PRIMES IN GEOMETRIC SERIES AND FINITE PERMUTATION GROUPS

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ABSTRACT. As a consequence of the classification of finite simple groups, the classification of permutation groups of prime degree is complete, apart from the question of when the natural degree $(q^n - 1)/(q - 1)$ of $L_n(q)$ is prime. We present heuristic arguments and computational evidence to support a conjecture that for each prime $n \ge 3$ there are infinitely many primes of this form, even if one restricts to prime values of q.

1. INTRODUCTION

The study of transitive permutation groups of prime degree goes back to the work of Galois on polynomials of prime degree. It is sometimes asserted that the groups of prime degree are now completely known, as a consequence of the classification of finite simple groups. This assertion is true only if one ignores an apparently difficult number-theoretic problem, namely the existence or otherwise of infinitely many primes of a particular form. The list of such permutation groups includes several easily described infinite families, three relatively small sporadic examples, and one other family which will be the subject of this note.

Let *p* be a prime, and let $q = p^e$, $e \ge 1$, be a prime power. The projective special linear groups $L_n(q) = PSL_n(q)$ and some closely related groups act doubly transitively, with degree

(1)
$$m = \frac{q^n - 1}{q - 1} = 1 + q + q^2 + \dots + q^{n-1},$$

on the points or hyperplanes of the projective space $\mathbb{P}^{n-1}(\mathbb{F}_q)$ for integers $n \ge 2$ and prime powers $q \ge 2$.

Definition 1.1 (Projective prime). If the number *m* of points and of hyperplanes of the projective space $\mathbb{P}^{n-1}(\mathbb{F}_q)$, defined by (1), is prime, then we call it a *projective prime*.

Remark 1.2 (*n* prime). A necessary condition for *m* in (1) to be prime is the primality of the exponent *n* since otherwise the polynomial $1 + t + t^2 + \cdots + t^{n-1}$ would be reducible over \mathbb{Z} .

The only projective primes with n = 2 are the Fermat primes $m = 2^{2^k} + 1$, while those with q = 2 are the Mersenne primes, of the form $m = 2^n - 1$ with *n* prime. However, there are many others, such as m = 13 with n = q = 3. An interesting case is the Mersenne prime

$$m = 31 = 1 + 2 + 4 + 8 + 16 = 1 + 5 + 25$$

Of course, it is an open problem whether there are infinitely many Fermat or Mersenne primes; at the time of writing, only five Fermat primes (with k = 0, ..., 4) and 51 Mersenne primes are

²⁰¹⁰ Mathematics Subject Classification. 11A41, 11N05, 11N32, 20B05, 20B25.

Key words and phrases. Permutation group, prime degree, projective space, Bunyakovsky conjecture, Goormaghtigh conjecture.

known to exist. More generally, the existence of infinitely many projective primes seems to be an open problem.

As in the case of Mersenne primes, there is plausible heuristic evidence, given in Section 4, to support a conjecture that there are infinitely many projective primes. Indeed, there is much stronger computational support for this, even in restricted cases such as when n = 3 and q is prime (see Section 6). Our aim in this note is to put forward such evidence, in the hope of inspiring specialists in number theory to address this problem. Thus, we formulate the following

Conjecture 1.3 (Projective primes). There are infinitely many projective primes.

Remark 1.4 (Addition of primes). More than once, some mathematicians (and physicists), among them very eminent ones, have expressed the opinion that the two famous Goldbach conjectures are completely devoid of interest since they concern the addition of primes whereas primes are created to be multiplied, not added. Our note may be considered as a strong case for the interest, in certain contexts, of the addition of primes and of numbers related to them (such as prime powers). By the way, one of the Goldbach conjectures has already been proved by Harald Helfgott [28].

Notation 1.5 (Primes). Notation for primes depends on the context. We follow the tradition of denoting a prime number by the letter p when this number is treated alone. If, however, there are more than one prime number involved, as is the case, for example, in formula (1), then one of these numbers may be denoted by m or by some other letter.

2. TRANSITIVE PERMUTATION GROUPS OF PRIME DEGREE

This section summarises the background in finite permutation groups from which the problem stated in the Introduction arises. Any reader who is interested only in the problem itself can safely omit this section. For more details on permutation groups of prime degree, see [13, §3.5] or [30, §V.21].

An elementary but important fact about transitive groups of prime degree is that they are all primitive, that is, they leave invariant no non-trivial equivalence relations. In particular, this means that rational or meromorphic functions of prime degree cannot be compositions of those of lower degree. Groups of prime degree are also rather rare: for example, there are 2 801 324 transitive groups of degree 32 (all but seven of them imprimitive), and only twelve groups of degree 31; similarly there are 315 842 groups of degree 40, but only ten of degree 41, and six of degree 47. The number of transitive groups of degree m \leq 47, $m \neq$ 32. The GAP system [21] contains a list of all primitive groups of degrees $m \leq$ 2499, and therefore, in particular, a list of all the groups of prime degrees up to the same limit.

One of the sections of the memoir by Évariste Galois [19]¹ is called "Application to irreducible equations of prime degree". If we translate the work of Galois on polynomials and their roots into

¹In his preface of 16 January 1831, Galois writes that this text is an "extrait d'un ouvrage que j'ai eu l'honneur de présenter à l'Académie il y a un an". The French word "ouvrage" means either a book or just a large piece of work. This larger text was sent to Fourier for refereeing. Fourier had suddenly died, and the manuscript was never found among his papers.

modern terminology, he showed that the solvable groups of prime degree p are the subgroups G of the 1-dimensional affine group

$$AGL_1(p) = \{t \mapsto at + b \mid a, b \in \mathbb{F}_p, a \neq 0\}$$

containing the translation subgroup $C_p = \{t \mapsto t + b \mid b \in \mathbb{F}_p\}$. Such groups G are semidirect products $G = C_p \rtimes C_d$ where C_d acts as a subgroup of \mathbb{F}_p^* for some divisor d of p - 1. There is one group G for each such d, including $G = C_p$ for d = 1 and $G = AGL_1(p)$ for d = p - 1.

This directs our attention to the nonsolvable groups of prime degree. Burnside [8, §251] showed that any such group G must be doubly transitive (as is $AGL_1(p)$, unlike its proper subgroups). In fact, in this case elementary arguments show that a minimal normal subgroup S of G is a nonabelian simple group, which is also transitive of degree p, with trivial centraliser $C_G(S)$ in G. Thus G acts faithfully by conjugation on S, so

$$S \leq G \leq \operatorname{Aut} S$$
.

This reduces the problem to that of determining the nonabelian simple groups S of prime degree p, and then studying their automorphism groups for possible subgroups G of degree p (the action of S need not extend to all subgroups of Aut S).

The classification of finite simple groups was announced around 1980, though not completely proved until over twenty years later. One consequence (see [9], for example) was the classification of doubly transitive finite permutation groups. There are eight families, described in some detail in [13, §7.7] and summarised in [10, §7.4]². As far as our problem is concerned, most of them can be ignored, as their degrees are composite: for example, the symplectic groups $\text{Sp}_{2n}(2)$ have degrees $2^{n-1}(2^n \pm 1)$, while the unitary and 'small' Ree groups over \mathbb{F}_q have degree $q^3 + 1$, divisible by q + 1. The groups which survive this elimination process are listed in the following theorem:

Theorem 2.1 (Transitive groups of prime degree). The nonabelian simple permutation groups S of prime degree, together with any transitive groups $G \le \text{Aut } S$ of degree equal to that of S, are as follows:

- (a) alternating groups $S = A_p$ for primes $p \ge 5$, together with the corresponding symmetric groups $Aut(A_p) = S_p$;
- (b) $G = S = L_2(11)$ for p = 11, acting on the cosets of a subgroup A_5 (two representations, on two conjugacy classes of such subgroups, equivalent under Aut $S = PGL_2(11)$), and the Mathieu groups $S = M_{11}$ and M_{23} , acting on Steiner systems with p = 11 and p = 23 points;
- (c) groups G such that $S = L_n(q) \le G \le P\Gamma L_n(q) \le Aut(L_n(q))$, acting on the points or hyperplanes of the projective space $\mathbb{P}^{n-1}(\mathbb{F}_q)$ when the degree $m = (q^n 1)/(q 1)$ is a projective prime, $m \ge 5$.

If we also include the affine groups, where

(d) $C_p \leq G \leq AGL_1(p)$ for primes p,

²Note that Aut M_{22} , of degree 22, is omitted from [13, p. 252]. Similarly, $L_2(11)$, of degree 11, is omitted from the list of groups of prime degree in [30, §V.21.2], though it is mentioned in II.8.28(6) and in the Errata in the 2nd printing.

then we have a complete list of the transitive groups of prime degree.

Remark 2.2 (Commentaries on Theorem 2.1).

- (1) The group $L_2(11)$ in (b) is one of three cases, known already to Galois, in which the simple group $L_2(p)$ has a non-trivial transitive representation of degree less that p + 1, specifically of degree p = 5, 7 or 11 on the cosets of a subgroup isomorphic to A_4 , S_4 or A_5 . The first case appears in both (a) and (c), via the isomorphisms $L_2(5) \cong A_5 \cong L_2(4)$, while the second appears in (c) via the isomorphism $L_2(7) \cong L_3(2)$. The group Aut ($L_2(11)$) = PGL₂(11) does not have a representation of degree 11; hence, only the group $L_2(11)$ is a member of our list.
- (2) For the two Mathieu groups in (b) we have $Aut(M_{11}) = M_{11}$ and $Aut(M_{23}) = M_{23}$.
- (3) In (c), we have Aut $(L_n(q)) = P\Gamma L_n(q)$ if n = 2, but if $n \ge 3$ then Aut $(L_n(q))$ contains $P\Gamma L_n(q)$ with index 2, the 'extra automorphism' arising from the point-hyperplane duality of $\mathbb{P}^{n-1}(\mathbb{F}_q)$.
- (4) In (d), the group $AGL_1(p)$ is *not* the automorphism group of C_p (indeed, $Aut(C_p) \cong C_{p-1}$). The case (d) does not correspond to the general scheme of the cases (a), (b), (c) since, as explained above, the group $AGL_1(p)$ is solvable.

For a given projective prime m, the groups G in (c) are easily determined: they correspond bijectively to the subgroups of

$$\mathrm{PFL}_n(q)/\mathrm{L}_n(q) \cong (\mathrm{PGL}_n(q)/\mathrm{L}_n(q)) \rtimes \mathrm{Gal}\,\mathbb{F}_q \cong \mathrm{C}_d \rtimes \mathrm{C}_e$$

where $d = \gcd(q - 1, n)$ and $q = p^e$ for some prime p. In fact, if m is prime then n is prime (see Remark 1.2) and $q \not\equiv 1 \mod (n)$, so d = 1, the groups $PGL_n(q)$ and $L_n(q)$ coincide, and $P\Gamma L_n(q)/L_n(q) \cong C_e$. The real problem is to know which primes are projective, and thus correspond to groups in (c), and in particular whether or not there are infinitely many of them.

Although this paper concentrates on those cases where $L_n(q)$ has prime degree, there is also interest in cases such as $L_5(3)$ where its natural degree *m* is a prime power (11² in this case). For example, Guralnick [25] has shown that if a nonabelian simple group *S* has a transitive representation of prime power degree, then *S* is an alternating group or $L_n(q)$ acting naturally, or $L_2(11)$, M_{11} or M_{23} acting as in Theorem 2.1(b), or the unitary group $U_4(2) \cong Sp_4(3) \cong O_5(3)$ permuting the 27 lines on a cubic surface. In particular, *S* is doubly transitive in all cases except the last, where it has rank 3, that is, three orbits on ordered pairs. See also [14], where Estes, Guralnick, Schacher and Straus have shown that for each prime *p* there are only finitely many *e*, *q*, *n* \ge 3 such that $p^e = (q^n - 1)/(q - 1)$.

Another related topic which we will not address here is the Feit–Thompson Conjecture [17] (see also [26, Problem B25]), that if p and q are distinct primes then $(p^q - 1)/(p - 1)$ does not divide $(q^p - 1)/(q - 1)$. A proof of this would significantly shorten the (very long) proof of the theorem [18] that groups of odd order are solvable.

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3. The Bunyakovsky Conjecture

Viktor Bunyakovsky (1804–1889) was a Russian mathematician and a disciple of Cauchy. In Russia he is mainly known for the Cauchy–Bunyakovsky inequality which, in the Western tradition, is named after Cauchy–Schwarz. (As is stated in the Wikipedia, Bunyakovsky "... is credited with an early discovery of the Cauchy–Schwarz inequality, proving it for the infinite dimensional case in 1859, many years prior to Hermann Schwarz's works on the subject.")

In 1857, Bunyakovsky formulated the following conjecture (see [6, 7]).

Conjecture 3.1 (Bunyakovsky Conjecture). *The following fairly obvious necessary conditions for* a polynomial $f(t) \in \mathbb{Z}[t]$ to have infinitely many prime values for $t \in \mathbb{N}$ are also sufficient:

- *the leading coefficient of f should be positive,*
- *f* should be irreducible,
- the integers f(t) for $t \in \mathbb{N}$ should have greatest common divisor 1.

The last condition is needed in order to avoid examples such as $f(t) = t^2 + t + 2$, which satisfies the first two conditions but has even values for all $t \in \mathbb{N}$; Bunyakovsky gives the surprising example $f(t) = t^9 - t^3 + 2520$, which is irreducible but has all its values divisible by 504. His conjecture is a special case of Schinzel's Hypothesis H [40], which concerns finite sets of polynomials simultaneously taking prime values.

Remark 3.2 (Verification of the coprimality of f(t) for $t \in \mathbb{N}$). The existence of examples like the one above leads to the following question: how to verify that the greatest common divisor of f(t) for $t \in \mathbb{N}$ is 1? A method (today we would say, an algorithm) proposed by Bunyakovsky is based on the following observations.

- (1) Let $f(t) = c_n t^n + \cdots + c_1 t + c_0$. If a prime *p* divides all the values of f(t) for $t \in \mathbb{N}$, then *p* is a divisor of c_0 . Indeed, substituting t = p in f(t) and taking the result modulo *p* we get $c_0 \equiv 0 \mod (p)$. Thus, we have only a finite number of primes *p* to test.
- (2) Let h(t) be a polynomial of degree k < p. Then all the values of h(t), $t \in \mathbb{N}$, are divisible by p if and only if all the coefficients of h are divisible by p. Indeed, otherwise, reducing h(t) modulo p we would get a non-zero polynomial of degree less that p which would have p roots.
- (3) All the values of the polynomial $t^p t$ are obviously divisible by p. Let h(t) be the remainder of f(t) on division by $t^p t$. All that remains is to determine whether the coefficients of h(t) are all divisible by p.

The conjecture is true for deg(f) = 1: this is Dirichlet's Theorem on primes in an arithmetic progression (see [4, §5.3.2] for a proof). However, it has not been proved for any polynomial of degree greater than 1, including the case $f(t) = t^2 + 1$ (see [26, §A1], [27, §2.8] or [38, Ch. 3.IVD]); this is sometimes called Landau's problem, though in fact it goes back to Euler [15]. In our case we have the advantage that we are not restricted to a single polynomial: we may consider polynomials $f(t) = 1 + t + t^2 + \cdots + t^{n-1}$ for any prime $n \ge 3$. (Since we have nothing to add to the current state of knowledge or ignorance concerning Fermat primes, we will assume for the rest of this paper that $n \ne 2$.) On the other hand, we require prime values of f(t) where t is a prime power, so a proof

of the Bunyakovsky Conjecture for such a polynomial would not necessarily yield infinitely many projective primes.

Finally, we note that the Bunyakovsky Conjecture has recently arisen in a similar way in the construction by Amarra, Devillers and Praeger [2] of block-transitive point-imprimitive 2-designs with specific parameters.

4. HEURISTIC ARGUMENTS

In this section we will present some heuristic arguments to support the conjecture that there are infinitely many projective primes $m = (q^n - 1)/(q - 1)$. They are based on heuristic arguments used elsewhere in considering the distribution and number of primes of a given form. In particular, some of the arguments in this section are adapted from Wagstaff's treatment [41] of conjectures of Gillies, Lenstra and Pomerance about Mersenne primes, and its summary in Prime Pages³. Of course heuristic arguments, based on assumptions which, although plausible, cannot be rigorously justified, do not prove anything (in particular, see the warning in Section 4.5). However, they may suggest results which one could attempt to prove by more legitimate means. Authors of classic texts did not disdain such kind of arguments: see, for example, Sections 2.5 and 22.20 of the book [27] by Hardy and Wright, where they present heuristic evidence that there are only finitely many Fermat primes whereas there are infinitely many prime pairs; see also the discussion of probabilistic methods in [34, Notes on Ch. 8.3] and Pólya's carefully-qualified defence of heuristic reasoning in number theory in [35].

Beside the "general" conjecture of infinitely many projective primes we also formulate a number of "specific" (and therefore stronger) conjectures concerning projective primes of some specific forms. Their plausibility is based mainly on a series of computational results presented in Sections 6 and 7.

4.1. **Prime divisors of** *m*. We consider firstly the case of any fixed prime $n \ge 3$, and secondly that of any fixed prime power *q*. In each case, we will need the following lemma in order to give better estimates for the number of projective primes up to some bound.

Lemma 4.1 (Prime divisors of *m*). Let $m = (q^n - 1)/(q - 1)$ for some integer *q* and prime $n \ge 3$, and let *r* be a prime dividing *m*. Then either $r \equiv 1 \mod (2n)$ (so in particular $r \ge 2n + 1$), or r = n with $q \equiv 1 \mod (n)$. Conversely, if $q \equiv 1 \mod (n)$ then *m* is divisible by *n*.

Proof. If a prime *r* divides *m* then $q^n \equiv 1 \mod (r)$. Since *n* is prime, it follows that either *n* divides the order r - 1 of the multiplicative group \mathbb{F}_r^* , or $q \equiv 1 \mod (r)$.

If *n* divides r - 1 then $r \equiv 1 \mod (n)$. Clearly $m = 1 + q + \dots + q^{n-1}$ is odd, and hence so is *r*, so $r \equiv 1 \mod (2n)$ since *n* is odd and hence $r \ge 2n + 1$.

If $q \equiv 1 \mod (r)$ then

$$m = 1 + q + \dots + q^{n-1} \equiv \underbrace{1 + 1 + \dots + 1}_{n \text{ times}} \equiv n \mod (r).$$

³See https://primes.utm.edu/mersenne/heuristic.html.

However, $m \equiv 0 \mod (r)$, so $n \equiv 0 \mod (r)$ and hence r = n since n is prime. The converse is obvious.

Thus *m* is not divisible by any prime $r \le 2n$, except the prime r = n if $q \equiv 1 \mod (n)$.

Example 4.2 (For Lemma 4.1). Let n = 3. If $q = 11 \neq 1 \mod (3)$ then $m = 133 = 7 \cdot 19$; this is divisible by the primes r = 7 and 19, both greater than 2n = 6. However, if $q = 16 \equiv 1 \mod (3)$ then $m = 273 = 3 \cdot 7 \cdot 13$, divisible by the prime r = n = 3 in addition to r = 7 and 13. Note that the 'large' primes 7, 13 and 19 appearing here as divisors of m are all congruent to 1 mod (2n).

4.2. Fixed *n*, while $q = p \rightarrow \infty$. Let us fix a prime $n \ge 3$, and consider whether $m = (q^n - 1)/(q - 1)$ is prime. For simplicity we will restrict q to be prime, rather than a prime power; therefore, from now on we will denote it by p instead of q. By the Prime Number Theorem (see [27, Theorem 6 and Ch. XXII] for example), the number of primes p in the range $1 \le p \le x$ is approximately $x/\ln(x)$ for large x. However, if $p \equiv 1 \mod (n)$ then m cannot be prime by Lemma 4.1, so we should restrict attention to the primes $p \not\equiv 1 \mod (n)$. Since primes are approximately evenly distributed between the non-zero congruence classes mod (n) (see [4, §5.3.2], for example), the number of primes p we should consider is therefore approximately $(n-2)x/(n-1)\ln(x)$.

Now $1 \le m \le (x^n - 1)/(x - 1)$, and the probability that a randomly-chosen integer m in this range is prime is approximately

(2)
$$\frac{1}{\ln((x^n - 1)/(x - 1))} \approx \frac{1}{n\ln(x) - \ln(x)} = \frac{1}{(n - 1)\ln(x)}$$

However, we know from Lemma 4.1 that $m \neq 0 \mod (r)$ for each prime $r \leq 2n$, including r = nsince $p \not\equiv 1 \mod (n)$. For each such r, excluding this one congruence class mod (r) multiplies the probability of m being prime by r/(r-1). If we regard congruences modulo distinct primes as statistically independent, then we should multiply the probability in (2) by P(2n), where

(3)
$$P(y) := \prod_{\text{prime } r \le y} \left(1 - \frac{1}{r}\right)^{-1}$$

for $y \ge 2$ and the product, as indicated, is over all primes $r \le y$. This gives an approximate probability

(4)
$$\frac{P(2n)}{(n-1)\ln(x)}$$

that *m* is prime. For fixed *n* this has the form $c_n/\ln(x)$ for a constant

$$c_n := \frac{P(2n)}{(n-1)}$$

If *n* is small one can easily calculate c_n : for instance $c_3 = 15/8$ and $c_5 = 35/32$. For large *n* one can approximate c_n by using a theorem of Mertens (see [32], [27, §22.9] or [34, Theorem 8.8(e)]) that

$$\prod_{\text{prime } r \le y} \left(1 - \frac{1}{r} \right) \sim \frac{\mu}{\ln(y)} \quad \text{as} \quad y \to \infty,$$

where $\mu := e^{-\gamma} = 0.561459...$ and γ is the Euler–Mascheroni constant 0.577215... This gives an approximate probability

(5)
$$\frac{c_n}{\ln(x)} \sim \frac{e^{\gamma} \ln(2n)}{(n-1) \ln(x)}$$

that *m* is prime.

If we multiply this probability by the approximate number of primes $p \neq 1 \mod (n)$ in the range $1 \leq p \leq x$, namely $(n-2)x/(n-1)\ln(x)$, we see that the expected number of primes *m* arising in this way is approximately

(6)
$$\frac{c_n(n-2)x}{(n-1)(\ln(x))^2} \sim \frac{e^{\gamma}(n-2)\ln(2n)x}{(n-1)^2(\ln(x))^2} \approx \frac{1.781(n-2)\ln(2n)x}{(n-1)^2(\ln(x))^2}.$$

Since this number tends to $+\infty$ as $x \to \infty$ for fixed *n*, this suggests that we should obtain infinitely many projective primes *m* in this way for any fixed prime $n \ge 3$ (and likewise if we allow *q* to be an arbitrary prime power).

For each fixed *n* the estimate in (6) has the form $C_n x/(\ln(x))^2$ for some constant C_n depending only on *n*. This is analogous to the Hardy–Littlewood estimate $Cx/(\ln(x))^2$ for the number $\pi_2(x)$ of twin prime pairs p, p + 2 with $p \le x$ (see [36]), where

$$C = 2 \cdot \prod_{\text{prime } r \ge 3} \frac{r(r-2)}{(r-1)^2} \approx 1.320323632.$$

Example 4.3 (n = 3). If we take n = 3, so that $c_n = 15/8$, then the number of primes $m = 1 + p + p^2$ for primes $p \le x$ should be approximately

(7)
$$\frac{15x}{16\ln(x)^2}$$

for large *x*. This estimate is compared with computational evidence in Section 6.4 (see Table 2).

4.3. Fixed prime power $q = p^e$ with $e \ge 2$. Instead, let us now fix q and let $n \to \infty$.

Lemma 4.4. If $e \ge 2$ then there are no projective primes $m = (q^n - 1)/(q - 1)$ with n > e.

Proof. If $e \ge 2$, so that q is a prime power but not itself a prime, we have

$$m = \frac{q^{n} - 1}{q - 1} = \frac{(1 + p + \dots + p^{n-1})(1 + p^{n} + \dots + p^{n(e-1)})}{1 + p + \dots + p^{e-1}}.$$

This is clearly composite if n > e since the two factors in the numerator are each larger than the denominator.

Thus, for a fixed q with $e \ge 2$ we can have only a finite number of projective primes m.

Remark 4.5 (e = 2). If e = 2 and m is prime then n = 2 (and hence $q = p^2$ is even, so that p = 2 and m = 1 + q = 5), against our earlier assumption; thus in dealing with prime powers $q = p^e > p$ (as in Section 6.8, where a number of examples are given) we will generally assume that $e \ge 3$.

Since the ultimate goal of this line of research is to complete the classification of the permutation groups of prime degree, it would still be of interest to know the finitely many projective primes arising for each proper prime power q.

4.4. Fixed prime p while $n \to \infty$. Let us therefore take e = 1, so we fix a prime p, and consider the primality of $m = (p^n - 1)/(p - 1)$ for odd primes n as $n \to \infty$. By Lemma 4.1 we may exclude any primes n dividing p - 1, since they cannot give prime values of m. Then m is not divisible by any prime $r \le 2n$.

By the Prime Number Theorem, for large *n* a randomly-chosen integer close to $(p^n - 1)/(p - 1)$ is prime with probability approximately

(8)
$$\frac{1}{\ln((p^n - 1)/(p - 1))} \approx \frac{1}{n\ln(p) - \ln(p - 1)}$$

However, *m* is not uniformly distributed, since it is coprime to each prime $r \le 2n$. As before, for each such *r* this excludes a proportion 1/r of the integers close to $(p^n - 1)/(p - 1)$, so we should multiply the probability in (8) by $M(2n) \sim e^{\gamma} \ln(2n)$, giving an approximate probability

(9)
$$\frac{e^{\gamma}\ln(2n)}{n\ln(p) - \ln(p-1)} \sim \frac{e^{\gamma}\ln(n)}{n\ln(p)}$$

that m is prime, for large n. (For small p, as in Section 6.7, this last approximation could induce significant errors.)

If we choose the prime *n* uniformly and randomly from the range $p \le n \le x$ for some large *x* (so that most such *n* are large as above), then the expected number of primes *m* arising is the sum of the probabilities in (9), that is

$$\frac{e^{\gamma}}{\ln(p)}\sum_{n}\frac{\ln(n)}{n}$$

where the sum is over all primes *n* such that $p \le n \le x$. Now

$$\sum_{n} \frac{\ln(n)}{n} \approx \ln(x) - \ln(p-1).$$

(see [27, Theorem 425]), so the expected number of primes *m* is approximately

(10)
$$\frac{e^{\gamma}(\ln(x) - \ln(p-1))}{\ln(p)} \sim \frac{e^{\gamma}\ln(x)}{\ln(p)} \approx \frac{1.781\ln(x)}{\ln(p)}$$

Since this tends to $+\infty$ with x, we may expect to obtain infinitely many projective primes m from any given prime $p \ge 2$. This estimate is compared with computational evidence in Section 6.7 (see Table 4).

4.5. A warning. Invoking the independence of congruences modulo different primes in order to make heuristic estimates, as we did in Section 4.2, has previously generated controversy: for instance, Wagstaff discusses this in [41], citing criticism by Lenstra in [31]. This is best illustrated with Pólya's discussion in [35] of the following well-known paradox.

Based on the type of argument used in Section 4.2, one can attempt a heuristic proof of the Prime Number Theorem. An integer x is prime if and only if $x \neq 0 \mod (r)$ for each prime $r \leq x$. For each

such *r* this event has probability (r - 1)/r, so by regarding these events as mutually independent, and by using Mertens's Theorem, one might expect *x* to be prime with probability

(11)
$$\prod_{\text{prime } r \le x} \left(1 - \frac{1}{r} \right) \sim \frac{\mu}{\ln(x)} \quad \text{as} \quad x \to \infty,$$

where $\mu = e^{-\gamma} = 0.561459...$ However, the correct asymptotic probability is $1/\ln(x)$, so this argument underestimates the probability of *x* being prime (and hence the values of the prime-counting function $\pi(x)$) by a factor of μ . Of course, it is sufficient to eliminate prime factors $r \le x^{1/2}$, rather than $r \le x$, so this alternative approach gives a second estimate

(12)
$$\prod_{\text{prime } r \le x^{1/2}} \left(1 - \frac{1}{r} \right) \sim \frac{\mu}{\ln(x^{1/2})} = \frac{2\mu}{\ln(x)} \quad \text{as} \quad x \to \infty.$$

This overestimates the correct probability by a factor of $2\mu = 1.122918...$, that is, by about 12%. If, as suggested by Pólya in [35], one takes the product over all primes $r \le x^{\mu}$ then the correct formula is obtained. Pólya confesses that it is not clear why what he calls this "trick of the magic μ " works here (Wagstaff [41] calls it a "fudge factor"), but he goes on to argue that mathematicians should imitate physicists by adapting their theories to fit experimental data when such paradoxes arise. Similar phenomena are discussed by Pólya [35] in relation to prime pairs and their generalisations, and by Wagstaff [41] in relation to the distribution of divisors of Mersenne numbers.

The great Russian mathematician Andreĭ Kolmogorov used to mention the following episode (the second author heard it directly from him). Kolmogorov was once present at a talk given by a prominent Russian physicist. The latter, basing his reasoning on some physical ideas, introduced the density of a probability distribution on a certain space. Then, he integrated this density and obtained π . At this point, Kolmogorov used to say, I would conclude that we had got a contradiction, and therefore all the reasoning was wrong. But the conclusion of the physicist was different. Thus, he said, we must divide the initial formula for the density by π . It seems that Pólya would rather line up with the physicist.

In our case, an appropriate choice of prime factors of *m* to avoid is also an intricate matter. For example, for n = 3, as we will see later, in Section 6.4 and Table 2, formula (7) overestimates the number of projective primes. But we know that, beside the "small primes" 2, 3 and 5, Lemma 4.1 also forbids all primes of the form $r \equiv -1 \mod (6)$. However, even if we adjoin to the product

$$P(6) = \left(\left(1 - \frac{1}{2}\right) \left(1 - \frac{1}{3}\right) \left(1 - \frac{1}{5}\right) \right)^{-1}$$

not all such corresponding terms but only $(1 - 1/11)^{-1}$, we will get $c'_3 = 33/16$ instead of $c_3 = 15/8$, and estimate (7) will be replaced with $33x/32 \ln(x)^2$, which will overestimate the number of projective primes even more than (7) does. On the other hand, if we remove the factor $(1 - 1/5)^{-1}$ from P(6), we will underestimate the desired number.

But this is not yet the end of the story. As we will see in Section 6.5 and Figure 2, the ratio of our estimates to the true number of projective primes grows, that is, the estimates grow faster

than the numbers they are supposed to estimate. There are even reasons to believe the overestimate grows to infinity (though slowly). Obviously, the constant factor does not play any part in this process: the rate of growth of the estimates depends only on the behavior of the function $x/\ln(x)^2$. Therefore, it is reasonable to suppose that, for very large x, we will have indeed to eliminate 5 from the set of forbidden primes. Note however that this process of eliminating or adjoining forbidden primes is not based on any solid theoretical foundation: it is purely empirical. Therefore, instead of making artificial choices of which primes to include in the product, it seems at least as reasonable to consider other functions instead of $x/\ln(x)^2$. This will be done in Section 6.5.

5. PRIMALITY TESTING

Before presenting the experimental results aiming to support our main conjecture (that there are infinitely many projective primes), let us briefly discuss two problems: the factorization of integers into prime factors, and the testing of primality. The problems are, evidently, related to each other, but there is an abyss between their complexities.

5.1. **Integer factorization in modern times.** According to [3], in 1977, Ronald Rivest, in a letter to Martin Gardner, estimated that

"... factoring a 125-digit number which is the product of two 63-digit prime numbers would require at least 40 quadrillion years using the best factoring algorithm known, assuming that $a \cdot b \mod(c)$ could be computed in 1 nanosecond, for 125-digit numbers a, b, and c."

It sounded like a solemn chorus from Purcell's *Dido and Aeneas*: "Never! Never! Never!" The same year, Gardner [22] launched a challenge: it was proposed to factor a 129-digit number which was the product of a 64-digit and a 65-digit prime. Apparently hopeless, whatever the future progress in computer technology would be.

Subsequent years saw spectacular progress in factorization algorithms. Finally, 17 years later, in 1994, the above 129-digit number was successfully factored. The project involved some 600 volunteers, 1600 computers, and six months of computation. An account may be found in [3]; the strange expression "squeamish ossifrage" in the title was the message encrypted using this number by the RSA cryptographic method.

We skip a number of important developments during the next 15 years and go directly to a milestone of 2009. In December of that year, a 232-digit number was factored: it was a product two 116-digit primes. This result was the outcome of two years of work by a team of 13 researchers, and was crowned with a \$50 000 prize. As stated in [39],

"The CPU time spent on finding these factors by a collection of parallel computers amounted approximately to the equivalent of almost 2000 years of computing on a single-core 2.2 GHz AMD Opteron-based computer."

Certainly, 2000 years for a 232-digit number as compared to 40 000 000 000 000 000 years for a 125-digit one, is incredible progress. A convenient, and machine-independent measure of an effort in a large-scale computation is GHz-years; in the present case we have 4400 GHz-years of computing.

Ten more years have passed, and a 240-digit number was factored in November 2019, and a 250-digit number in February 2020. And here lies the current frontier of the capability of factoring algorithms. The next challenge, a 260-digit number, still waits for its turn to be factored.

5.2. **Testing.** The above examples show how difficult, in practical terms, the problem of factorization can be. However, there exist algorithms which establish whether a given integer is prime or composite, and this without ever trying to factor it. The most well-known, and the most used in practice, is the Rabin–Miller algorithm [37] (see also the compendium [11]). In particular, it is implemented in the Maple command *isprime*. Let us take the above-mentioned 260-digit number (which, we recall, is not yet factored) and see how this command works. The computation is carried out on a very modest laptop.

001539441\
031400051
910067030
106587591

We see that the correct answer is given, almost literally, "in no time at all". In fact, the Maple time-counter outputs the CPU time within an accuracy of 0.001 seconds. Therefore, 0. seconds time displayed in the above session means < 0.0005 seconds, rounded downwards.

Now consider a harder example, a 6153-digit number we will encounter in Section 6.2.

$$\begin{bmatrix} > M := \frac{1201^{1999} - 1}{1200} : \\ > number_of_digits := ceil(evalf(log10(M))); \\ number_of_digits := 6153 \\ > time(isprime(M)); \\ 13.314 \\ > isprime(M); \\ true \\ \end{bmatrix}$$

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This number is prime, and the computation took more than 13 seconds. A good result, but to perform this testing on a large scale, that is, with large series of numbers, can turn out to be time-consuming.

5.3. How the test works. The Rabin–Miller algorithm is probabilistic. In order to determine whether a given number *m* is prime it takes a random element $t \in \mathbb{Z}_m$ and verifies a necessary primality condition. The condition itself is simple, so we give it here.

The number m-1 is even; suppose it is equal to $m-1 = (2l+1) \cdot 2^k$. Compute in \mathbb{Z}_m

 $a_0 = t^{2l+1}$, and then $a_i = a_{i-1}^2$ for i = 1, ..., k, so that $a_k = t^{m-1}$.

If one of the following holds then *m* is composite:

- (1) While computing the sequence a_i , we come for the first time to $a_i = 1$ but the previous number $a_{i-1} \neq -1$. Indeed, in this case the equation $a^2 = 1 \mod (m)$ has, beside two obvious roots 1 and -1, a third root a_{i-1} .
- (2) We get $a_k = t^{m-1} \neq 1$. This contradicts Fermat's little theorem.

Thus, if the test tells us that m is composite then this statement is true, and no probability is involved. If, however, neither of the two above conditions is satisfied, we conclude that m is *probably prime*. Rabin [37] showed that the probability of an erroneous answer is bounded by 1/4; usually it is much smaller. For large m, in the majority of cases this probability is infinitesimally small. A dialogue from Gilbert and Sullivan's *I am the captain of the Pinafore* comes to mind: "What, never? No, never. What, never? Well, hardly ever". Nevertheless, in order to be on the safe side, the test is repeated many times with different (random) values of t. This, by the way, explains why the treatment of a prime number takes much more time than that of a composite number of the same size.

Notice that raising a number *a* to a power a^r needs $O(\log r)$ arithmetic operations: we compute first a, a^2, a^4, a^8, \ldots (taking squares every time), and then multiply the terms corresponding to the binary expression of the exponent *r*. Notice also that, in our case, all computations are made modulo *m*, so that the size of the numbers remains bounded.

5.4. **Polynomial-time algorithms.** The subject of primality testing and factorization has many ramifications. We only mention very briefly a few of them. We recommend, for an interested reader, a very concise and clear overview [5] and a more modern and advanced exposition in [16] (especially Chapter 5, "Primality testing—an overview").

There are several algorithms for primality testing whose complexity is polynomial in the size of tested numbers. However, for most of them the estimation of complexity is based on some as yet unproved hypotheses.

Notation 5.1 (Simplified measure of complexity). Denote $k := \log m$, and denote $\widetilde{O}(k^s) := O(k^{s+\varepsilon})$ for all $\varepsilon > 0$.

This notation allows one to simplify complexity estimates for algorithms such as, for example, the Schönhage–Strassen algorithm of multiplication of long integers: we may now write just $\tilde{O}(k)$ instead of $O(k \cdot \log k \cdot \log \log k)$.

The complexity of the Rabin–Miller algorithm is $\widetilde{O}(k^2)$: here O(k) is the number of arithmetic operations, and $\widetilde{O}(k)$ is the complexity of an individual operation.

Four years before Rabin, Miller used the same test but in a deterministic way. Namely, it suffices to make the test for all $t \le 2 \log(m)^2 = 2k^2$, provided that the Extended Riemann Hypothesis is true. Thus, this algorithm, of complexity is $\tilde{O}(n^4)$, while being deterministic, is based on an unproved conjecture. Also, the factor $2k^2$ is not innocuous. For $k \sim 6000$, as in the above example, it transforms seconds into years.

Remark 5.2 (Are long computations reliable?). It is important to note that in a long computation there is a significant probability of a hardware error. This probability is much greater than that in the Rabin–Miller test.

In a revolutionary work [1], an *unconditional* polynomial time algorithm for primality testing was given for the first time. Here 'unconditional' means that the estimate of its complexity does not depend on any unproved statement. After several improvements its complexity is now established as $\tilde{O}(k^{15/2})$. It may also be $\tilde{O}(k^6)$ if another as yet unproved conjecture is valid. Its theoretical impact is great but its practical utility is very limited.

Another method is based on the theory of elliptic curves. It is commonly known as the ECPP algorithm, which means Elliptic Curve Primality Proving. The names we must mention here are Sh. Goldwasser, J. Kilian, A. Atkin and F. Morain. This algorithm is probabilistic; however, it is not of the "Monte-Carlo type" but of the "Las Vegas type". The latter means that it always gives the correct answer; it is the computation time which is random. It is polynomial *on average* if certain as yet unproved conjectures are true. Beside the correct answer, this algorithm also creates a *primality certificate*. A certificate is "something" which may be difficult to find but, once found, allows one to make a verification easily.

In [5], the following example is given. Consider the number $m = 4405^{2638} + 2638^{4405}$. It has 15071 digits. The *proof* of its primality by the ECPP algorithm was achieved in 5.1 GHz-years. This is a truly remarkable result if we compare it with other error-free algorithms. Note, however, that the Rabin–Miller algorithm gives the correct (though unproved) answer in less than two minutes.

5.5. A few comments. Since 1980, when Michael Rabin published his algorithm, not a single case of an erroneous answer has been observed. Even financiers, in their cryptographic protocols, rely entirely on this test. However, a mathematical mind resists accepting a "proof" which in principle might be wrong, even if the probability of such an event is infinitesimally small. What then to do if we have doubts about the validity of the conclusion 'prime' given by the probabilistic test? In our opinion, the most reasonable way to proceed is to run this test again once or twice. The test does not repeat exactly the same operations since it chooses different random elements of \mathbb{Z}_m every time. In this way the probability α of an error, already infinitesimal, will be replaced with α^2 or α^3 . (We may ask, rather provocatively: how many times can you repeat a two-minute test if you have 5.1 years at your disposal?)

And what if, by an incredible combination of chances, we take a composite number for a prime one? Well, let us recall that the aim of our particular study is to collect evidence that there are infinitely many projective primes. Therefore, one prime less or one prime more does not change much.

6. Computational evidence

6.1. Bunyakovsky's conjecture from an experimental perspective. The data in favor of the veracity of this conjecture abound. If we take, for example, $f(t) = t^2 + t + 1$ and count the number of integers $t \le 10^7$ for which f(t) is prime, we get 745 582 solutions. A very "modest" particular case of Bunyakovsky's conjecture is known as Landau's conjecture: it concerns $f(t) = t^2 + 1$. In this case the number of $t \le 10^7$ for which $t^2 + 1$ is prime is 456 362. There is little doubt that, at least in these two cases, the conjecture is true. No proof is, however, in view.

The main motivation of this note comes from group theory. Therefore, we will mainly consider not arbitrary values of t but only prime powers $t = p^e$, $e \ge 1$, and not arbitrary polynomials f(t)but only those of the form

$$f(t) = \frac{t^n - 1}{t - 1} = 1 + t + t^2 + \dots + t^{n-1}.$$

Remark 6.1 (Terminological). While speaking of *prime powers*, according to the context we may mean p^e with $e \ge 1$, that is, including "pure" primes, or, sometimes, with $e \ge 2$, in order to put prime powers in contrast with the pure primes whose exponent is e = 1.

6.2. First series of projective primes. A computer search has revealed 668 projective primes with $2 \le q \le 2000$ and $3 \le n \le 2000$, including one with 6153 decimal digits, arising from q = 1201 and n = 1999. It is interesting to note that only five pairs (q, n) out of 668 correspond to prime powers $q = p^e$ with $e \ge 2$, namely,

$$(q, n) = (2^3, 3), (2^7, 7), (2^9, 3), (3^3, 3), (11^3, 3).$$

All the other values of q are "pure" primes.

6.3. Number 31. A computer search of prime degrees up to 10^{12} reveals $L_3(5)$ and $L_5(2)$ as the only pair of groups $L_n(q)$ with the same natural degree in this range; it would be interesting to know whether any other such pairs exist.

Conjecture 6.2 (Number 31). *Beside* 31, *there are no other natural degrees common to two different projective groups* $L_n(q)$.

Remark 6.3 (Goormaghtigh conjecture). The Diophantine equation

$$\frac{x^n - 1}{x - 1} = \frac{y^k - 1}{y - 1}$$

has been studied by many authors (see [26, Problem B25], for example). In 1917, a Belgian engineer and amateur mathematician René Goormaghtigh⁴ (1893–1960) conjectured [23] that this equation, for $n \neq k$, $n, k \ge 3$, has only two solutions in \mathbb{N} : 1 + 2 + 4 + 8 + 16 = 1 + 5 + 25 = 31 and $1 + 2 + 4 + \cdots + 2^{12} = 1 + 90 + 90^2 = 8191$. However, 90 is not a prime power, so that there

⁴See https://forvo.com/word/ren%C3%A9_goormaghtigh to learn how to pronounce this name.

is no field with 90 elements. By the way, the number 8191 is prime. Therefore, it is an instance of Bunyakovsky's conjecture for two different polynomials (and certainly for many other ones, like $t^2 + 91$, for example), but it is a projective prime for only one of them, namely, $1 + t + t^2 + \cdots + t^{12}$.

In [12], it is proved that for fixed exponents $n \neq k$, $n, k \geq 3$, there can be only a finite number of solutions. For additional information about this equation see [24].

6.4. **Projective planes over prime fields.** Let us take only prime values p, not taking into account the prime powers $q = p^e$ with $e \ge 2$, let us fix n = 3 and consider projective primes $m = 1 + p + p^2$. Our colleague Jean Bétréma examined all primes $p \le 10^{11}$ using the package Primes.jl of the language Julia. It turns out that Julia is much more efficient than Maple for problems of this sort. We partially reproduce Bétréma's results in Table 1.

Segment	#(prime p)	#(prime <i>m</i>)	ratio	max p
$2, \ldots, 10^{10}$	455 052 511	15 801 827	3.473%	9 999 999 491
$10^{10},\ldots,2\cdot10^{10}$	427 154 205	13 882 936	3.250%	19 999 999 757
$2 \cdot 10^{10}, \ldots, 3 \cdot 10^{10}$	417 799 210	13 279 095	3.178%	29 999 999 921
$3 \cdot 10^{10}, \ldots, 4 \cdot 10^{10}$	411 949 507	12913713	3.135%	39 999 999 719
$4 \cdot 10^{10}, \ldots, 5 \cdot 10^{10}$	407 699 145	12 645 233	3.102%	49 999 999 619
$5 \cdot 10^{10}, \ldots, 6 \cdot 10^{10}$	404 383 577	12439618	3.076%	59 999 999 429
$6 \cdot 10^{10}, \ldots, 7 \cdot 10^{10}$	401 661 384	12 274 191	3.056%	69 999 999 287
$7 \cdot 10^{10}, \ldots, 8 \cdot 10^{10}$	399 359 707	12136112	3.039%	79 999 999 679
$8 \cdot 10^{10}, \ldots, 9 \cdot 10^{10}$	397 369 745	12010780	3.023%	89 999 999 981
$9 \cdot 10^{10}, \ldots, 10^{11}$	395 625 822	11910803	3.011%	99 999 999 977
Total	4 118 054 813	129 294 308	3.140%	99 999 999 977

TABLE 1. The second column gives the number of primes in the corresponding segment, while the third column gives the number of those primes p which create a projective prime $m = 1 + p + p^2$. The proportion of such primes among all the primes of the second column is given in the fourth column.

We may see from this table that the number of primes $p \le 10^{11}$ which produce a prime value of *m* is 129 294 308, the largest of them being 99 999 999 977. The corresponding projective prime is $m = 9\,999\,999\,5500\,000\,000\,507$. Such primes *p* represent approximately 3.140% of the total number 4 118 054 813 of primes up to 10^{11} .

Of course, this percentage diminishes together with the growth of the upper limit. For example, if we count the proportion of such primes up to 10^6 , we get 5.97%. Nevertheless, it is quite reasonable to conjecture that even in this very restricted situation there are infinitely many projective primes.

Table 2 compares the numbers of projective primes $m = 1 + p + p^2$ for primes $p \le x_i = i \cdot 10^{10}$, i = 1, ..., 10, with the heuristic estimates given by (7) in Section 4.2. It can be seen that the latter are of the right order of magnitude, but that they consistently over-estimate the number of such

x	#(prime $m \mid p \leq x$)	estimate (7)	ratio	estimate (13-14)	ratio
$1 \cdot 10^{10}$	15 801 827	1.7683×10^{7}	1.1190	1.5799306×10^{7}	0.999841
$2 \cdot 10^{10}$	29 684 763	3.3328×10^{7}	1.1227	2.9686686×10^7	1.000065
$3 \cdot 10^{10}$	42 963 858	4.8096×10^{7}	1.1195	4.2969637×10^{7}	1.000135
$4 \cdot 10^{10}$	55 877 571	6.2736×10^{7}	1.1227	5.5881270×10^{7}	1.000066
$5\cdot 10^{10}$	68 522 804	7.7239×10^{7}	1.1272	6.8526763×10^7	1.000058
$6 \cdot 10^{10}$	80 962 422	9.1332×10^{7}	1.1281	8.0965961×10^{7}	1.000044
$7 \cdot 10^{10}$	93 236 613	1.0524×10^{8}	1.1287	9.3237376×10^7	1.000008
$8\cdot 10^{10}$	105 372 725	1.1900×10^{8}	1.1293	1.0536780×10^{8}	0.999953
$9 \cdot 10^{10}$	117 383 505	1.3262×10^{8}	1.1298	1.1737691×10^{8}	0.999944
10^{11}	129 294 308	1.4614×10^{8}	1.1303	1.2927974×10^{8}	0.999887

TABLE 2. The second column gives the cumulative totals from the second column in Table 1, i.e. the number of projective primes m with n = 3 arising from primes $p \le x_i = i \cdot 10^{10}$ (i = 1, ..., 10); the third column gives an approximation for the estimate for this number from Section 4.2, while the fourth column gives the ratio of these two numbers. The meaning of the last two columns is explained in Section 6.5.

primes by about 12%. In Section 6.5 we present another estimate. We do not have any theoretical bases to support it, only empirical ones, but it approximates the values we need much better than the previous estimate. The results of this estimate are represented in the two last columns of Table 2.

6.5. An empirical estimate vs. the theoretical one: rectifying the anomaly. We see that for large x estimate (7) systematically overestimates the number of projective primes. Hence, let us instead consider an estimate of the following form:

(13)
$$y = \frac{Cx}{\ln(x)^{\alpha}}$$

where the constants C and α are to be found from empirical data.

Denote $z = x/y = \frac{1}{C} \ln(x)^{\alpha}$. Then, taking the logarithm of each side of this equation we get

$$\ln(z) = -\ln(C) + \alpha \cdot \ln(\ln(x)).$$

Thus, in the coordinates

 $u = \ln(\ln(x)), \qquad v = \ln(z)$

equation (13) takes the form of an equation of a straight line

$$v = a + \alpha u$$
 where $a = -\ln(C)$.

Our next steps are as follows:

(1) Take a number of pairs (x_i, y_i) where x_i are at our choice while y_i are the numbers of projective primes $m = 1 + p + p^2$ created from primes $p \le x_i$.

- (2) Compute the corresponding pairs (u_i, v_i) .
- (3) Put the points (u_i, v_i) on the plane of the coordinates (u, v), in the hope that they will be reasonably close to a straight line $v = a + \alpha u$.
- (4) Find the equation of this straight line and thus obtain the values of $C = e^{-a}$ and of α .

The points should be taken with care. As we will see in Section 7, numbers of projective primes behave rather randomly. Therefore, in order to get a reasonable estimate, the values of x_i must be large enough, and the spaces between them must also be large.

The numbers $x_i = i \cdot 10^{10}$, i = 1, ..., 10 satisfy both conditions. Taking x_i and y_i from the first two columns of Table 2 we get the results which are shown in Figure 1. We would say that they are even better than one might hope. The corresponding constants, found by the method of least squares, are as follows:

(14) a = -0.150383694, $C = e^{-a} = 1.162280117$, $\alpha = 2.104419156$.

The estimates given by (13) with the constants (14), and their ratios to the true number of projective primes are shown in the last two columns of Table 2.

Remark 6.4 (Overestimate). If the estimates (13-14) are correct, and it seems that they are, or at least if they are close to the correct asymptotic, then the *overestimate* of the ratio given by formula (7) tends to infinity, though rather slowly: it is proportional to $\ln(x)^{\alpha-2}$.



FIGURE 1. The horizontal axis corresponds to the variable $u = \ln(\ln(x))$, the vertical one to the variable $v = \ln(x/y)$ where y is meant to count projective primes. The ten distinguished points correspond to $x_i = i \cdot 10^{10}$, while y_i is the number of prime $p \le x_i$ such that $m = 1 + p + p^2$ is prime.

It would be interesting to understand the nature of the above constants. For example, α is reasonably close to $1 + 2\mu = 2.1229$ (the constant μ is defined in Section 4.2). It would be tempting to conjecture that they are equal, especially since μ appears in several other conjectures related to



FIGURE 2. Comparison of two estimates. The horizontal scale is logarithmic: an abscissa k corresponds to the number 10^k . White circles represent the ratios of estimates (7) to the true numbers of projective primes $m = 1 + p + p^2$ obtained from the primes $p \le 10^k$, k = 3, ..., 11. The solid circles represent the similar ratios for estimates (13-14). Notice that for smaller numbers both estimates *underestimate* the number of projective primes (both black and white points are below the level 1).

prime numbers (see, for example, [35], and also the chapter on the Mersenne primes in [36]). But no: experience shows that the estimate of the exponent α diminishes when the the bound x grows. And let us not forget that for the time being we are unable to prove even that there are infinitely many projective primes, to say nothing of their asymptotic behavior.

6.6. When the exponent *n* grows. We considered the prime exponents $n \le 100$ and counted the number of primes $p \le 10^6$ (omitting prime powers) such that $m = (p^n - 1)/(p - 1)$ is also prime. The results are presented in Table 3: *N* denotes the number of primes *p* with the above property, and max *p* is the largest $p \le 10^6$ which, for a given *n*, produces a prime value of *m*. The total number of primes up to one million is 78 498. The proportion of "good" primes thus varies (in our table) from 6% (for n = 3) to 0.5% (for n = 83). There is a general tendency for this proportion to decrease, but without any apparent regularity.

We do not present the corresponding projective primes *m* since their decimal representations are too long: for example, the number *m* corresponding to the last cell of the table, namely, $m = (998471^{97} - 1)/998470$, has 576 digits.

The following conjecture seems quite reasonable:

Conjecture 6.5 (Projective primes for a fixed *n*). For any fixed prime $n \ge 3$ there are infinitely many prime values $m = 1 + p + p^2 + ... + p^{n-1}$, where *p* ranges over all prime numbers.

n	N	max p	n	N	max p	n	N	max p	n	N	max p
3	4684	999773	19	2933	999067	43	1119	999961	71	848	999 907
5	4034	999653	23	1150	999 287	47	1212	999 491	73	577	999 307
7	4436	999961	29	1032	998111	53	694	999 007	79	689	996811
11	2243	999631	31	1980	997 463	59	1106	999 953	83	390	993 557
13	2658	999 863	37	1285	999 269	61	913	999763	89	430	995 339
17	2527	999 287	41	862	999233	67	821	999727	97	571	998 471

TABLE 3. For a given prime exponent $n \le 100$, the number N shows how many primes $p \le 10^6$ there are such that $m = (p^n - 1)/(p - 1)$ is also prime. The column "max p" shows the largest such p.

р	estimation (10)	true number	exponents <i>n</i>
3	13.832	12	3, 7, 13, 71, 103, 541, 1091, 1367, 1627, 4177, 9011, 9551
5	8.669	11	3, 7, 11, 13, 47, 127, 149, 181, 619, 929, 3407
7	6.796	5	5, 13, 131, 149, 1699
11	5.136	9	17, 19, 73, 139, 907, 1907, 2029, 4801, 5153
13	4.675	9	5, 7, 137, 283, 883, 991, 1021, 1193, 3671
17	4.052	7	5, 7, 11, 47, 71, 419, 4799

TABLE 4. The last column gives the list of exponents $n \le 10^4$ such that the number $m = (p^n - 1)/(p - 1)$ is prime. The number of such exponents is given in column 3, while the estimation of this number by formula (10) is presented in column 2.

6.7. **Projective primes with a fixed** p. What if we fix p and allow n to tend to infinity (taking only prime values), as in Section 4.4? Here the computational evidence, presented in Table 4, is less convincing, which is not surprising given the small number of primes m involved.

For the primes p = 3, 5, 7, 11, 13 and 17, the second column of Table 4 gives the estimates based on the first expression in (10) for the number of primes $m = (p^n - 1)/(p - 1)$ with $n \le x = 10^4$. The third column gives the true figures, found by a computer search, and the relevant exponents nare listed in the fourth column.

We see that the estimation (10) of the number of exponents has a reasonably good correspondence with their actual number. However, we have a feeling that the above data are not entirely convincing. For example, there are only three exponents in the table, out of 53, which are greater than 5000. Also, for p = 3, the next "good" exponent, after the one given in the table, is rather far away: n = 36913. Therefore, in this case we prefer to formulate not a conjecture but a question:

Question 6.6 (Generalized Mersenne). Let *p* be a prime. Do there exist infinitely many values of *n* such that the number $m = (p^n - 1)/(p - 1)$ is prime?

6.8. **Prime powers once again.** We now consider fixed prime powers $q = p^e$ with $e \ge 2$, as $n \to \infty$. By Remark 4.5, apart from the example $q = 2^2$ and m = 5 we can ignore the case e = 2. An extensive search of prime powers producing projective primes has given the following results:

p = 2: the search for $q = 2^e$, $e \ge 2$, producing projective primes, up to $q \le 10^{60}$, gives eight solutions:

• There are four solutions for which m = 1 + q is a Fermat prime:

$$(q, n) = (2^2, 2), (2^4, 2), (2^8, 2), (2^{16}, 2)$$

(the other known Fermat prime $m = 1 + 2^1 = 3$ does not correspond to $e \ge 2$).

- There are three relatively small solutions: $(q, n) = (2^3, 3), (2^7, 7), (2^9, 3).$
- A rather unexpected solution is $(q, n) = (2^{59}, 59)$. The corresponding projective prime

$$m = 1 + 2^{59} + 2^{118} + \dots + 2^{59 \cdot 58}$$

has 1031 digits. It is generally believed that there are only five Fermat primes. However, this example prevents us from conjecturing that there are only finitely many powers of 2 which yield projective primes.

p = 3: the search for $q = 3^e$, $e \ge 3$, producing projective primes, up to $q \le 10^{60}$, gives only one solution: $(q, n) = (3^3, 3)$, $m = 1 + 27 + 27^2 = 757$. We are inclined to believe that for p = 3 this solution is unique.

 $q \le 10^{15}$: the total search for all prime powers $q \le 10^{15}$ with $e \ge 3$ producing projective primes gives 337 solutions. Only eight of them have the exponent e > 3, namely,

$$(q, n) = (5^7, 7), (11^9, 3), (43^5, 5), (67^7, 7), (167^5, 5), (313^5, 5), (509^5, 5), (859^5, 5).$$

For all the other 329 solutions q is the cube of a prime.

 $q = p^3 \le 10^{18}$: the total search for cubes of primes up to 10^{18} reveals 2121 solutions, the largest one being p = 999953, $q = p^3 = 999859006626896177$, and

$$m = 1 + q + q^2 = 999718033132923614193697947364111507$$

The following conjecture seems to be very plausible:

Conjecture 6.7 (Cubes of primes). *There are infinitely many values of* $q = p^3$, *with p being prime, such that* $m = 1 + q + q^2$ *is prime.*

Since we did not find any examples where the same prime power q yields more than one projective prime, we ask:

Question 6.8 (Generalized Mersenne). Does there exist a proper prime power q such that the number $m = (q^n - 1)/(q - 1)$ is prime for more than one value of n?

Of course, Table 4 in Section 6.7 gives a number of examples of this phenomenon where q = p is prime.

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7. STOCHASTIC BEHAVIOR OF PROJECTIVE PRIMES

In this section, we present a number of observations concerning the stochastic behavior of the (numbers of) projective primes. Our feeling is that this subject, while being excitingly interesting, is not yet ready for a profound statistical analysis. However, we would like to share our observations with the community of specialists in probabilistic number theory in the hope that they may clarify certain points of our study.

We deal here exclusively with projective primes of the form $m = 1 + p + p^2$, where p is prime.

7.1. Local estimates. Let *a*, *b* be two integers, a < b. Then, according to (13-14), the primes $p \in [a, b]$ should create, approximately,

(15)
$$C \cdot \left(\frac{b}{\ln(b)^{\alpha}} - \frac{a}{\ln(a)^{\alpha}}\right)$$

projective primes of the type $m = 1 + p + p^2$, where C = 1.162280117 and $\alpha = 2.104419156$. In Figure 3, left, we consider the primes $p \le 10^9$. We subdivide this range into 10^4 segments $S_i = [(i - 1) \cdot 10^5, i \cdot 10^5]$ $(i = 1, 2, ..., 10^4)$ of equal size 10^5 . Horizontally, we mark the order number *i* of a segment (from 1 to 10^4). For each of these segments, we divide estimate (15) by the true number of projective primes $m = 1 + p + p^2$ for $p \in S_i$; this ratio is the ordinate of the corresponding point in the picture. The picture thus contains 10^4 points.

On the right of Figure 3, the construction is similar, but now the range considered is $p \le 10^{10}$, and the length of each of the 10^4 segments is 10^6 .



FIGURE 3. Both pictures contain 10^4 points. Each point corresponds to a segment in \mathbb{N} of length 10^5 (left) or 10^6 (right). The abscissa of a point is the order number of the corresponding segment. The ordinate is the ratio of the estimate (15) to the true number of projective primes $m = 1 + p + p^2$ generated by primes p in this segment.

We may make the following observations.

- There are wild fluctuations in the ratio. Obviously, they are due not to estimate (15) itself but to the fluctuations in the numbers of projective primes $m = 1 + p + p^2$ with p belonging to the corresponding segments.
- Comparing the vertical scales shows that the right-hand band of points is narrower than the left-hand one. This is natural since considering larger segments leads to smoothing the fluctuations.
- The interesting fact is, however, that the variations in both pictures do not diminish when we let *i* increase. We may even say that they increase.

7.2. **Histogram.** Let us take the right-hand picture of Figure 3. The minimum value of the ordinate (i. e., of the ratio) in this picture is $r_{min} = 0.9217$, the maximum is $r_{max} = 1.1244$. We subdivide the segment $[r_{min}, r_{max}]$ into 100 parts and count the number of points whose ordinates belong to each part. The resulting histogram is shown in Figure 4, left.

The mean of this distribution is 46.66, the standard deviation is 13.66. On the right of the same figure we show the density of the normal distribution with the same parameters. Note that the height of the left picture, which is 300 points out of 10 000, corresponds well to the height of the density, which is approximately 0.03.

The resemblance of the two graphs is visible. We leave it to the specialists to use, if necessary, more sophisticated statistical tools.



FIGURE 4. On the left: the histogram of the distribution of heights of the points in the right-hand picture of Figure 3. On the right: the density of the normal distribution with the same mean 46.66 and standard deviation 13.66.

7.3. Conclusion. Our main aim in this note has been to give heuristic and computational evidence that there are infinitely many projective primes, especially in the simplest and apparently most abundant case, where n = 3 and q is prime. We will not pursue these speculations further and will

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leave the question of more exact estimates of the number and distribution of projective primes to the community of experts in probabilistic number theory. (Our own backgrounds and motivation for this investigation lie in the areas of dessins d'enfants and permutation groups.)

Acknowledgements. We are greatly indebted to Yuri Bilu who acquainted us with the Bunyakovsky conjecture, which became a crucial point of our study, and to Peter Cameron, Robert Guralnick and Cheryl Praeger for some very useful comments. Jean Bétréma helped us with some computations which were too heavy for our Maple package on a laptop computer. Alexander Zvonkin was partially supported by the ANR project COMBINÉ (ANR-19-CE48-0011).

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⁵Numerous publications, as well as the English Wikipedia (visited on June 30, 2020), give the following wrong title for Bunyakovsky's paper: "Nouveaux théorèmes relatifs à la distinction des nombres premiers et à la décomposition des entiers en facteurs". According to the French Wikipedia (see [7]), an article with this title does indeed exist, but it was published in 1840 and not in 1857, and it does not discuss the conjecture in question. The reader may also consult the original paper reproduced in the Google archive.

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