2: Connection Φ on the torus $W \times H$

| kj | ki | | | | |
|------|----------|----------|--------|------|---|
| | 0 | 1 | 2 | 3 | 4 |
| 0 | | 31300528 | 541732 | 1528 | 1 |
| 1 | 31300528 | 5427200 | 31232 | 4 | |
| 2 | 541732 | 31232 | 6 | | |
| 3 | 1528 | 4 | | | |
| 4 | 1 | | | | |

Figure 3: The number of NCRSF on a toroidal grid 4×4 in each direction where $k = \gcd(ki, kj)$

oriented edges except those crossing the blue (connection w or w^{-1}) and red sides (connection z or z^{-1}) of a fix fundamental rectangle as in figure 2. Then $\Delta(v) = \sum_{u \rightarrow v} v - \Phi_{u \rightarrow v} u$ where $\Phi_{u \rightarrow v}$ is the connection value on edge $u \rightarrow v$.

Due to planarity of the grid, a cycle cannot cross itself on the toroidal grid. Thus the homology class (i, j) of a cycle in this graph has $\gcd(i, j) = 1$. Moreover the length of such a cycle is at least $|Wj| + |Hi|$. That gives the first part of the following proposition.

Proposition 2. *Given γ a NCRSF on a toroidal grid $W \times H$, if γ has k cycles with homology class (i, j) then i and j are co-prime and $|kjW| + |kiH| \leq WH$. Reciprocally for any $(i, j) \in \mathbf{Z}^2$ co-prime and any $k > 0$ such that $|kjW| + |kiH| \leq WH$, there exists a NCRSF of parameters (i, j, k) .*

The second part of this proposition is achieved by k repetitions of a digital line from Bresenham's line algorithm (with corners).

We denote $Q_{i,j,k} = (2 - z^i w^j - z^{-i} w^{-j})^k$ for any $k > 0$ and any $(i, j) \in S = \{(0, 1)\} \cup \{(a, b) \mid a > 0 \text{ and } \gcd(a, b) = 1\}$, the $(Q_{i,j,k})_{i,j,k}$ are linearly independent. So the number of NCRSFs is the sum of the coefficients $(\alpha_{i,j,k})_{i,j,k}$ of $\det \Delta$ in the decomposition in $(Q_{i,j,k})_{i,j,k}$: $\det \Delta = \sum_{i,j,k} \alpha_{i,j,k} Q_{i,j,k}$. One can show that $\alpha_{i,j,k} = \alpha_{i,-j,k}$ for $i > 0$.

Proposition 3. *The number of biperiodic recurrent configurations in direction \vec{s} of size $W \times H$ is $\sum_{(i,j,k) \in S_{\vec{s}}} \alpha_{i,j,k}$ where $S_{\vec{s}} = \{(i, j, k) \mid (i, j) \in S, k > 0, |kjW| + |kiH| \leq WH, iW s_x + jH s_y \neq 0\}$.*

This formula enhances the counting the recurrent configurations in a direction that was limited to enumeration. Some explicit results are given on Figure 3. Some extra results up to $m \times n$ with $m, n \leq 9$ can be found at https://www.labri.fr/perso/hderycke/biperiodic_recurrent/.

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RETURN WORDS AND DERIVATED SEQUENCES TO ROTE SEQUENCES

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1. INTRODUCTION

We study complementary symmetric Rote sequences, which are sequences over the binary alphabet $\{0, 1\}$ with factor complexity $C(n) = 2n$ and with language closed on exchange of letters $0 \leftrightarrow 1$. We refer about the work in progress [7].

Rote in [8] proved that a sequence $\mathbf{v} = v_0v_1\cdots$ is complementary symmetric Rote sequence if and only if its first difference sequence $\mathbf{u} = u_0u_1\cdots$, which is defined by $u_i = v_i - v_{i+1} \pmod 2$, is Sturmian sequence.

Our aim is to describe return words and derivated sequences of complementary symmetric Rote sequences. In the sequel, we will call them simply Rote sequences. The study is based on the link between Rote and Sturmian sequences and recent work about derivated sequences of Sturmian sequences [1, 6].

Let $\mathbf{u} = u_0u_1u_2\cdots$ be an infinite sequence and let $w = u_iu_{i+1}\cdots u_{i+n-1}$ be its factor. The integer i is called an occurrence of the factor w . A return word to a factor w is a word $u_iu_{i+1}\cdots u_{j-1}$ with i and j being two consecutive occurrences of w such that $i < j$.

Let w be a prefix of \mathbf{u} which has k return words r_0, r_1, \dots, r_{k-1} . Then the sequence \mathbf{u} can be written as a concatenation of these return words: $\mathbf{u} = r_{d_0}r_{d_1}r_{d_2}\cdots$ and the derivated sequence of \mathbf{u} to prefix w is the sequence $\mathbf{d}_{\mathbf{u}}(w) = d_0d_1d_2\cdots$. For simplicity, we consider two derivated sequences to be the same if they differ only by a permutation of letters. We work only with sequences which are uniformly recurrent, i.e. each prefix w of \mathbf{u} occurs in \mathbf{u} infinitely many times and the set of all return words to w is finite.

Recall that the factor w of \mathbf{u} is right special, if both words $w0$ and $w1$ are factors of \mathbf{u} . Analogously the left special factor is defined. The factor is bispecial, if it is both left and right special. To find all derivated sequence it suffices to study only right special prefixes. Indeed, if the prefix w is not right special, then there is a unique letter a such that wa is a factor of \mathbf{u} . Thus the occurrences of w and wa in \mathbf{u} coincides and they have the same return words and derivated sequences. If \mathbf{u} is aperiodic, then w is always a prefix of some right special prefix of \mathbf{u} .

We focus on standard Sturmian sequences, i.e. the sequences whose each prefix is left special. In that case, we can take into consideration only bispecial prefixes to find all derivated sequences.

2. ROTE SEQUENCES AND ASSOCIATED STURMIAN SEQUENCES

We define the mapping \mathcal{S} which maps factors of Rote sequence to factors of associated Sturmian sequence.

Definition 1. *The mapping $\mathcal{S} : \{0, 1\}^+ \rightarrow \{0, 1\}^*$ is for every $v = v_0v_1\cdots v_n \in \{0, 1\}^+$ of length at least 2 defined by $\mathcal{S}(v_0v_1\cdots v_n) = u_0u_1\cdots u_{n-1}$, where $u_i = v_i + v_{i+1} \pmod 2$ for all $i \in \{0, 1, \dots, n-1\}$, $\mathcal{S}(v_0) = \varepsilon$.*

We can naturally extend the domain of \mathcal{S} to $\{0, 1\}^{\mathbb{N}}$ and write the associated Sturmian sequence \mathbf{u} to the Rote sequence \mathbf{v} as $\mathbf{u} = \mathcal{S}(\mathbf{v})$. To each Sturmian sequence, there are two associated Rote sequences \mathbf{v} and \mathbf{v}' . Since we have $\mathbf{v}' = E(\mathbf{v})$, we work only with Rote sequences starting with the letter 0 without lose of generality. The factors of Rote sequence \mathbf{v} and associated Sturmian sequence \mathbf{u} are closely related.

Proposition 2. *Let \mathbf{u} be a Sturmian sequence and \mathbf{v} be a Rote sequence such that $\mathbf{u} = \mathcal{S}(\mathbf{v})$. The word u is a factor of \mathbf{u} if and only if both words v, v' such that $u = \mathcal{S}(v) = \mathcal{S}(v')$ are the factors of \mathbf{v} . Moreover, for every $i \in \mathbb{N}$, i is an occurrence of u in \mathbf{u} if and only if i is an occurrence of v in \mathbf{v} or an occurrence of v' in \mathbf{v} .*

To study return words and derivated sequences of a given Rote sequence, we have to examine these objects in the case of associated Sturmian sequence. In [6] the authors describe derivated sequences of fixed points of Sturmian morphisms. Their basic idea is to suitably decompose a given Sturmian morphism onto some elementary morphisms.

A morphism is a mapping $\psi : \mathcal{A}^* \rightarrow \mathcal{B}^*$ such that $\psi(uv) = \psi(u)\psi(v)$ for all $u, v \in \mathcal{A}^*$. If $\mathcal{A} = \mathcal{B}$, ψ is a morphism over \mathcal{A}^* . The domain of the morphism ψ can be naturally extended to $\mathcal{A}^{\mathbb{N}}$. The matrix of a morphism ψ over \mathcal{A}^* is a matrix M_ψ defined by $(M_\psi)_{ab} = |\psi(a)|_b$ for all $a, b \in \mathcal{A}$. The morphism is primitive, if there is a positive integer k such that all elements of $(M_\psi)^k$ are positive.

A fixed point of a morphism ψ is a sequence \mathbf{u} such that $\psi(\mathbf{u}) = \mathbf{u}$. The sequence \mathbf{u} is substitutive if $\mathbf{u} = \sigma(\mathbf{v})$ for a morphism σ and a sequence \mathbf{v} which is a fixed point of morphism θ . Moreover, \mathbf{u} is substitutive primitive if θ is primitive. Durand in [3] proved that sequence \mathbf{u} is substitutive primitive if and only if \mathbf{u} has finite number of distinct derivated sequences.

A morphism ψ is Sturmian if $\psi(\mathbf{u})$ is Sturmian sequence for any Sturmian sequence \mathbf{u} . Consider the following elementary Sturmian morphisms φ_b and φ_β defined by

$$\varphi_b : \begin{array}{l} 0 \rightarrow 0 \\ 1 \rightarrow 01 \end{array} \quad \text{with } M_b = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \quad \text{and} \quad \varphi_\beta : \begin{array}{l} 0 \rightarrow 10 \\ 1 \rightarrow 1 \end{array} \quad \text{with } M_\beta = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}.$$

By [4] to a given standard Sturmian sequence \mathbf{u} we can uniquely assign the pair: directive sequence $\mathbf{z} \in \{b, \beta\}^{\mathbb{N}}$ and the sequence $(\mathbf{u}^{(n)})_{n \geq 0}$, where $\mathbf{u}^{(n)} \in \{0, 1\}^{\mathbb{N}}$ is a standard Sturmian sequence, such that for every $n \in \mathbb{N}$ we have

$$\mathbf{u} = \varphi_{z_0 z_1 \dots z_{n-1}}(\mathbf{u}^{(n)}).$$

The directive sequence \mathbf{z} contains infinitely many letters b and infinitely many letters β . In addition, \mathbf{z} is purely periodic, i.e. $\mathbf{z} = z^\infty$, if and only if \mathbf{u} is the fixed point of the morphism φ_z . This fixing morphism φ_z is always primitive.

3. RETURN WORDS TO PREFIXES OF ROTE SEQUENCES

Vuillon in [13] showed that an infinite sequence is Sturmian if and only if each non-empty factor has exactly two distinct return words. Using some results from [2] we show that all derivated sequences of Rote sequences to non-empty prefixes are over a ternary alphabet.

Theorem 3. *Let \mathbf{v} be a Rote sequence. Then every non-empty prefix x of \mathbf{v} has exactly three distinct return words.*

In addition, we can construct the return words in Rote sequence using the relevant return words in associated Sturmian sequence. We need an auxiliary definition of stability.

Definition 4. *The word $u = u_0 u_1 \dots u_{n-1} \in \{0, 1\}^*$ is called stable (S) if $\sum_{i=0}^{n-1} u_i = 0 \pmod{2}$. Otherwise, u is unstable (U).*

Definition 5. *Let w be a prefix of a Sturmian sequence \mathbf{u} with return words r, s . Due to Vuillon's result [9] \mathbf{u} is a concatenation of blocks $r^k s$ and $r^{k+1} s$ or blocks sr^k and sr^{k+1} for some positive integer k . With respect to the return words of w we distinguish three cases:*

- i) w is of type $SU(k)$, if r is stable and s is unstable;
- ii) w is of type $US(k)$, if r is unstable and s is stable;
- iii) w is of type $UU(k)$, if both r and s are unstable.

The type of prefix w is denoted \mathcal{T}_w .

We do not define the type SS since it cannot appear in Sturmian sequences. We use these prefix types to describe the return words to corresponding prefixes in Rote sequences.

Theorem 6. *Let x be a prefix of a Rote sequence \mathbf{v} . Denote by \mathbf{u} the Sturmian sequence $S(\mathbf{v})$ and by w its prefix $S(x)$ with return words r and s . Then the prefix x of \mathbf{v} has three return words $A, B, C \in \{0, 1\}^+$ such that:*

- i) $r = S(A0)$, $sr^{k+1}s = S(B0)$ and $sr^k s = S(C0)$ if w is of type $SU(k)$;

- ii) $rr = \mathcal{S}(A0)$, $rsr = \mathcal{S}(B0)$ and $s = \mathcal{S}(C0)$ if w is of type $US(k)$;
- iii) $rr = \mathcal{S}(A0)$, $rs = \mathcal{S}(B0)$ and $sr = \mathcal{S}(C0)$ if w is of type $UU(k)$.

Remark 7. We can also determine the type of w using the matrix P_w composed of the Parikh vectors of return words to w . Let w be a prefix of a Sturmian sequence \mathbf{u} with return words r, s , where r is the more frequent return word. Then the matrix P_w is defined by:

$$P_w = \begin{pmatrix} |r|_0 & |s|_0 \\ |r|_1 & |s|_1 \end{pmatrix} \pmod{2},$$

where $|u|_a$ denotes the number of letters a in the word u . Then the type \mathcal{T}_w of the prefix w is

$$\text{i) } SU \text{ if } P_w = \begin{pmatrix} p & q \\ 0 & 1 \end{pmatrix} \quad \text{ii) } US \text{ if } P_w = \begin{pmatrix} p & q \\ 1 & 0 \end{pmatrix} \quad \text{iii) } UU \text{ if } P_w = \begin{pmatrix} p & q \\ 1 & 1 \end{pmatrix}$$

for some numbers $p, q \in \{0, 1\}$.

4. DERIVATED SEQUENCES OF ROTE SEQUENCES

Theorem 8. Let \mathbf{v} be a Rote sequence with non-empty bispecial prefix x . Then the derived sequence $\mathbf{d}_{\mathbf{v}}(x)$ is uniquely determined by derived sequence $\mathbf{d}_{\mathbf{u}}(w)$ of $\mathbf{u} = \mathcal{S}(\mathbf{v})$ to the prefix $w = \mathcal{S}(x)$ and by type \mathcal{T}_w of prefix w .

Moreover, we are able to construct derived sequence $\mathbf{d}_{\mathbf{v}}(x)$ of a given Rote sequence \mathbf{v} to a prefix x if we know the type of $\mathcal{S}(x)$ and the derived sequences of $\mathcal{S}(\mathbf{v})$ to $\mathcal{S}(x)$.

First, we explain how to determine the type of a given prefix w of a standard Sturmian sequence \mathbf{u} . Clearly it suffices to study only bispecial prefixes. Let us order the bispecial prefixes of \mathbf{u} by their length starting from the shortest one. Their types can be determined using results from Section 3 of [6], where the authors explain how prefixes and their return words change under application of morphisms φ_b and φ_β .

Proposition 9. Let \mathbf{u} be a standard Sturmian sequence with directive sequence $\mathbf{z} \in \{b, \beta\}^{\mathbb{N}}$. Then the type of its n -th bispecial prefix w is given by the natural number k and the matrix P_w , where

- i) $P_w = M_{z_0} M_{z_1} \cdots M_{z_{n-1}} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \pmod{2}$ if the sequence $z_n z_{n+1} z_{n+2} \cdots$ has a prefix $b^k \beta$,
- ii) $P_w = M_{z_0} M_{z_1} \cdots M_{z_{n-1}} \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \pmod{2}$ if the sequence $z_n z_{n+1} z_{n+2} \cdots$ has a prefix $\beta^k b$.

The derived sequences of a given Sturmian sequence \mathbf{u} are described in detail in [6]. In particular, if \mathbf{u} is a standard Sturmian sequence with directive sequence $\mathbf{z} = z_0 z_1 z_2 \cdots$, then the derived sequence $\mathbf{d}_{\mathbf{u}}(w)$ to the n -th bispecial prefix w is standard Sturmian sequence with directive sequence $z_n z_{n+1} z_{n+2} \cdots$.

It is well known that all Sturmian sequences are 2iet sequences. We show that derived sequences of Rote sequences are 3iet sequences. A three interval exchange transformation $T : [0, 1) \rightarrow [0, 1)$ is given by partition of interval $[0, 1)$ into three subintervals $I_A = [0, \alpha)$, $I_B = [\alpha, \alpha + \beta)$ and $I_C = [\alpha + \beta, 1)$ and by permutation π on the set $\{1, 2, 3\}$ which expresses how these subintervals are rearranged, see [5] for more details. The 3iet sequence $\mathbf{u} = u_0 u_1 u_2 \cdots \in \{A, B, C\}^{\mathbb{N}}$ of transformation T with initial point $\rho \in [0, 1)$ is defined by $u_n = L$ if $T^n(\rho) \in I_L$ for all $n \in \mathbb{N}$.

Proposition 10. Let \mathbf{v} be a Rote sequence with non-empty prefix x , let $\mathbf{u} = \mathcal{S}(\mathbf{v})$ be a standard Sturmian sequence with prefix $w = \mathcal{S}(x)$. If $\mathbf{d}_{\mathbf{u}}(w)$ is 2iet sequence of transformation with intervals $[0, \alpha)$ and $[\alpha, 1)$ and initial point ρ , then the derived sequence $\mathbf{d}_{\mathbf{v}}(x)$ is 3iet sequence of transformation T with initial point ρ , where T is given by:

- i) intervals $[0, \alpha)$, $[\alpha, 2\alpha - k(1 - \alpha))$, $[2\alpha - k(1 - \alpha), 1)$ and permutation $(3, 2, 1)$ if w is of type $SU(k)$;
- ii) intervals $[0, 2\alpha - 1)$, $[2\alpha - 1, \alpha)$, $[\alpha, 1)$ and permutation $(3, 2, 1)$ if w is of type $US(k)$;
- iii) intervals $[0, 2\alpha - 1)$, $[2\alpha - 1, \alpha)$, $[\alpha, 1)$ and permutation $(2, 3, 1)$ if w is of type $UU(k)$.

Now we suppose that directive sequence \mathbf{z} is purely periodic, i.e. \mathbf{u} is fixed point of morphism φ_z .

By [6], sequence \mathbf{u} has at most $|z|$ distinct derivated sequences. We prove that the associated Rote sequence \mathbf{v} has a finite number of distinct derivated sequences too and so by Durand's result from [3] it is substitutive primitive.

Theorem 11. *Let \mathbf{v} be a Rote sequence and let $\mathcal{S}(\mathbf{v}) = \mathbf{u}$ be a standard Sturmian sequence which is a fixed point of a morphism φ_z . Then*

- i) \mathbf{v} has at most $3|z|$ distinct derivated sequences, each of them is fixed by a morphism;
- ii) \mathbf{v} is substitutive primitive.

Moreover, using our results and Durand's construction from [3] we can construct the fixing morphisms of these derivated sequences algorithmically.

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Formal Intercepts of Sturmian words

Caius Wojcik

Abstract :

We study Sturmian words, and particularly the second parameter describing this class of words.

Sturmian words are infinite words over a 2-letter alphabet. They are defined as the infinite words having lowest unbounded complexity. Namely, a theorem of Morse and Hedlund states that an infinite word is ultimately periodic if and only if it has bounded complexity. Sturmian words are characterised by the properties of not being ultimately periodic, and being balanced, that is, the number of 1 appearing in factors of a given length only takes two values.

The first parameter describing a Sturmian word is its slope, defined as the asymptotic proportion of 1, and characterises the set of factors of the sturmian word. Through the continued fraction expansion of this irrational number one can construct a distinguished Sturmian word of the corresponding slope, called the characteristic word, which is the only sturmian word of this slope admitting two sturmian extensions on the left.

Sturmian words are obtained geometrically by drawing a line on the plane and coding the crossing through a vertical line by a 1 and the horizontal line by a 0. The slope of the sturmian word being obtained through the slope of the considered line, and the second parameter, the intercept, is usually presented as the real number on the y -axis intersecting the line.

The characteristic word may be described using so-called standard and central words. The first ones are obtained by a concatenation process similar to Euclidean's algorithm on integers. The second ones are obtained by removing the last two letters of a standard word, and are palindromic words with

two relatively prime periods, realizing the sharp bound of Fine and Wilf's theorem.

On the other hand, we study the repetition function of sturmian words, defined as the number of distinct factors of a given length appearing at the beginning of an infinite word. This function is related to the diophantine exponent of infinite words, and in the case of sturmian words, it has been shown that the diophantine exponent is bounded if and only if the partial quotients of the continued fraction expansion of the slope are bounded.

We study the repetition function using so-called Rauzy graphs of infinite words, also called factor graphs. Rauzy graphs are defined as a sequence of directed graphs, with vertexes the factors of a fixed infinite word and arrows linking two factors if one is obtained by the other by a one-letter shift. The Rauzy graphs of arbitrary words are in general difficult to compute, but in the case of sturmian words they are particularly simple.

Indeed, the rauzy graphs of sturmian words are constituted of two cycles, patched together by a common part. The length of these two cycles are linked to the denominator of the sequence of convergents in the continued fraction expansion of the slope. Also, the evolution of the Rauzy graphs in the case of sturmian word may be fully described.

Any infinite word defines an infinite path on its Rauzy graph. Since these paths must be coherent with each others, there are a lot of restrictions on the possible paths taken by a given word. The consideration of these paths for different element of the shift orbit of the base word is a natural point of view of combinatorics on words. The repetition function can be read on those path, since it is the length of the longest hamiltonian path at the beginning of the base word's path.

Among the two cycles of the Rauzy graph of a sturmian word there a distinguished one, that we call the referent cycle, defined as the one through whom the characteristic word first passes. Indeed, we can compute the number of times the characteristic word turns around a cycle, and take it as a reference for other sturmian words. We use this idea to compute the length of cycles in Rauzy graphs, using the fact that the first repeated factor at the beginning of the characteristic word is its prefix.

These description of Rauzy graphs give another understanding of the so-called three-gap theorem, stating that an arithmetic sequence on the circle divides it into arcs with at most three possible lengths, corresponding to the frequencies of factors in the sturmian word. Indeed, two factors belonging

to the same branch of the Rauzy graph will have same frequencies, and the frequencies of the factors of the common part will be the sum of the two other frequencies.

The dynamical point of view of sturmian words as coding of rotations on the circle gives an analytic description of the intercept. Namely, an intercept can be expressed as an infinite alternating sum of irrational number obtained by shifting the original sequence of partial quotients, pounded with integer coefficients satisfying Ostrowski conditions. However, computations difficulties aside, there remains the problem of combinatorially define the intercept, in addition to the non-injectivity of this representation.

Ostrowski conditions may be view as a condition of injectivity of the number system associated to the continuants of the continued fraction expansion. They describe the rule that are being applied when one wants to sum two integers written in Ostrowski expansion, an operation that is very far from being understood.

In the dynamical setting given by an infinite word, and given an element in its dynamical subshift, realized as an approximation by suffixes of the base word, one can naturally consider the sequence of natural integers encoding each of the corresponding suffixes. In the context of Sturmian words, this sequence consists of partial expansions of an infinite Ostrowski expansion, linking the chaos of dynamics to the rigidity of combinatorics.

This infinite Ostrowski expansion is obtained by reading the operation of shifting on the Rauzy graphs of Sturmian words. Namely, the partial sums are obtained by considering the case of a length that is the increment of the length of bispecials that are purely central (that is, coming from standard words that are not semi-standard).

This observation allows us to define what we call the *formal intercept* of a sturmian word. They are formally defined as the projective limit of the system of integers with a given maximal number of terms in their Ostrowski expansion. Our main result is that this definition fully describes in a combinatorial and dynamical point of view the set of Sturmian words, in a bijective way.

To prove this result, we use the particularly simple shape of Rauzy graphs of Sturmian words. We use the computations of lengths of cycles previously obtained to check the congruences relations to build a Sturmian words associated to a given formal intercept. For the converse, we characterize the formal intercept as an exact approximation of a given Sturmian word.

We will compute the formal intercept of the two sturmian extensions of the characteristic word. They have the particularity of having their support in even places and odd places respectively, the support being defined as the set of indices with non-zero Ostrowski coefficients.

We will present the content of a work in progress, that goes as follows.

The action of the shift on Sturmian words allows us to add 1 to a formal intercept, therefore defining the addition of a formal intercept with a natural number. We will say that two formal intercept are equivalent whenever they coincide up to addition by some integers. For Sturmian words that are not equivalent to the characteristic word, this reduces to equality of almost all Ostrowski coefficients.

On the other hand, given a formal intercept one can use product formulas on reversal of standard words to define a Sturmian word. However this process will not reach the sturmian words that are equivalent to the characteristic word without being one of its suffixes. This operation, presented as a product formula, is convenient when considering questions of computations.

The sturmian words obtained by this process have a formal intercept that is complementary of the base intercept. That is, for a sum law defined on the set of classes of equivalence of Sturmian words, this amounts to the consideration of the opposite of a given Sturmian word.

The sum law defined this way would give us an isomorphism between the set of equivalence classes of Sturmian words and the quotient of the real numbers by the subgroup generated by 1 and the slope of the sturmian word. We will present how certain formulas on continuants of continued fraction expansions of quadratic numbers can be interpreted as elements of torsion in the group of equivalence class of Sturmian words, giving an asymptotic realisation of the operation of division by a natural integer in Ostrowski expansion.

UN ENSEMBLE INFINI DÉCIDABLE DE POINTS FIXES SANS CUBE ADDITIF

DAMIE JAMET, FLORIAN LIETARD, AND THOMAS STOLL

RÉSUMÉ. Dans ce travail, nous exhibons une classe infinie de points fixes de morphismes pour lesquels le problème d'évitabilité des cubes additifs est décidable. Pour cela, nous étendons l'algorithme proposé par Cassaigne *et al.* [1] à l'ensemble des morphismes semblables (au sens des matrices d'incidence) au morphisme étudié dans leur article. Comme dans l'article de Cassaigne *et al.*, la preuve est en partie informatique : autrement dit, elle nécessite une étude exhaustive, par ordinateur, d'un grand nombre de cas.

Les premières études sur les problèmes d'évitement remontent aux travaux de A. Thue [2, 3]. On sait ainsi qu'il existe un mot infini sans carré sur un alphabet de trois lettres (un *carré* est un mot fini de la forme w_1w_2 où $w_1 = w_2$).

Dans [4], F.M. Dekking montre qu'il existe un mot infini sur 3 lettres sans cube abélien (un mot fini w_1w_2 sur l'alphabet Σ est un *carré abélien* si w_1 et w_2 sont de même longueur et si w_1 est l'image de w_2 par une permutation de Σ). Dans [5], V. Keränen a montré l'existence d'un mot infini sur 4 lettres sans carré abélien.

Dans [1], Cassaigne *et al.* montrent que le point fixe $w_0 = \lim_{n \rightarrow \infty} \varphi_0^n(0) = 031430\dots$ du morphisme $\varphi_0 : 0 \mapsto 03, 1 \mapsto 43, 3 \mapsto 1, 4 \mapsto 01$ est sans cube additif : un *cube additif* est un mot fini sur $\Sigma \subset \mathbb{N}$ de la forme $w_1w_2w_3$ où w_1, w_2 et w_3 sont trois mots finis de même longueur et :

$$\sum_{i_1 \in w_1} i_1 = \sum_{i_2 \in w_2} i_2 = \sum_{i_3 \in w_3} i_3.$$

Dans notre exposé nous montrerons dans un premier temps que l'algorithme proposé dans [1] s'étend de manière naturelle à l'ensemble des morphismes semblables (au sens des matrices d'incidence) au morphisme φ_0 . En particulier, nous exhibons ainsi une classe potentiellement infinie (à un facteur multiplicatif et à une translation des lettres près sur l'alphabet) de points fixes sur 4 lettres sans cube additif.

Bien que les morphismes considérés soient semblables deux à deux, il est intéressant de constater que leurs points fixes ne possèdent pas tous les mêmes propriétés additives. Par exemple, soit $w_1 = \lim_{n \rightarrow \infty} \varphi_1^n(6) = 602106\dots$ le point fixe du morphisme $\varphi_1 : 0 \mapsto 2, 1 \mapsto 62, 2 \mapsto 10, 6 \mapsto$

60 semblable au morphisme φ_0 . Le mot infini w_1 est l'image du mot infini w_0 par la transformation de $\{0, 1, 3, 4\}$ dans $\{0, 1, 2, 6\}$ comme suit : $0 \mapsto 6, 3 \mapsto 0, 1 \mapsto 2, 4 \mapsto 1$. Notre implémentation de l'algorithme de Cassaigne *et al.*¹ nous donne une preuve informatique du fait que w_1 est sans cube additif. Cependant, alors que nous conjecturons que tous les carrés additifs dans w_0 sont des carrés abéliens, nous remarquons que le mot w_1 possède des carrés additifs propres, autrement dit, non abéliens, comme par exemple

$$w_1 = 6021062260101 \underline{06026} \underline{22622} 602 \dots$$

Cette partie sera consacrée à une étude numérique détaillée de ce phénomène.

Nous terminerons notre exposé par une présentation des pistes envisagées pour nous attaquer aux problèmes ouverts suivants :

- (1) Existe-t-il un mot infini sur l'alphabet $\{0, 1, 2, 3\}$ sans cubes additifs ?
- (2) Quel est le plus petit morphisme sur 3 lettres dont le point fixe ne possède pas de cubes additifs ?
- (3) Existe-t-il un mot infini sur un alphabet fini sans carrés additifs ?

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Autour du déséquilibre des mots C-adiques

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Résumé

Nous étudions une propriété combinatoire, le déséquilibre, d'une classe particulière de mots sur l'alphabet $\{a, b, c\}$: les mots C-adiques. En particulier, nous exhibons des familles de mots C-adiques de déséquilibres arbitrairement grands, et même des mots C-adiques de déséquilibre infini. Ces constructions ont été obtenues par l'exploration d'un automate et l'étude de ses chemins.

1 Motivations

À l'algorithme de fraction continue soustractif décrit par l'itération de l'application

$$\begin{array}{lll} (\mathbb{R}^+)^2 & \rightarrow & (\mathbb{R}^+)^2 \\ (x, y) & \mapsto & \begin{array}{ll} (x - y, y) & \text{si } x \geq y \\ (x, y - x) & \text{sinon} \end{array} \end{array}$$

est associée une classe particulière de mots infinis binaires : les mots sturmiens. Ceux-ci jouissent de deux caractérisations combinatoires : d'une part, ce sont exactement les mots de complexité $n + 1$, c'est-à-dire les mots qui admettent $n + 1$ facteurs de longueur n pour tout entier n ; d'autre part, ce sont les mots apériodiques dont le déséquilibre vaut 1, c'est-à-dire les mots apériodiques dans lesquels chaque lettre apparaît, à une unité près, un même nombre de fois dans tous les facteurs d'une longueur donnée.

Plusieurs tentatives ont été faites pour généraliser les fractions continues à des triplets de réels positifs. Un tel algorithme pourrait permettre d'approcher simultanément deux réels par une suite de couples de nombres rationnels.

Dans ce document, nous nous interrogeons sur les mots C-adiques, qui sont les mots ternaires associés à l'algorithme [3] :

$$\begin{array}{lll} (\mathbb{R}^+)^3 & \rightarrow & (\mathbb{R}^+)^3 \\ (x, y, z) & \mapsto & \begin{array}{ll} (x - z, z, y) & \text{si } x \geq z \\ (y, x, z - x) & \text{sinon.} \end{array} \end{array}$$

Tout comme les mots d'Arnoux-Rauzy, ces mots sont de complexité $2n+1$. L'intérêt de cet algorithme est qu'il admet comme instance n'importe quel triplet de réels positifs, contrairement à l'algorithme d'Arnoux-Rauzy qui n'est défini que sur un ensemble de mesure de Lebesgue nulle.

Aussi, il est naturel de s'interroger sur l'existence d'une borne uniforme pour le déséquilibre. Hélas, comme pour les mots d'Arnoux-Rauzy [2], nous pouvons construire des familles de mots C-adiques de déséquilibre aussi grand que souhaité, et même, par un lemme de pompage, des mots de déséquilibre infini.

2 Déséquilibre, mots C-adiques

Soit u un mot fini sur l'alphabet ternaire $A = \{a, b, c\}$ et $\alpha \in A$ une lettre. On désigne par $|u|_\alpha$ le nombre d'occurrences de la lettre α dans le mot u . Le *vecteur de Parikh* de u est le vecteur $\chi(u) = (|u|_a, |u|_b, |u|_c)$, qui compte les multiplicités de chacune des lettres de l'alphabet. Remarquons que la somme des coordonnées de ce vecteur est égale à la longueur du mot u , que l'on note $|u|$. Étant donné deux mots de même longueur u et v , on appelle *vecteur de déséquilibre de u et v* la différence de leurs vecteurs de Parikh. La somme des coordonnées d'un tel vecteur est nulle. Le déséquilibre d'un mot infini w est la quantité (éventuellement infinie) :

$$d = \sup_{n \in \mathbb{N}} \sup_{u, v \in F_n(w)} \|\chi(u) - \chi(v)\|_\infty,$$

qui s'écrit encore :

$$d = \sup_{n \in \mathbb{N}} \sup_{u, v \in F_n(w)} \max_{\alpha \in \{a, b, c\}} ||u|_\alpha - |v|_\alpha|.$$

Elle mesure les iniquités de répartition entre les lettres dans un mot donné.

Dans toute la suite, on s'intéressera aux substitutions $c_1 : a \mapsto a, b \mapsto ac, c \mapsto b$ et $c_2 : a \mapsto b, b \mapsto ac, c \mapsto c$, qui proviennent de l'algorithme de fraction continue 1 [3]. On notera aussi $C = \{c_1, c_2\}$.

Un *mot C-adique* est un mot infini de la forme $w = \lim_{n \rightarrow \infty} s_1 \circ \dots \circ s_n(w')$, où $(s_k)_{k \in \mathbb{N}} \in C^{\mathbb{N}}$ porte le nom de *suite directrice* de w , et où w' est un mot infini quelconque sur A ; avec la condition supplémentaire que chacune des substitutions c_1 et c_2 apparaisse une infinité de fois dans la suite directrice.

Remarque : pour S un ensemble de substitutions, on peut étudier les mots S-adiques dans un cadre plus général [1].

Enfin, pour $s \in C$, nous introduisons les applications ^-s et s^- qui, à un mot fini non vide u , associent le mot $s(u)$ auquel on efface la première (resp. la dernière) lettre. Ces applications ne sont pas des morphismes.

3 Construction de mots de C-adiques de déséquilibres arbitrairement grands

Pour tout entier n , nous souhaitons exhiber un mot C-adique de déséquilibre supérieur à n . Pour ce faire, nous allons construire par récurrence une suite $(w_n)_{n \in \mathbb{N}}$ de mots C-adiques, et donner sur chacun d'eux deux facteurs u_n et v_n de même longueur, dont le vecteur de déséquilibre est de norme n . Cela assure que w_n a pour déséquilibre au moins n . L'intuition de ces constructions provient de l'exploration d'un automate et de l'étude de ses chemins.

Construction de (w_n) .

Soit w_0 n'importe quel mot C-adique, par exemple $c_1 \circ c_2 \circ c_1 \circ c_2 \circ \dots(a)$. Posons $w_1 = c_2 \circ c_2 \circ c_2(w_0)$ et pour tout $n \geq 1$:

$$\begin{cases} w_{n+1} = c_1^{2n+2} \circ c_2(w_n) & \text{si } n \text{ est impair} \\ w_{n+1} = c_2^{2n+2} \circ c_1(w_n) & \text{sinon.} \end{cases}$$

Pour tout entier n , w_n est un mot C-adique.

Construction des suites de facteurs (u_n) et (v_n) .

La lettre b apparaît dans w_0 (la complexité nous le garantit), donc $acc = c_2 \circ c_2 \circ c_2(b)$ apparaît dans w_1 . Posons $u_1 = ac$, $v_1 = cc$ et, pour tout $n \geq 1$:

$$\begin{cases} u_{n+1} = c_1^- \circ (c_1 \circ c_1)^n \circ c_1 \circ c_2(u_n) \\ v_{n+1} = c_1 \circ (c_1 \circ c_1^-)^n \circ c_1 \circ c_2(v_n) \end{cases} \quad \text{si } n \text{ est impair}$$

$$\begin{cases} u_{n+1} = c_2 \circ (c_2 \circ c_2^-)^n \circ c_2 \circ c_1(u_n) \\ v_{n+1} = c_2^- \circ (c_2 \circ c_2)^n \circ c_2 \circ c_1(v_n) \end{cases} \quad \text{sinon.}$$

Les mots u_n et v_n sont bien facteurs de w_n , pour chaque n .

Lemme 1. *Soit $n \geq 1$.*

- *Si n est impair, alors u_n commence par a et termine par c , v_n commence et termine par c , et $\chi(u_n) - \chi(v_n) = (n, 1 - n, -1)$.*
- *Sinon, alors u_n commence et termine par a , v_n commence par a et termine par c , et $\chi(u_n) - \chi(v_n) = (1, n - 1, -n)$.*

Theorème 1. *Il existe des mots C-adiques de déséquilibre arbitrairement grand.*

4 Construction de mots C-adiques de déséquilibre infini

Nous venons de construire une famille de mots C-adiques pour laquelle le déséquilibre n'est pas borné. Toutefois, le déséquilibre des mots pris individuellement peut l'être. Pour construire un mot dont le déséquilibre est infini, nous allons recourir à un lemme de pompage.

La première étape consiste à montrer que lorsque l'on compose un mot C-adique par l'une ou l'autre des substitutions c_1 et c_2 , on ne rééquilibre pas le mot au-delà d'une certaine proportion.

Proposition 1. *Si w est un mot C-adique de déséquilibre $Des(w) \geq 3n$, alors $c_1(w)$ (resp. $c_2(w)$) est un mot C-adique de déséquilibre $Des(w) \geq n$.*

Pour tout entier naturel m , on note désormais $C^m := \{s_0 \circ \dots \circ s_{m-1} \in \{c_1, c_2\}^m\}$, et $C^* := \bigcup_{m \in \mathbb{N}} C^m$.

Corollaire 1 (Lemme de pompage). $\forall s \in C^*, \forall n \in \mathbb{N}, \exists \sigma \in C^*$ tel que pour tout mot C-adique w , $Des(s \circ \sigma(w)) \geq n$.

Theorème 2. *Il existe un mot C-adique de déséquilibre infini.*

Démonstration. D'après le corollaire 1, je peux construire une suite $(\sigma_k)_{k \in \mathbb{N}} \in (C^*)^{\mathbb{N}}$ telle que pour tout entier naturel n et pour tout mot C-adique w , le déséquilibre du mot $\sigma_0 \circ \dots \circ \sigma_n(w)$ vaut au moins n . Ainsi, le mot limite $w_\infty = \lim_{n \rightarrow \infty} \sigma_0 \circ \dots \circ \sigma_n(w_0)$ (s'il existe, sinon prendre une valeur d'adhérence de la suite), où w_0 est un mot C-adique quelconque, est lui même un mot C-adique (chaque σ_k fourni par le corollaire contient c_1 et c_2), et son déséquilibre vaut au moins n , pour tout n ; il est donc de déséquilibre infini. \square

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Number of valid decompositions of Fibonacci prefixes

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Abstract

We establish several recurrence relations and an explicit formula for the number of factorizations of the length- n prefix of the Fibonacci word into a (not strictly) decreasing sequence of standard Fibonacci words (OEIS sequence A300066).

1 Introduction

Extended Ostrowski numeration systems were introduced in [4] to solve a problem on palindromes in Sturmian words. A representation of n in such a system related to a given Sturmian slope corresponds to a factorization of the prefix of length n of the standard Sturmian word of this slope as a concatenation of finite standard words in a non-strictly decreasing order. Since in this abstract we consider only the Fibonacci case, it is reasonable to give a Fibonacci example: consider the prefix $abaababaabaaba$ of the Fibonacci word of length 14 and its decompositions to standard words $s_0 = a$, $s_1 = ab$, $s_2 = aba$, $s_3 = abaab$, $s_4 = abaababa$, $s_5 = abaababaabaab$ in a decreasing order. We see that

$$\begin{aligned} abaababaabaaba &= (abaababaabaab)(a) = s_5 s_0 \\ &= (abaababa)(abaab)(a) = s_4 s_3 s_0 \\ &= (abaababa)(aba)(ab)(a) = s_4 s_2 s_1 s_0 \\ &= (abaababa)(aba)(aba) = s_4 s_2 s_2 \\ &= (abaab)(aba)(aba)(aba) = s_3 s_2 s_2 s_2 \\ &= (abaab)(aba)(aba)(ab)(a) = s_3 s_2 s_2 s_1 s_0. \end{aligned}$$

These six factorizations correspond to six *valid* representations of 14:

$$14 = \overline{100001} = \overline{11001} = \overline{10111} = \overline{10200} = \overline{1300} = \overline{1211}.$$

If we restrict ourselves to representations corresponding to strictly decreasing sequences, or, which is the same, to the representations only containing zeros

and ones, their number for each n is equal to the well-studied OEIS sequence A000119 (see, e.g., [2]). In particular, the lower limit of the sequence is 1, and the upper asymptotics grows as $O(\sqrt{n})$. But here we consider the number of all valid representations of n , denoted by $T(n)$, so that, for example, $T(14) = 6$, the obtained sequence is new and was just recently uploaded to the OEIS as A300066. Here we prove a series of recurrence relations and an explicit formula for it.

2 Result

Let φ denote the golden ratio, $\varphi = \frac{1+\sqrt{5}}{2}$. The Fibonacci word is a Sturmian word $s = s[1]s[2]\dots$ of the slope $1/(\varphi + 1) = 1/\varphi^2$ and of zero intercept, that is, for all n , we have

$$s[n] = \begin{cases} a, & \text{if } \{n/\varphi^2\} < 1 - 1/\varphi^2, \\ b, & \text{otherwise.} \end{cases} \quad (1)$$

Here $\{x\}$ denotes the fractional part of x . Another way to construct s is to consider it as a limit $s = \lim s_n$ of finite standard words

$$s_{-1} = b, s_0 = a, s_{n+1} = s_n s_{n-1} \text{ for all } n \geq 0. \quad (2)$$

We write $N = \overline{k_n \dots k_0}$ and call this representation of N *valid* if $k_i \geq 0$ for all i and $s(0..N] = s_n^{k_n} s_{n-1}^{k_{n-1}} \dots s_0^{k_0}$, where $s(0..N]$ is the prefix of length N of the Fibonacci word. The number of valid representations of N is denoted by $T(N)$.

Proposition 1. *If $s[n] = a$, all valid representations of n end with an even number of 0s. If $s[n] = b$, all of them end with an odd number of 0s.*

The main result of this abstract is the following

Theorem 1. *If $s[n] = a$, then $T(n) = \lceil n/\varphi^2 \rceil$, or, which is the same, $T(n)$ is equal to the number of occurrences of b to $s(0..n]$ plus one. If $s[n] = b$, then $T(n) = \lceil n/\varphi^3 \rceil$, or, which is the same, $T(n)$ is equal to the number of occurrences of aa to $s(0..n]$ plus one.*

The proof of the theorem is based on several recurrence relations on $T(n)$:

Proposition 2. *For all \bar{s} , $T(\bar{r}\bar{0}) \geq T(\bar{r})$. If $r = r'10^{2k}$ for some $k \geq 0$, then $T(\bar{r}\bar{0}) = T(\bar{r})$.*

Proposition 3. *For all $z \in \{0,1\}^*$ and for all $k \geq 1$,*

$$T(\overline{z10^{2k}}) = T(\overline{z10^{2k-2}}) + T(\overline{z(01)^k}).$$

Proposition 4. *For all $z \in \{0,1\}^*$ and for all $k \geq 1$,*

$$T(\overline{z10^k1}) = \begin{cases} T(\overline{z10^{k+1}}), & \text{if } k \text{ is odd,} \\ T(\overline{z10^k}) + T(\overline{z(01)^{k/2}}), & \text{if } k \text{ is even.} \end{cases}$$

Propositions 2 to 4 give a full list of recurrence relations sufficient to compute $T(n)$ for every $n > 1$, starting from $T(1) = 1$. In particular, as corollaries, we get simple formulas on the values of T on Fibonacci numbers and their predecessors: starting from $F_1 = 1$, $F_2 = 2$, $F_{n+2} = F_{n+1} + F_n$, we get

$$T(F_{2n-1}) = T(F_{2n}) = F_{2n-3} + 1$$

and

$$T(F_{2n} - 1) = T(F_{2n+1} - 1) = F_{2n-2}.$$

The same recurrence relations serve to prove Theorem 1.

3 Acknowledgement

We are deeply grateful to J. O. Shallit for computing the first values of the considered sequence and submitting it to the On-Line Encyclopedia of Integer Sequences.

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Repetition avoidance in products of factors

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Abstract

We consider a variation on a classical avoidance problem from combinatorics on words that has been introduced by Mousavi and Shallit at DLT 2013. Let $\mathbf{pexp}_i(w)$ be the supremum of the exponent over the products of i factors of the word w . The repetition threshold $\mathbf{RT}_i(k)$ is then the infimum of $\mathbf{pexp}_i(w)$ over all words $w \in \Sigma_k^\omega$. Mousavi and Shallit obtained that $\mathbf{RT}_i(2) = 2i$ and $\mathbf{RT}_2(3) = \frac{13}{4}$. We show that $\mathbf{RT}_i(3) = \frac{3i}{2} + \frac{1}{4}$ if i is even and $\mathbf{RT}_i(3) \geq \frac{3i}{2} + \frac{1}{6}$ if i is odd and $i \geq 3$.

Keywords: Words; Repetition avoidance.

1 Main results

Mousavi and Shallit [2] have considered two generalizations of the avoidance of fractional repetitions in infinite words. A word is circularly r^+ -power-free if it does not contain a factor pxs such that sp is a repetition of exponent strictly greater than r . Let $\Sigma_k = \{0, 1, \dots, k-1\}$. The smallest real number r such that w is r^+ -power-free is denoted by $\mathbf{cexp}(w)$. Let $\mathbf{RTC}(k)$ denote the

minimum of $\text{cexp}(w)$ over every $w \in \Sigma_k^\omega$. Similarly, $\text{pexp}_i(w)$ is the smallest real number r such that every product of i factors of w is r^+ -power-free word and $\text{RT}_i(k)$ is the minimum of $\text{pexp}_i(w)$ over every $w \in \Sigma_k^\omega$.

In this paper, we consider the ternary alphabet. We obtain bounds on $\text{RT}_i(3)$ which extend the result of Mousavi and Shallit that $\text{RT}_2(3) = \frac{13}{4}$.

Proposition 1. $\text{RT}_2(k) = \text{RTC}(k)$.

Proof. The language of words in Σ_k^* avoiding circular repetitions of exponent at least e (or strictly greater than e) is a factorial language. As it is well-known [1], if a factorial language is infinite, then it contains a uniformly recurrent word w . By Proposition 14 in [2], $\text{pexp}_2(w) = \text{cexp}(w)$. This implies that $\text{RT}_2(k) = \text{RTC}(k)$. \square

Proposition 2. *If i is even and $i \geq 2$, then $\text{RT}_i(3) \geq \frac{3i}{2} + \frac{1}{4}$.*

Proof. Mousavi and Shallit [2] have proved that $\text{RT}_2(3) = \frac{13}{4}$, which settles the case $i = 2$. We have double checked their computation of the lower bound $\text{RT}_2(3) \geq \frac{13}{4}$. Suppose that i is a fixed even integer and that w_3 is an infinite ternary word. The lower bound for $i = 2$ implies that there exists two factors u and v such that $uv = t^e$ with $e \geq \frac{13}{4}$. Thus, the prefix t^3 of uv is also a 2-terms product of factors of w_3 . So we can form the i -terms product $(t^3)^{i/2-1}uv$ which is a repetition of the form t^x with exponent $x = 3(\frac{i}{2} - 1) + e \geq 3(\frac{i}{2} - 1) + \frac{13}{4} = \frac{3i}{2} + \frac{1}{4}$. This is the desired lower bound. \square

Proposition 3. *If i is odd and $i \geq 3$, then $\text{RT}_i(3) \geq \frac{3i}{2} + \frac{1}{6}$.*

Proof. Suppose that $i \geq 3$ is a fixed odd integer, that is, $i = 2j + 1$. Suppose that w_3 is a recurrent ternary word such that the product of i factors of w_3 is never a repetition of exponent at least $\frac{3i}{2} + \frac{1}{6} = 3j + \frac{5}{3}$. First, w_3 is square-free since otherwise there would exist an i -terms product of exponent $2i$. Also, w_3 does not contain two factors u and v with the following properties:

- $uv = t^3$,
- $u = t^e$ with $e \geq \frac{5}{3}$.

Indeed, this would produce the i -terms product $(uv)^j u$ which is a repetition of the form t^x with exponent $x = 3j + e \geq 3j + \frac{5}{3}$.

So if a , b , and c are distinct letters, then w_3 does not contain both $u = abcab$ and $v = cabc$ and w_3 does not contain both $u = abcbabc$ and $v = babcb$. A computer check shows that no infinite ternary square-free word satisfies this property. This proves the desired lower bound. \square

Proposition 4. *If i is even and $i \geq 2$, then $\text{RT}_i(3) \leq \frac{3i}{2} + \frac{1}{4}$.*

Proof. Let i be any even integer at least 2. To prove this upper bound, it is sufficient to construct a ternary word w satisfying $\text{pexp}_i(w) \leq \frac{3i}{2} + \frac{1}{4}$. The ternary morphic word used in [2] to obtain $\text{RT}_2(3) \leq \frac{13}{4}$ seems to satisfy the property. However, it is easier for us to consider another construction. Let us show that the image of every $7/5^+$ -free word over Σ_4 by the following 45-uniform morphism satisfies $\text{pexp}_i \leq \frac{3i}{2} + \frac{1}{4}$.

$0 \mapsto 010201210212021012102010212012101202101210212$
 $1 \mapsto 010201210212012101202101210201021202101210212$
 $2 \mapsto 010201210120212012102120210121021201210120212$
 $3 \mapsto 010201210120210121021201210120212012102010212$

First, we check that such ternary images are $\left(\frac{202}{135}^+, 36\right)$ -free using the method in [3]. Since $\frac{202}{135} < \frac{3}{2}$, the period of every repetition formed from i pieces and with exponent at least $\frac{3i}{2}$ must be at most 35. Then we check exhaustively that the ternary images do not contain two factors u and v such that

- $uv = t^e$,
- $e > 3$,
- $9 \leq |t| \leq 35$.

Thus, the period of every repetition formed from i pieces and with exponent strictly greater than $\frac{3i}{2}$ must be at most 8. Finally, we check exhaustively that $\text{pexp}_i \leq \frac{3i}{2} + \frac{1}{4}$ by considering only i -terms products that are repetitions of period at most 8. \square

2 Concluding remarks

We conjecture that $\text{RT}_i(3) = \frac{3i}{2} + \frac{1}{6}$ for every odd $i \geq 3$, based on numerical evidence. We hope to get a suitable morphism and a proof of this case in the near future. Then the next step would be to consider the 4-letter alphabet. A quick computer check shows that $\text{RT}_i(4) \geq i + \frac{1}{2}$ for every $i \geq 2$ and we conjecture that this is best possible. However, a proof of an upper bound of the form $\text{RT}_i(4) \leq i + c$ cannot be similar to the proof of Proposition 4. That is because the multiplicative factor of i , which drops from $\frac{3}{2}$ when $k = 3$ to 1 when $k = 4$, forbids that the constructed word is a morphic image of a Dejean word.

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Permutations sortable by dequeues, two stacks in parallel and two stacks in series

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This talk is based on joint work with Tony Guttman.

In his seminal work *The Art of Computer Programming* [4], Knuth was the first to consider a number of classic data structures from the point of view of the permutations they could produce from the identity permutation, or equinumerously, the permutations which the data structure can sort. In the case of a single stack, there are two possible operations: I , which moves the next element from the input onto the stack, and O , which outputs the top element of the stack (see Figure 1, top left). A permutation $\{p_1, p_2, \dots, p_n\}$ of the numbers $\{1, 2, \dots, n\}$ is said to be stack-sortable if it is possible for the elements to be input in the order p_1, p_2, \dots, p_n and output in the order $1, 2, \dots, n$. Knuth noticed that the permutations sortable using a single stack are exactly those which avoid the pattern 231, and that these are enumerated by the Catalan numbers. Knuth posed the same question for a number of other data structures, in particular, he asked how many permutations of each length can be sorted using a double ended queue (deque), two stacks in parallel (tsip) or two stacks in series (tsis).

In [1], Albert and Bousquet-Mélou characterised the counting function $P(t)$ for tsip-sortable permutations in terms of a generating function $Q(u, a)$ for weighted quarter plane loops. This solves the enumeration problem for tsip-sortable permutations in the sense that P is characterised by a system of functional equations. Using a similar method we discovered a simple relationship between the counting function $D(t)$ for deque-sortable permutations and the counting function $P(t)$, thereby solving the enumeration problem for deque-sortable permutations in the same sense. The enumeration of tsis-sortable permutations appears to be more difficult, as we have so far only been able to compute the first 20 coefficients of the associated counting function $S(t)$.

It is convenient to consider these permutations from the perspective of sequences of moves which sort them. For example, A stack-sortable permutation is uniquely determined by the sequence of operations I, O which sorts it. This operation sequence is then a word over the alphabet $\{I, O\}$ which contains n occurrences of each letter, with the restriction that the number of I 's preceding any O is at least the number of O 's preceding that O . Using this classification, it is easy to see that these words, and hence stack-sortable permutations, are enumerated by the Catalan numbers.

The operation sequence for a deque-sortable or tsip-sortable permutation is a word over the alphabet $\{I_1, I_2, O_1, O_2\}$, while for a tsis-sortable permutation, the associated alphabet is $\{\rho, \lambda, \mu\}$ (see Figure 1). We call these operation sequences deque-words, tsip-words and tsis-words, respectively. In these cases the correspondence between operation sequences and sortable permutations is not bijective because, in general, multiple different operation sequence can sort the same permutation.

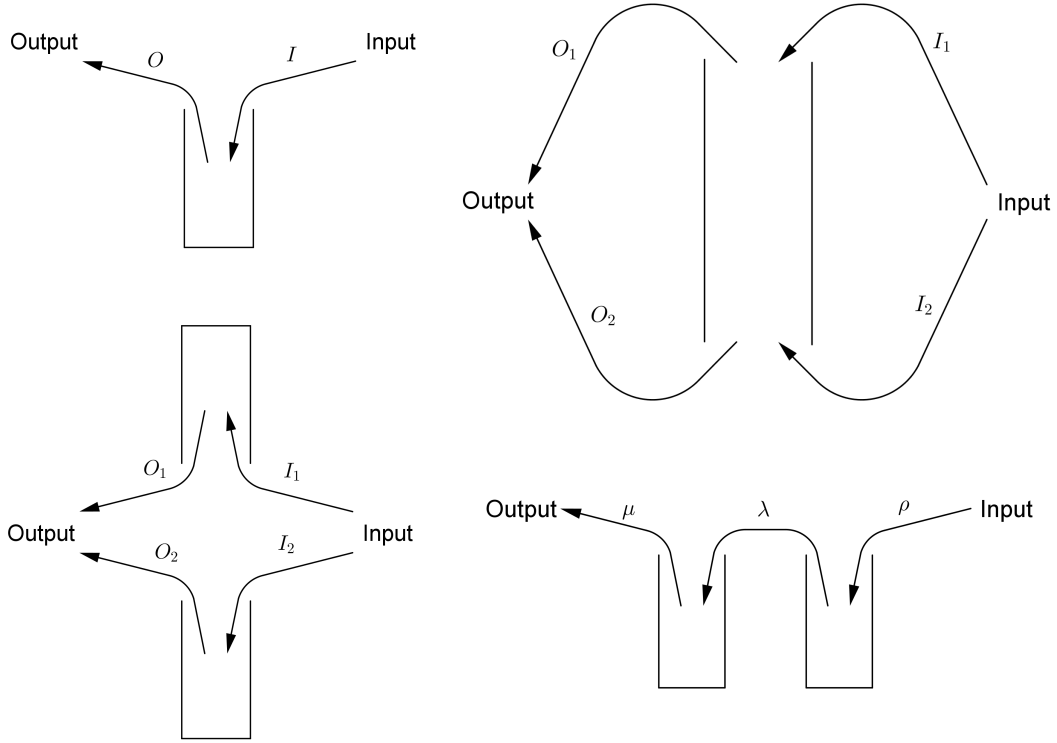


Figure 1: The allowed operations for a stack (top left), deque (top right), two stacks in parallel (bottom left) and two stacks in series (bottom right).

For two stacks in parallel, Albert and Bousquet-Mélou found a restricted class of operation sequences which are in bijection with their associated permutations. These *canonical* operation sequences are those with the following two properties:

1. The subwords I_1O_2 and I_2O_1 are forbidden
2. Any subword of the operation sequence which is itself a tsip-word begins with I_1 .

This classification allowed them to relate the generating function $P(t)$ for tsip-sortable permutations to a generating function $Q(u, a)$ for weighted quarter plane walks. Subsequently, we found a similar class of canonical operation sequences for deques, which are those deque-words with the above two properties as well as the following additional property: when the deque contains at most 1 element, only the moves I_1 and O_1 are allowed, not I_2 or O_2 .

By enumerating canonical deque-words, we derive the following relationship between the counting function $D(t)$ for deque-sortable permutations, and the counting function $P(t)$ for tsip-sortable permutations:

$$2D(t) = 2 + t + 2Pt - 2Pt^2 - t\sqrt{1 - 4P + 4P^2 - 8P^2t + 4P^2t^2 - 4Pt}.$$

Despite the fact that both problems are now in some sense solved, it is still not proven that the radii of convergence t_p of $P(t)$ and t_d of $D(t)$ are equal. From the expression above, we can deduce that if

$$1 - 4P + 4P^2 - 8P^2t + 4P^2t^2 - 4Pt > 0$$

for all $t \in [0, t_p)$, then $t_d = t_p$. It is conjectured, however, that the expression above is equal to 0 at $t = t_p$, so it is not possible to prove the inequality by simply approximating the left hand side. Since the function P has a simple characterisation in terms of the generating function $Q(u, a)$ for weighted quarter plane loops, the question of whether $t_d = t_p$ can be reduced to questions about Q . We show that the following three conjectures from [1] imply that $t_d = t_p$.

Conjecture 1. *The series $Q(u, a)$ is $(a + 1)$ -positive. That is, Q takes the form*

$$Q(u, a) = \sum_{n \geq 0} u^n P_n(a + 1),$$

where each polynomial P_n has positive coefficients.

Conjecture 2. *the radius of convergence $\rho_Q(a)$ of $Q(\cdot, a)$ is given by*

$$\rho_Q(a) = \begin{cases} \frac{1}{(2 + \sqrt{2 + 2a})^2}, & \text{if } a \geq -1/2, \\ \frac{-a}{2(a - 1)^2}, & \text{if } a \in [-1, -1/2]. \end{cases}$$

Conjecture 3. *The series $Q_u(u, a) = \frac{\partial Q}{\partial u}$ is convergent at $u = \rho_Q(a)$ for $a \geq -1/3$.*

Finally, using the solutions to these two problems, we have computed over 1000 coefficients of each of the series $P(t)$ and $D(t)$. We find that the coefficients of $P(t)$ behave as $\kappa_p \cdot \mu^n \cdot n^\gamma$, where $\mu = 8.281402207\dots$ and $\gamma \approx -2.473$, while the coefficients of $D(t)$ behave as $\kappa_d \cdot \mu^n \cdot n^{-3/2}$. In particular, we find that the growth rates of these two sequences agree to 10 significant digits, adding credence to the conjecture that $t_p = t_d$.

In the case of tsis-sortable permutations, we have computed the first 20 coefficients of the associated generating function $S(t)$. Our analysis suggests that the coefficients behave like $\kappa_s \cdot \mu_s^n \cdot n^{\gamma_s}$, where $\mu_s \approx 12.4$ and $\gamma_s \approx -2.5$.

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Rigidity for some dynamical systems of arithmetic origin

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The rigidity property for a measure-theoretic dynamical systems is the convergence to the identity of a sequence of powers of the map. We look at examples of rigid and non-rigid systems in the class of interval exchanges. Following those coming from square-tiled surfaces, which will be mentioned in P. Hubert's lecture, we consider the famous Veech example of 1969 and some generalizations, which are finite extensions of rotations of angle α with marked points β_i : by the same word-combinatorial methods as in those previous cases, we can prove they are rigid if α has unbounded partial quotients, non-rigid if the coding by the partition defined by the β_i is linearly recurrent. In the intermediate case when α has bounded partial quotients but the coding is not linearly recurrent, we have partial results using the Ostrowski expansions of the β_i related to α : there are rigid examples, including Veech 1969 in this case, and non-rigid ones providing the first known examples of non-rigid not linearly recurrent interval exchanges.

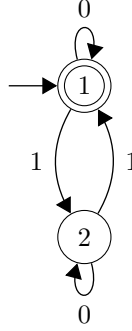
Etude de l'automate minimal des écritures en base 2^p des multiples de l'ensemble de Thue-Morse

Adeline MASSUIR

30 mai 2018

Ce travail a été réalisé en collaboration avec Emilie Charlier et Célia Cisternino.

L'ensemble de Thue-Morse, \mathcal{T} , est l'ensemble des nombres naturels dont l'écriture en base 2 contient un nombre pair de 1. Les premiers éléments sont 0, 3, 5, 6, 9, ... Cet ensemble est clairement 2-reconnaissable. En effet, l'automate suivant accepte les représentations en base 2 de ces nombres :



De plus, on peut facilement, à partir de cet automate, décrire l'automate qui accepte les représentations de ces nombres dans une base qui est une puissance de 2. Enfin, \mathcal{T} n'étant clairement pas ultimement périodique, on sait par le théorème de Cobham qu'il ne peut être b -reconnaissable pour un b qui n'est pas une puissance (entière positive) de 2.

Nous nous sommes donc intéressées à l'automate minimal de chacun des langages

$$\text{rep}_b(m\mathcal{T})$$

où $b = 2^p$ (avec $p > 0$) et $m \in \mathbb{N}_0$.

Pour ce faire, nous avons commencé par construire un automate qui acceptait ce langage. Plusieurs étapes sont nécessaires. Tout d'abord, on construit un automate qui accepte les couples $(\text{rep}_b(n), \text{rep}_b(mn))$ pour tous les naturels n (en prenant pour convention qu'on ajoute des 0 de tête au plus court des deux mots afin que les deux composantes aient la même longueur). Ensuite, on construit l'automate acceptant le langage

$$\{(\text{rep}_b(t), \text{rep}_b(n)) : t \in \mathcal{T}, n \in \mathbb{N}\}$$

(avec la même convention que précédemment). Ensuite, on effectue le produit des deux automates, pour obtenir un automate acceptant le langage suivant (toujours avec la même convention) :

$$\{(\text{rep}_b(t), \text{rep}_b(mt)) : t \in \mathcal{T}\}$$

Finalement, il nous reste à projeter le label de chaque arc sur sa deuxième composante.

A partir de cet automate, nous avons pu prouver que, si on note $b = 2^p$ et $m = k2^i$ où $p > 0$, k impair et $i \geq 0$, alors le nombre d'états de l'automate minimal du langage

$$\text{rep}_b(m\mathcal{T})$$

est $2k + \lceil \frac{i}{p} \rceil$.

Ce travail fait écho à l'article de B. Alexeev¹, dans lequel il détermine le nombre d'états de l'automate minimal de chaque langage $\text{rep}_b(m\mathbb{N})$ pour tous $m, b \in \mathbb{N}_0$. Toutefois, nos procédés de démonstration sont différents des siens.

1. Boris Alexeev, *Minimal DFA for testing divisibility*, Journal of Computer and System Sciences **69** (2004), 2, pp. 235–243.

Generalized Beatty sequences and complementary triples

Abstract

A Beatty sequence is the sequence $A(n) = \lfloor n\alpha \rfloor$ for $n \geq 1$, where α is a positive real number. What Beatty observed is that when B is the sequence $B(n) = \lfloor n\beta \rfloor$, with α and β satisfying

$$\frac{1}{\alpha} + \frac{1}{\beta} = 1, \quad (1)$$

then $(A(n))$ and $(B(n))$ are *complementary* sequences, that is, the sets $\{A(n) : n \geq 1\}$ and $\{B(n) : n \geq 1\}$ are disjoint and their union is the set of positive integers.

A *generalized Beatty sequence* is a sequence v defined by $v(n) = p\lfloor n\alpha \rfloor + qn + r$, where α, p, q, r are real numbers. These occur in several problems, as for instance in homomorphic embeddings of Sturmian languages in the integers ([1]).

Question 1 Let α be an irrational number, and let A defined by $A(n) = \lfloor n\alpha \rfloor$ for $n \geq 1$ be the Beatty sequence of α . Let Id defined by $\text{Id}(n) = n$ be the identity map on the integers. For which sextuples of integers p, q, r, s, t, u are the two sequences

$$v = pA + q\text{Id} + r \quad \text{and} \quad w = sA + t\text{Id} + u$$

complementary sequences?

Question 2 For which nonatuples of integers $(p_1, q_1, r_1, p_2, q_2, r_2, p_3, q_3, r_3)$ the three sequences

$$v_i = p_i A + q_i \text{Id} + r_i, \quad i = 1, 2, 3$$

are a complementary triple?

Here a *complementary triple* are three sequences, with the property that the sets they determine are disjoint with union the positive integers.

In this talk, based on joint work with Jean-Paul Allouche I give (incomplete) answers to these questions.

[1] Michel Dekking, The Frobenius problem for homomorphic embeddings of languages into the integers, Theoretical Computer Science 732(2018),73-79.

Rigidity and Substitutive Dendric Words

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Joint work with V. Berthé, F. Dolce, F. Durand and D. Perrin

Dendric words are infinite words defined in terms of extension graphs that describe the left and right extensions of their factors. Extension graphs are bipartite graphs that can be roughly described as follows: given an infinite word x , and given a finite factor w of x , one puts an edge between left and right copies of letters a and b such that awb is a factor of x . Dendric words are defined by requiring that the extension graph of each of its factor is a tree. This class of words with linear factor complexity includes classical families of words such as Sturmian words, codings of interval exchanges, or else, Arnoux-Rauzy words. Dendric words have striking combinatorial, ergodic and algebraic properties. This includes in particular algebraic properties of their return words [4], and of maximal bifix codes defined with respect to their languages [2, 5, 6]. They have been introduced in [4] and studied in several papers (as, for instance, [5, 6]), under the name of tree words. We have chosen to call them here dendric words, and the subshifts they generate dendric subshifts, in order to avoid any ambiguity with respect to the notion of tree shift that refers to shifts defined on trees, and not on words (see e.g. [1]).

In this talk, based on [7], I investigate the properties of substitutive dendric words and prove some rigidity properties. Rigidity of an infinite word x has to do with the algebraic properties of its *stabilizer* $\text{Stab}(x)$, that is the monoid of substitutions that fix it: an infinite word generated by a substitution is rigid if all the substitutions in $\text{Stab}(x)$ are powers of a unique substitution. In this work, we concentrate on the iterative stabilizer according to the terminology of [9]: we focus on non-erasing morphisms and on infinite words generated by iterating a substitution.

There are numerous results on the two-letter case concerning rigidity (see [10, 11] and also [3]). It is indeed well known that Sturmian words generated by substitutions are rigid [10, 11]. The situation is more contrasted as soon as the size of the alphabet increases. For instance, over a ternary alphabet, the stabilizer of a given infinite word can be infinitely generated, even when the word is generated by iterating an invertible primitive morphism (see [8, 9]).

Our main results are the following, where \S -adic expansions correspond to the limit of compositions of substitutions of the form $\sigma_1 \circ \dots \circ \sigma_n$, \S_e stands for

the set of elementary positive automorphisms of the free group generated by the alphabet of the word and tame substitutions are elements of \S_e^* .

Theorem 1. *A recurrent dendric word over an alphabet A is primitive substitutive if and only if it has an eventually periodic primitive \S_e -adic representation.*

Theorem 2. *Let x be a dendric word. Primitive substitutions in the stabilizer $\text{Stab}(x)$ of x coincide up to powers. More precisely, if x is a fixed point of both σ and τ primitive substitutions, then there exist $i, j \geq 1$ such that $\tau^i = \sigma^j$.*

Let x be a recurrent substitutive dendric word. There is a primitive tame substitution θ such that any primitive substitution $\sigma \in \text{Stab}(x)$ has a power that is (tamely) conjugate to a power of θ , that is, there exists a tame substitution τ such that $\sigma^i = \tau\theta^j\tau^{-1}$, for some $i, j \geq 1$.

In particular, if x is a dendric word, any primitive substitution in $\text{Stab}(x)$ is a tame substitution.

Theorem 3. *Let (X, S) be an aperiodic minimal dendric subshift. Then it admits no rational topological eigenvalue.*

Corollary 4. *Let (X, S) be an aperiodic minimal dendric subshift. Then, it can neither be generated by a primitive constant length substitution, nor be a Toeplitz subshift.*

Our proofs rely on the notion of return words and on the so-called Return Theorem [4] that states that for every infinite dendric word defined over the alphabet A , the set of (right) return words is a basis of the free group generated by the alphabet A .

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